



Xsfmm Family of Attached Matrix Extensions

Version 0.6.6

Copyright © by SiFive, Inc. All rights reserved.

Xsfmm Family of Attached Matrix Extensions

Proprietary Notice

Copyright © 2026 by SiFive, Inc. All rights reserved.

Information in this document is provided “as is,” with all faults.

SiFive expressly disclaims all warranties, representations, and conditions of any kind, whether express or implied, including, but not limited to, the implied warranties or conditions of merchantability, fitness for a particular purpose and non-infringement.

SiFive does not assume any liability arising out of the application or use of any product or circuit, and specifically disclaims any and all liability, including without limitation indirect, incidental, special, exemplary, or consequential damages.

SiFive reserves the right to make changes without further notice to any products herein.

Revision History

Version	Date	Changes
0.6.6	Jan 8, 2026	Clarify illegal instruction exception for reserved altfmt/vtype
0.6.5	Sep 2, 2025	Added implementation note for non-zero vstart for vtdiscard
0.6.4	Jun 6, 2025	Added implementation notes for non-zero vstart
0.6.3	May 13, 2025	Sync supported SEW/TWIDEN/ALTFMT table with Zvma spec.
0.6.2	Apr 29, 2025	Remove deprecated Xsfmm32a4i
0.6.1	Mar 13, 2025	Xsfmmbase depends on Zve32x. FP depends on Zve32f/Zve64d.
0.6	May 31, 2024	ISA INCOMPATIBLE Use BF16 vector opcodes; split operand extensions
0.5	May 10, 2024	Replace FP8 P3109 formats with OCP
0.4	Apr 26, 2024	Rework VRG, configuration instructions, context control

Version	Date	Changes
0.3	Apr 23, 2024	Remove tk and tm CSRs.
0.2	Mar 26, 2024	Ext. naming, vset* naming, add missing descriptions
0.1	Mar 15, 2024	Initial Revision

Contents

List of Figures	3
1 Overview	4
1.1 Matrix Tile State	5
1.1.1 Tile Punning	6
1.2 Additions to vtype CSR	8
1.3 Matrix-Multiply Operations	9
1.4 Matrix Configuration Instructions	11
1.4.1 sf.vsettm, sf.vsettn, and sf.vsettk instructions	13
1.4.2 Supported SEW, TWIDEN, and KMAX values	13
1.5 Tile Subset Specifier (TSS)	14
1.6 Load/Store Tile Subset to Memory	14
1.7 Move Tile Subset between Tile State and Vector Registers	16
1.8 Matrix Arithmetic Instructions	16
1.8.1 Matrix Arithmetic Instruction Encoding	17
1.9 Matrix Tile-Zero Instruction	18
1.10 Context Status	18
1.10.1 Xsfmm Context Status in mstatus	18
1.10.2 Xsfmm Context Status in vsstatus	19
1.10.3 Xsfmm Context Discard Instruction	19
1.11 C Intrinsics	20
1.11.1 Matrix Configuration Instructions	20
1.11.2 Load/Store Tile Subset to Memory	21
1.11.3 Move Tile Subset between Tile State and Vector Registers	22
1.11.4 Matrix Arithmetic Instructions	23
1.11.5 Matrix Tile-Zero Instruction	24
1.11.6 Context Status	24
2 ABI Attributes	25

2.1	__riscv_in("xsfmm")	26
2.2	__riscv_out("xsfmm")	26
2.3	__riscv_inout("xsfmm")	27
2.4	__riscv_new("xsfmm")	27
2.5	__riscv_preserves("xsfmm").....	28
2.6	Example Scenario.....	28
2.7	Analogy Using General Program	28

List of Figures

Figure 1	Tile Punning Aliasing.....	7
Figure 2	A or B Matrix Operand	9
Figure 3	Matrix Operand Layout in Vector Register Group	10

Overview

The Xsfmm family of extensions supports high-throughput matrix-multiplication computation. These extensions add new architectural state to each hart in the form of square matrix tiles. They provide instructions that perform tall-skinny by short-wide matrix multiplications, i.e. outer products or wide outer products. These instructions source their operands from vector registers and accumulate their results into the new square matrix tiles. The inputs' data width is specified by SEW, while the accumulators' data width is given by $TEW = SEW * TWIDEN$, where TWIDEN can be 1, 2, or 4, depending on the instruction.

The new computational instructions perform the operation $C[M,N] += A[K,M]^T * B[K,N]$, where C is stored in the tile state. Each of the input matrices A and B is held in a set of vector registers, where each row of the matrix is held in a separate vector register group and the number of rows is given by K. The inner dimension K can be a maximum (KMAX) of 1, 2, or 4, depending on the data type of A and B.

In addition, the Xsfmm extensions provide instructions to move rows and columns of the matrix arrays, either to and from the vector registers or to and from memory.

The Xsfmmbase extension provides the new architectural state and the new configuration, move, load, store, and context-manipulation instructions. These instructions support tile EEW (TEW) values corresponding to all supported SEW settings (8..ELEN).

The Xsfmm extensions build upon the vector extensions. The vector registers serve as a data source for matrix operands, and the existing vector instructions are used for pre- and post-processing of matrices. Hence, the Xsfmmbase extension requires the Zve32x extension.

The Xsfmm extensions are compatible with the RV32E, RV32I, RV64E, and RV64I base ISAs, but it is expected that RV64I will be by far the most common base-ISA choice.

The following table lists the computational extensions in the Xsfmm family.

Extension	Dependencies	Multiplicand Types	Accumulator Type
Xsfmm32a8i	Xsfmmbase	[U]Int8	Int32
Xsfmm32a8f	Xsfmmbase_Zve32f	OCP MX FP8	FP32
Xsfmm32a16f	Xsfmmbase_Zve32f	BF16, FP16	FP32
Xsfmm32a32f	Xsfmmbase_Zve32f	FP32	FP32
Xsfmm64a64f	Xsfmmbase_Zve64d	FP64	FP64

An additional extension, Xsfmm32a, is defined as a shorthand for the combination of Xsfmm32a8i, Xsfmm32a16f, and Xsfmm32a32f.

1.1 Matrix Tile State

The new matrix state is organized into *tiles*, where a tile is a separately addressable two-dimensional square matrix. The fixed matrix state is viewed as a different number of tiles, depending on TEW.

The ISA is optimized for the case where TEW=32, as this is the most common accumulator size. When TEW=32, the tiles have dimensions TExTE elements, where TE is a parameter of the implementation. TE is constrained to be a power of 2, $VLEN/4 \geq TE \geq 4$. The upper bound is set by the requirement that a tile row or column must fit within a single vector register group ($VLEN \times 8$ bits). Implementations may choose less than the maximum TE to reduce implementation cost. The configuration instructions, described below, support writing of TE-agnostic code.

The ISA fixes the number of TEW=32 tiles to four (mt0, mt4, mt8, mt12) to support all cases at reasonable performance. The matrix state needs to hold at least one accumulator as well as transposes of one or both inputs to support efficient execution of all matrix operations. Also, providing more separately addressable tiles allows more flexible usage of tile accumulators enabling optimizations such as register blocking of the operands, reducing memory traffic.

Tile EEW	#Tiles	Tile Dimensions	Tile Specifiers
8	16	TExTE	mt0,mt1, ..., mt14,mt15
16	8	TExTE	mt0,mt2,mt4,mt6,mt8,mt10,mt12,mt14
32	4	TExTE	mt0, mt4, mt8, mt12
64	8	(TE/2)x(TE/2)	mt0,mt2,mt4,mt6,mt8,mt10,mt12,mt14

For smaller TEW<32, the tile dimensions (TExTE) are kept the same as TEW=32. Every halving of TEW halves the state required for a single tile and so the number of tiles increases by a factor of two. The storage of each TEW=32 tile is divided to form two TEW=16 tiles (e.g. mt0 for TEW=32 forms mt0 and mt2 for TEW=16), and the storage of each TEW=16 tile is divided to form two TEW=8 tiles (e.g. mt0 for TEW=16 forms mt0 and mt1 for TEW=8).

For TEW=64, the tile dimensions are reduced by a factor of 2 in each dimension (i.e. tiles are TE/2 x TE/2 elements), but the number of tiles increases to 8 as the tile size shrinks quadratically while the element width only doubles. The storage of each TEW=32 tile is divided to form two TEW=64 tiles (e.g. mt0 for TEW=32 forms mt0 and mt2 for TEW=64).

Additional ISA extensions describe the tile dimensions. The XsfmmNt extension indicates $TE \geq N$, e.g. Xsfmm32t indicates TE is at least 32.

Note

The accumulator array dimensions (TE) may be determined dynamically with the `vset[i]vl[i]` instructions to enable writing TE-agnostic software.

1.1.1 Tile Punning

The tile state is reshaped for each supported TEW which results in tiles being combined to support wider accumulators. For example, 4 TEW=8 tiles (mt0..3) are combined to support a single TEW=32 tile (mt0).

Given the layout of elements for each TEW in the tile, there is an aliasing of the differently sized elements in the combined space. The aliasing is a repetitive pattern of 4x4 (2x2 TEW=64) sub-matrices of elements as illustrated in the following figure.



Figure 1: *Tile Punning Aliasing*

Viewing the tile state as a linearized buffer of $16 * TE * TE$ bytes, an element's position can be determined with the following algorithm.

Given,

```
tile:      tile
tew:       element size
row,col:   element indices
```

```
case (tew) {
8: {
    ptile = tile
    minor_offset = (row % 4) * 4 + (col % 4)
    major_offset = (row / 4) * (TE / 4) + (col / 4)
}
16: {
    ptile = tile + ((row & 2) >> 1)
    minor_offset = (row % 2) * 4 + (col % 2) * 2 + ((col / 2) % 2) * 8
    major_offset = (row / 4) * (TE / 4) + (col / 4)
}
32: {
    ptile = tile + (row & 2) + ((col & 2) >> 1)
}
```

```
    minor_offset = (row % 2) * 8 + (col % 2) * 4
    major_offset = (row / 4) * (TE / 4) + (col / 4)
}
64: {
    ptile = tile + (row & 1)
    minor_offset = (col % 2) * 8
    major_offset = (row / 2) * (TE / 4) + (col / 2)
}
}

offset = (ptile * TE * TE) + (major_offset * 16) + minor_offset
```

1.2 Additions to vtype CSR

The `tm`, `tk`, and `vtwiden` fields are added to the existing `vtype` CSR in previously reserved space.

bits	field
XLEN-1	vill
30	reserved
29:16	tm[13:0]
15:14	reserved
13:11	tk[2:0]
10:9	vtwiden[1:0]
8	altfmt
7	vma
6	vta
5:3	vsew[2:0]
2:0	vlmul[2:0]

`tm` can hold values from 0-TE, inclusive.

`tk` can hold values from 0-4, inclusive.

`tn` is given by the value in the `v1` register.

When `altfmt=1` and `SEW=16`, all matrix-multiply instructions become reserved, except for the following, which is redefined to use the BF16 format for any operand that would otherwise have used the FP16 format:

- `sf.mm.f.f` (when the `Xsfmtm32a16f` extension is implemented)

Selecting `altfmt=1` for other `SEW` settings is reserved.

Note

SiFive implementations set `vill` in `vtype` when a reserved combination of `altfmt` and `SEW` is selected.

If `vtwiden==0`, the matrix unit is not configured.

1.3 Matrix-Multiply Operations

The general form of the matrix-multiply operations is:

$$C[tm,tn] += A[tk,tm]^T * B[tk,tn]$$

where the inputs A and B are each sourced from a vector register group, while the accumulator C is held in a matrix tile. A, B, and C are all specified using register specifier fields in the instruction encoding.

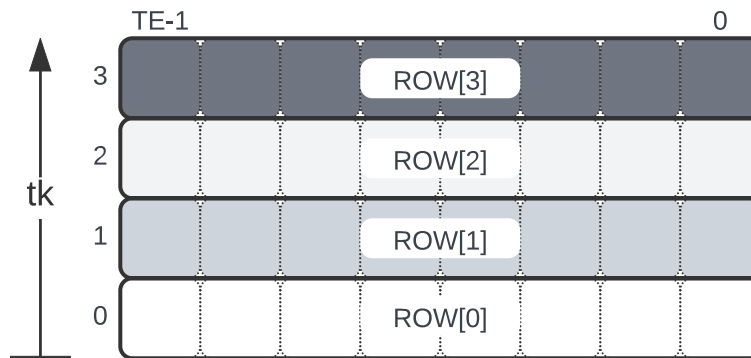


Figure 2: A or B Matrix Operand

The values of `tm`, `tn`, and `tk` are obtained from CSRs, with `tn` being an alias for the existing `v1` register. `tm` and `tk` are fields in the existing `vtype` register. The new configuration instructions described below allow the values of `tm`, `tn`, and `tk` to be set in an implementation-agnostic manner from the application matrix dimensions.

Each vector register group input has `tk` rows, where each row is held separately in one or more vector registers, with the total number of vector registers constrained to fit in one eight-register group. Each row fits within a vector register group (i.e. an aligned power-of-2 number of vector registers) and the collection of rows also fits in a (potentially larger) vector register group. The value of `tk` sets the dimension of the dot products performed for each accumulator element and the maximum value of `tk` (`KMAX`) depends on the element widths, with a maximum value of 4.

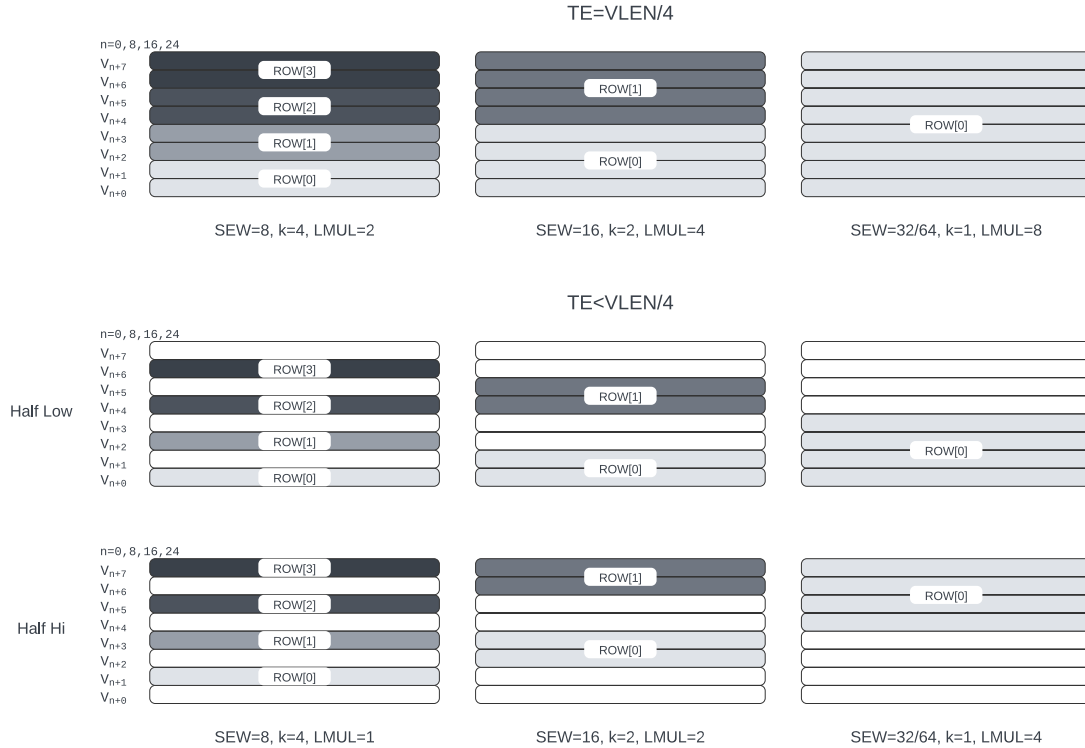


Figure 3: Matrix Operand Layout in Vector Register Group

This format allows existing vector unit-stride loads to be used to load each row of an input from a memory matrix held in row-major format.

Successive rows are separated by $(8/KMAX)$ vector registers. For example, for vector register group specifier $v8$, and assuming $KMAX=4$ and $LMUL=2$, then $v8..v9$ holds one matrix row, $v10..v11$ holds the next row, $v12..v13$ the next, and $v14..v15$ the last. But when $KMAX=4$ and $LMUL=1$, then $v8$, $v10$, $v12$, and $v14$ hold the successive rows of the matrix: the separation between rows is the same, but the odd-numbered vector registers are skipped.

Vector register specifiers must be divisible by $LMUL$. For example, when $KMAX=4$ and $LMUL=2$, $v8$ is valid but $v9$ is not. But when $KMAX=4$ and $LMUL=1$, both are valid: the former represents the even-numbered registers $v8..v14$, whereas the latter represents the odd-numbered registers $v9..v15$.

Note

These definitions facilitate writing $LMUL$ -agnostic software, but they do not preclude $LMUL$ -conscious software from using all of the vector registers.

Vector register specifiers, when taken modulo 8, must be less than 8/KMAX. For example, when KMAX=4, v10 is invalid.

Note

This constraint prevents register groups from spanning EMUL=8 register-group boundaries, simplifying vector-unit implementations.

1.4 Matrix Configuration Instructions

Application matrices have varied sizes, and implementations will vary in the TE dimensions they support. The matrix configuration instructions allow a single binary matrix-multiply routine to work for any size application matrix multiply on any sized matrix-multiply implementation.

Note

More optimized routines are possible when microarchitectural parameters are known to software.

The actual size of a matrix-multiply operation is determined by values in two CSRs, tk and tm are fields in the vtype register, tn is the same as the vl register. These CSRs are set using configuration instructions.

Encoding of vector and matrix configuration instructions

31	30	25	24	20	19	15	14	12	11	7	6	0
0		zimm[10:0]				rs1				rd	1010111	vsetvli
1		1		zimm[9:0]				uimm[4:0]				vsetivli
1		000000				rs2				rd	1010111	vsetvl
1		000010				00000				rs1/ATN	1 1 1	rd 1010111 sf.vsettn
1		000010				00001				rs1/ATM	1 1 1	rd 1010111 sf.vsettm
1		000010				00010				rs1/ATK	1 1 1	rd 1010111 sf.vsettk
1		6				5				5		7

ATM is application matrix m dimension

ATN is application matrix n dimension

ATK is application matrix k dimension

The existing vset[i]vl[i] instructions are used to configure the hart to execute Xsfmtm instructions. These instructions behave as defined in the V extension, but their behavior is modified when the requested vtype value has vtwidth != 0. In particular, the vset and vtwidth fields are set as requested, but other fields in vtype are determined using special rules described below. AVL is determined as described in the V spec, but the new vl is determined

using a special rule described below. `rd` is written with the new value of `v1`, as usual. Although `vill` is set under different rules, as described below, the behavior remains that if `vill` is set, then the other bits in `vtype`, as well as `v1`, are set to zero.

`vset[i]v1[i]` behavior when requested `vtype.vtwiden != 0`

```
TEW = SEW * TWIDEN
ETE = TEW < 64 ? TE : TE/2 // Effective number of elements along tile edge
EVE = VLEN/SEW             // Effective number of elements in a vector register

LMUL = min(8/KMAX, 8/TWIDEN, ceil(ETE/EVE))
      ^      ^      ^
      |      |      +--matrix engine size constraint
      |      +--matrix row/col must fit in vector register group
      +--leave space for K dimension

tn/vl = min(AVL, LMUL*EVE, ETE) // AVL is as defined in V spec
tm = min(ATM, LMUL*EVE, ETE)   // ATM comes from requested vtype.tm
tk = min(ATK, KMAX)            // ATK comes from requested vtype.tk
altfmt = ALTFMT                // ALTFMT comes from requested vtype.altfmt
vma = 1
vta = 1

vill = (TEW > ELEN) || (any reserved field in vtype != 0)

rd = tn/vl
```

A new assembler pseudo-op, `sf.vsettn` ("set `tn` and type") is provided to simplify using the `vsetvli` instruction for Xsfrm purposes.

```
// vsetvli rd, rs1, (vsew << 3) + (alt << 8) + (vtwiden << 9)
sf.vsettn rd, rs1, eX, wY
```

The valid ``wY`` and ``eX`` options are as follows.

vtwiden	assembler syntax
01	w1
10	w2
11	w4

vsew	alt	assembler syntax
00	0	e8
01	0	e16
01	1	e16alt
10	0	e32
11	0	e64

An ``eX, wY`` combination is valid if and only if $X * Y \leq ELEN$.

Note

`w0` is not supported—this behavior is achieved by the existing `vset[i]v1[i]` instructions—and assemblers should reject it.

1.4.1 sf.vsettm, sf.vsettn, and sf.vsettk instructions

These instructions set `tm`, `tn`, or `tk` to the value in `rs1`, constrained by the current matrix configuration. The new value of `tm`, `tn`, or `tk` is also written to the scalar destination register `rd`. If the matrix unit has not been configured, i.e. `vtwiden=0` in `vtype`, the instruction sets `vll` in `vtype`.

These instructions are included in the `Xsfrmmbase` extension.

```
sf.vsettm rd, rs1 // rd = tm = min(rs1=ATM, LMUL*EVE, ETE)
```

The `sf.vsettm` instruction sets `tm` to be the lesser of `rs1` (`ATM`), `LMUL*EVE`, and `ETE`.

```
sf.vsettn rd, rs1 // rd = tn/vl = min(rs1=ATN, LMUL*EVE, ETE)
```

The `sf.vsettn` instruction sets `tn` to be the lesser of `rs1` (`ATN`), `LMUL*EVE`, and `ETE`.

```
sf.vsettk rd, rs1 // rd = tk = min(rs1=ATK, KMAX)
```

The `sf.vsettk` instruction sets `tk` to the lesser of `rs1` (`ATK`) and `KMAX`.

1.4.2 Supported SEW, TWIDEN, and KMAX values

The inner dimension specified by `tk` is used to support fused dot-product operations on narrower datatypes to increase efficiency. The ISA fixes maximum `K` values (`KMAX`) of matrix-multiply instructions according to `SEW` and `TWIDEN` to simplify ISA specification of fused arithmetic operators.

Note

Unlike other matrix accelerator architectures that support fusion of narrower datatypes by flattening `K` into both the `M` and `N` dimensions, this approach allows software to use standard matrix formats.

TEW=8

SEW	TWIDEN	KMAX	
8	1	4	(fused 4-element dot product)

TEW=16

SEW	TWIDEN	KMAX	
8	2	4	(fused 4-element dot product)
16	1	2	(fused 2-element dot product)

TEW=32

SEW	TWIDEN	KMAX	
8	4	4	(fused 4-element dot product)
16	2	2	(fused 2-element dot product)

32 1 1

TEW=64

SEW	TWIDEN	KMAX
16	4	2 (fused 2-element dot product)
32	2	1
64	1	1

1.5 Tile Subset Specifier (TSS)

For operations that move rows or columns between the tiles and vector registers or memory, a scalar value is used to encode the set of tile elements that are accessed and the pattern by which they are accessed.

bits	meaning
XLEN-1:31	reserved
30:27	tile specifier (0-15)
26:24	pattern (0=row, 1=col, >1 reserved)
23:0	index within pattern (e.g. row/col index)

When there are fewer than 16 tiles due to the specified TEW, the $\log_2(16 / \text{\#Tiles})$ LSBs of the tile-specifier field are ignored, such that this field always refers to a valid tile.

The 3-bit field for pattern currently only has two encodings defined, row and column. The pattern field defines how the index field is interpreted. For the currently defined row and column patterns, the index field is treated as an unsigned integer providing a one-dimensional index of the row or column. Row and column indices greater than or equal to ETE are reserved.

Note

We expect most implementations will ignore reserved bits in TSS, taking pattern mod 2, and taking the index within the pattern mod ETE.

1.6 Load/Store Tile Subset to Memory

These instructions transfer tile subsets between the tile state and memory. Implementations must provide vector loads and stores with EEWs corresponding to all supported SEW settings (8..ELEN). These instructions are included in the Xsfgmmbase extension.

```
// Load tile subset from memory
sf.vlte8  rs2, (rs1) // rs1 is memory address, rs2 is TSS
sf.vlte16 rs2, (rs1) // rs1 is memory address, rs2 is TSS
sf.vlte32 rs2, (rs1) // rs1 is memory address, rs2 is TSS
sf.vlte64 rs2, (rs1) // rs1 is memory address, rs2 is TSS

// Store tile subset to memory
sf.vste8  rs2, (rs1) // rs1 is memory address, rs2 is TSS
```

```

sf.vste16 rs2, (rs1) // rs1 is memory address, rs2 is TSS
sf.vste32 rs2, (rs1) // rs1 is memory address, rs2 is TSS
sf.vste64 rs2, (rs1) // rs1 is memory address, rs2 is TSS

31 29 28 27 26 25 24    20 19    15 14 12 11    7 6    0
  nf  mew  mop  vm    rs2    rs1    width vd/vs3 opcode
e e e  1  0  0  1    rs2    rs1    1 1 1    0    0000111 // Load tile
e e e  1  0  0  1    rs2    rs1    1 1 1    0    0100111 // Store tile

eee
000 8
001 16
010 32
011 64
1xx reserved

```

The integer register specified by `rs1` holds the base address in memory. The integer register specified by `rs2` holds the tile subset specifier. These instructions do not support masks. EEW is statically encoded in the instruction.

These instructions always load and store the tile subset to a contiguous region of memory. For the currently defined row and column patterns, the elements of the row or column are stored in increasing index order in memory.

For these instructions, the body elements are those with indices in the range $[vstart, \min(vl, ETE) - 1]$, and the tail elements are those with indices $\geq \min(vl, ETE)$.

Note

Since $vl=tn$ holds the number of columns in the accumulator, i.e. the length of a row, these instructions are oriented towards loading and storing accumulator rows. When loading and storing accumulator columns, additional `vsetvli` or `sf.vsettn` instructions may be needed to program $vl=tn$ to hold the number of rows, i.e. the length of a column, instead.

For the purposes of exception handling, these instructions are considered to access memory in ascending order beginning at the byte determined by a function of the base address in `rs1`, the `vstart` value, and the instruction EEW. No memory is accessed below the elements specified by the `vstart` value.

The matrix extension follows the precise exception model of the vector extension, allowing memory to be updated past the element indicated by `vstart` on a trap.

Misaligned element addresses are handled the same as for unit-stride vector loads and stores.

1.7 Move Tile Subset between Tile State and Vector Registers

Tile subset operations transfer elements between a tile and a vector register group. These instructions are included in the Xsfmmbase extension.

```
// Move tile subset to vector register group, using TEW=SEW
sf.vtmv.v.t vd, rs1 // vd is destination vector register, rs1 is TSS of source

// Move vector register group to tile subset, using TEW=SEW
sf.vtmv.t.v rs1, vs2 // vs2 is source vector register, rs1 is TSS of destination
```

31		26	25	24		20	19		15	14	12	11		7	6		0
	funct6		vm		rs2		rs1		funct3	vd/vs3	opcode						
	010000		1		11111		rs1		1 1 0	vd	1010111						// sf.vtmv.v.t
	010111		1		vs2		rs1		1 1 0	0	1010111						// sf.vtmv.t.v

For these instructions, the body elements are those with indices in the range $[vstart, \min(vl, ETE) - 1]$, and the tail elements are those with indices $\geq \min(vl, ETE)$.

Note

SiFive implementations raise an illegal-instruction exception when executing `sf.vtmv` instructions with a nonzero `vstart`.

1.8 Matrix Arithmetic Instructions

These instructions all have the form $C += A^T * B$, where C is held in a tile and A and B are supplied by vector register groups.

Unsupported `vtype` settings are handled the same as for vector instructions. An illegal-instruction exception is raised if `vstart` is nonzero.

```
vtype holds SEW and TWIDEN
vs2 specifies A vector register group
vs1 specifies B vector register group
mtd is C destination tile specifier
```

# IEEE FP defined for:	SEW	TWIDEN	ALTFMT
#	16	2	0 FP16 (Xsfmm32a16f)
#	16	2	1 BF16 (Xsfmm32a16f)
#	32	1	0 FP32 (Xsfmm32a32f)
#	64	1	0 FP64 (Xsfmm64a64f)
#			

```
sf.mm.f.f mtd, vs2, vs1 # IEEE FP matmuls (FP16/BF16, FP32, FP64)
```

```
# OCP FP8 defined for SEW=8 TWIDEN=4, accumulate into FP32
```

```
sf.mm.e5m2.e5m2 mtd, vs2, vs1 # vs2=E5M2, vs1=E5M2 (Xsfmm32a8f)
sf.mm.e5m2.e4m3 mtd, vs2, vs1 # vs2=E5M2, vs1=E4M3 (Xsfmm32a8f)
sf.mm.e4m3.e5m2 mtd, vs2, vs1 # vs2=E4M3, vs1=E5M2 (Xsfmm32a8f)
sf.mm.e4m3.e4m3 mtd, vs2, vs1 # vs2=E4M3, vs1=E4M3 (Xsfmm32a8f)

# INT8 defined for SEW=8 TWIDEN=4, accumulate into INT32
sf.mm.u.u mtd, vs2, vs1 # unsigned vs2(A) unsigned vs1(B) (Xsfmm32a8i)
sf.mm.s.u mtd, vs2, vs1 # signed vs2(A) unsigned vs1(B) (Xsfmm32a8i)
sf.mm.u.s mtd, vs2, vs1 # unsigned vs2(A) signed vs1(B) (Xsfmm32a8i)
sf.mm.s.s mtd, vs2, vs1 # signed vs2(A) signed vs1(B) (Xsfmm32a8i)
```

Note

There are numerous FP8 formats being developed and standardized, but the OCP MX formats appear to be gaining the most traction. Additional formats would ideally be supported with explicit conversions to BF16 or FP16 in the vector unit, as supporting all combinations natively would consume substantial opcode space. But native support for additional formats remains possible if sufficient demand arises.

These instructions use the values of t_m , t_n , and t_k to determine what computations to perform. If any of these values are zero, then no computation is performed and no state is updated.

Any tile elements outside the range $[0, t_m-1] \times [0, t_n-1]$ are considered part of the tail and are handled using a tail-agnostic policy.

The invalid and overflow scalar floating-point exception flags are updated by the results of floating-point matrix operations. The inexact, underflow, and divide-by-zero flags are not updated, nor is the fixed-point saturation flag.

For $SEW \geq 32$, floating-point products are computed in IEEE 754 arithmetic then rounded to a TEW-bit value, using the rounding mode specified by frm . They are then added to the accumulator in IEEE 754 arithmetic, using the same rounding mode.

For $SEW \leq 16$, floating-point products are bulk-normalized, rounded to odd, accumulated in fixed-point arithmetic, then converted to FP32, rounding to odd. Finally, they are added to the accumulator in IEEE 754 arithmetic, using the rounding mode specified by frm .

1.8.1 Matrix Arithmetic Instruction Encoding

The matrix arithmetic instructions are encoded in the second vector major opcode, OP-VE.

The rd field encodes the accumulator tile, represented with t in the diagram below. The tile-selection bits are allocated from the MSB of the rd field (e.g. TEW=32 operations have 4 tiles, requiring 2 bits, $t[3:2]$). Tile-selection bits not present in the instruction encoding are defined to be zero. If t encodes an invalid tile specifier, i.e. $t \bmod (16 / \#Tiles)$ is nonzero, the instruction is reserved.

```

31  26 25 24  20 19  15 14 12  11 10 9 8 7 6   0
  funct6 vm   vs2   vs1 funct3      rd      opcode
  111100 1    vs2   vs1    001    t t t 0 0  1110111 // sf.mm.f.f

  11111a 1    vs2   vs1    001    t t 0 0 b  1110111 // sf.mm.<a>.<b>

```

Encoding of a/b fields

0 e5m2

1 e4m3

```

  11110a 1    vs2   vs1    000    t t 0 0 b  1110111 // sf.mm.<a>.<b>

```

Encoding of a/b fields

0 unsigned

1 signed

1.9 Matrix Tile-Zero Instruction

The `sf.vtzero.t` instruction writes 0 to each element of the `tm` by `tn` submatrix of the destination tile. Any tile elements outside the range `[0, tm-1] x [0, tn-1]` are considered part of the tail and are handled using a tail-agnostic policy. If the matrix has not been configured, i.e. `vtwiden=0` in `vtype`, an illegal-instruction exception is raised. NOTE: SiFive implementations raise an illegal-instruction exception when executing `sf.vtzero.t` instructions with a nonzero `vstart`. This instruction is included in the `Xsfrmmbase` extension.

```

# Write zeros to tile
sf.vtzero.t mtd

```

```

31  26 25 24  20 19  15 14 12  11 10 9 8 7 6   0
  010000 1  11110 00000  110  t t t t 0  1010111 // sf.vtzero.t

```

1.10 Context Status

1.10.1 Xsfrm Context Status in mstatus

A `Xsfrm` context status field, `MS`, is added to `mstatus[30:29]` and shadowed in `sstatus[30:29]`. It is defined analogously to the floating-point context status field, `FS`.

Attempts to execute any instruction that accesses the tile state raises an illegal-instruction exception when `mstatus.MS` is set to Off. Note, the `vset*` instructions do not access the tile state.

When `mstatus.MS` is set to Initial or Clean, executing any instruction that changes tile state will change `mstatus.MS` to Dirty. Implementations may also change `mstatus.MS` from Initial or Clean to Dirty at any time, even when there is no change in tile state.

If `mstatus.MS` is Dirty, `mstatus.SD` is 1; otherwise, `mstatus.SD` is set in accordance with existing specifications.

1.10.2 Xsfmm Context Status in vsstatus

When the hypervisor extension is present, a Xsfmm context status field, MS, is added to vsstatus[30:29]. It is defined analogously to the floating-point context status field, FS.

When V=1, both vsstatus.MS and mstatus.MS are in effect: attempts to execute any instruction that accesses tile state raise an illegal-instruction exception when either field is set to Off.

When V=1 and neither vsstatus.MS nor mstatus.MS is set to Off, executing any instruction that changes tile state will change both mstatus.MS and vsstatus.MS to Dirty. Implementations may also change mstatus.MS or vsstatus.MS from Initial or Clean to Dirty at any time, even when there is no change in tile state.

If vsstatus.MS is Dirty, vsstatus.SD is 1; otherwise, vsstatus.SD is set in accordance with existing specifications.

If mstatus.MS is Dirty, mstatus.SD is 1; otherwise, mstatus.SD is set in accordance with existing specifications.

For the purposes of the mstatus.VS and vsstatus.VS fields, all Xsfmm instructions (including configuration instructions) are considered to modify vector state. For the purposes of the mstatus.FS and vsstatus.FS fields, the Xsfmm instructions that implicitly access fcsr are considered to modify floating-point state.

1.10.3 Xsfmm Context Discard Instruction

```
31 26 25 24 20 19 15 14 12 11 7 6 0
funct6 vm vs2 vs1 funct3 rd opcode
010000 1 11100 00000 110 00000 1010111 // sf.vtdiscard
```

The sf.vtdiscard instruction is provided to inform the runtime that the tile state is no longer useful and need not be saved. sf.vtdiscard raises an illegal-instruction exception if mstatus.MS=Off (or if V=1 and vsstatus.MS=Off). Otherwise, it changes mstatus.MS to Initial (and, if V=1, changes vsstatus.MS to Initial). This instruction is included in the Xsfmmbase extension.

`sf.vtdiscard` raises an illegal-instruction exception if `vtype.vill=1`, but does not raise an exception as a result of `vtype.vtwiden` being zero.

Note

SiFive implementations raise an illegal-instruction exception when executing `sf.vtdiscard` instructions with a nonzero `vstart`.

Note

`sf.vtdiscard` does not actually write the tile state. For ABI purposes, the tile state becomes UNPREDICTABLE as a result of executing this instruction.

1.11 C Intrinsics

1.11.1 Matrix Configuration Instructions

```
size_t __riscv_sf_vsettn_e8w1(size_t atn);
size_t __riscv_sf_vsettn_e8w2(size_t atn);
size_t __riscv_sf_vsettn_e8w4(size_t atn);
size_t __riscv_sf_vsettn_e16w1(size_t atn);
size_t __riscv_sf_vsettn_e16w2(size_t atn);
size_t __riscv_sf_vsettn_e16w4(size_t atn);
size_t __riscv_sf_vsettn_e32w1(size_t atn);
size_t __riscv_sf_vsettn_e32w2(size_t atn);
size_t __riscv_sf_vsettn_e64w1(size_t atn);
```

```
size_t __riscv_sf_vsettm_e8w1(size_t atn);
size_t __riscv_sf_vsettm_e8w2(size_t atn);
size_t __riscv_sf_vsettm_e8w4(size_t atn);
size_t __riscv_sf_vsettm_e16w1(size_t atn);
size_t __riscv_sf_vsettm_e16w2(size_t atn);
size_t __riscv_sf_vsettm_e16w4(size_t atn);
size_t __riscv_sf_vsettm_e32w1(size_t atn);
size_t __riscv_sf_vsettm_e32w2(size_t atn);
size_t __riscv_sf_vsettm_e64w1(size_t atn);
```

```
size_t __riscv_sf_vsettm_e8w1(size_t atm);
size_t __riscv_sf_vsettm_e8w2(size_t atm);
size_t __riscv_sf_vsettm_e8w4(size_t atm);
size_t __riscv_sf_vsettm_e16w1(size_t atm);
size_t __riscv_sf_vsettm_e16w2(size_t atm);
size_t __riscv_sf_vsettm_e16w4(size_t atm);
size_t __riscv_sf_vsettm_e32w1(size_t atm);
```

```
size_t __riscv_sf_vsettm_e32w2(size_t atm);
size_t __riscv_sf_vsettm_e64w1(size_t atm);

size_t __riscv_sf_vsettk_e8w1(size_t atk);
size_t __riscv_sf_vsettk_e8w2(size_t atk);
size_t __riscv_sf_vsettk_e8w4(size_t atk);
size_t __riscv_sf_vsettk_e16w1(size_t atk);
size_t __riscv_sf_vsettk_e16w2(size_t atk);
size_t __riscv_sf_vsettk_e16w4(size_t atk);
size_t __riscv_sf_vsettk_e32w1(size_t atk);
size_t __riscv_sf_vsettk_e32w2(size_t atk);
size_t __riscv_sf_vsettk_e64w1(size_t atk);
```

1.11.2 Load/Store Tile Subset to Memory

Non-overload

```
void __riscv_sf_vlte8_i8(size_t tss, const int8_t *base_addr, size_t vl);
void __riscv_sf_vlte8_u8(size_t tss, const uint8_t *base_addr, size_t vl);
void __riscv_sf_vlte16_i16(size_t tss, const int16_t *base_addr, size_t vl);
void __riscv_sf_vlte16_u16(size_t tss, const uint16_t *base_addr, size_t vl);
void __riscv_sf_vlte16_f16(size_t tss, const _Float16 *base_addr, size_t vl);
void __riscv_sf_vlte16_bf16(size_t tss, const __bf16 *base_addr, size_t vl);
void __riscv_sf_vlte32_i32(size_t tss, const int32_t *base_addr, size_t vl);
void __riscv_sf_vlte32_u32(size_t tss, const uint32_t *base_addr, size_t vl);
void __riscv_sf_vlte32_f32(size_t tss, const float *base_addr, size_t vl);
void __riscv_sf_vlte64_i64(size_t tss, const int64_t *base_addr, size_t vl);
void __riscv_sf_vlte64_u64(size_t tss, const uint64_t *base_addr, size_t vl);
void __riscv_sf_vlte64_f64(size_t tss, const double *base_addr, size_t vl);

void __riscv_sf_vste8_i8(size_t tss, int8_t *base_addr, size_t vl);
void __riscv_sf_vste8_u8(size_t tss, uint8_t *base_addr, size_t vl);
void __riscv_sf_vste16_i16(size_t tss, int16_t *base_addr, size_t vl);
void __riscv_sf_vste16_u16(size_t tss, uint16_t *base_addr, size_t vl);
void __riscv_sf_vste16_f16(size_t tss, _Float16 *base_addr, size_t vl);
void __riscv_sf_vste16_bf16(size_t tss, __bf16 *base_addr, size_t vl);
void __riscv_sf_vste32_i32(size_t tss, int32_t *base_addr, size_t vl);
void __riscv_sf_vste32_u32(size_t tss, uint32_t *base_addr, size_t vl);
void __riscv_sf_vste32_f32(size_t tss, float *base_addr, size_t vl);
void __riscv_sf_vste64_i64(size_t tss, int64_t *base_addr, size_t vl);
void __riscv_sf_vste64_u64(size_t tss, uint64_t *base_addr, size_t vl);
void __riscv_sf_vste64_f64(size_t tss, double *base_addr, size_t vl);
```

Overload

```
void __riscv_sf_vlte8(size_t tss, const int8_t *base_addr, size_t vl);
void __riscv_sf_vlte8(size_t tss, const uint8_t *base_addr, size_t vl);
void __riscv_sf_vlte16(size_t tss, const int16_t *base_addr, size_t vl);
void __riscv_sf_vlte16(size_t tss, const uint16_t *base_addr, size_t vl);
void __riscv_sf_vlte16(size_t tss, const _Float16 *base_addr, size_t vl);
void __riscv_sf_vlte16(size_t tss, const __bf16 *base_addr, size_t vl);
void __riscv_sf_vlte32(size_t tss, const int32_t *base_addr, size_t vl);
```

```
void __riscv_sf_vlte32(size_t tss, const uint32_t *base_addr, size_t vl);
void __riscv_sf_vlte32(size_t tss, const float *base_addr, size_t vl);
void __riscv_sf_vlte64(size_t tss, const int64_t *base_addr, size_t vl);
void __riscv_sf_vlte64(size_t tss, const uint64_t *base_addr, size_t vl);
void __riscv_sf_vlte64(size_t tss, const double *base_addr, size_t vl);

void __riscv_sf_vste8(size_t tss, int8_t *base_addr, size_t vl);
void __riscv_sf_vste8(size_t tss, uint8_t *base_addr, size_t vl);
void __riscv_sf_vste16(size_t tss, int16_t *base_addr, size_t vl);
void __riscv_sf_vste16(size_t tss, uint16_t *base_addr, size_t vl);
void __riscv_sf_vste16(size_t tss, _Float16 *base_addr, size_t vl);
void __riscv_sf_vste16(size_t tss, __bf16 *base_addr, size_t vl);
void __riscv_sf_vste32(size_t tss, int32_t *base_addr, size_t vl);
void __riscv_sf_vste32(size_t tss, uint32_t *base_addr, size_t vl);
void __riscv_sf_vste32(size_t tss, float *base_addr, size_t vl);
void __riscv_sf_vste64(size_t tss, int64_t *base_addr, size_t vl);
void __riscv_sf_vste64(size_t tss, uint64_t *base_addr, size_t vl);
void __riscv_sf_vste64(size_t tss, double *base_addr, size_t vl);
```

1.11.3 Move Tile Subset between Tile State and Vector Registers

Non-overload

```
vint8m8_t      __riscv_sf_vtmv_v_t_i8m8(size_t tss, size_t vl);
vuint8m8_t     __riscv_sf_vtmv_v_t_u8m8(size_t tss, size_t vl);
vfloat8e4m3m8_t __riscv_sf_vtmv_v_t_f8e4m3m8(size_t tss, size_t vl);
vfloat8e5m2m8_t __riscv_sf_vtmv_v_t_f8e5m2m8(size_t tss, size_t vl);
vint16m8_t     __riscv_sf_vtmv_v_t_i16m8(size_t tss, size_t vl);
vuint16m8_t    __riscv_sf_vtmv_v_t_u16m8(size_t tss, size_t vl);
vfloat16m8_t   __riscv_sf_vtmv_v_t_f16m8(size_t tss, size_t vl);
vbf16m8_t     __riscv_sf_vtmv_v_t_bf16m8(size_t tss, size_t vl);
vint32m8_t     __riscv_sf_vtmv_v_t_i32m8(size_t tss, size_t vl);
vuint32m8_t    __riscv_sf_vtmv_v_t_u32m8(size_t tss, size_t vl);
vfloat32m8_t   __riscv_sf_vtmv_v_t_f32m8(size_t tss, size_t vl);
vint64m8_t     __riscv_sf_vtmv_v_t_i64m8(size_t tss, size_t vl);
vuint64m8_t    __riscv_sf_vtmv_v_t_u64m8(size_t tss, size_t vl);
vfloat64m8_t   __riscv_sf_vtmv_v_t_f64m8(size_t tss, size_t vl);

void __riscv_sf_vtmv_t_v_i8m8(size_t tss, vint8m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v_u8m8(size_t tss, vuint8m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v_f8e4m3m8(size_t tss, vfloat8e4m3m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v_f8e5m2m8(size_t tss, vfloat8e5m2m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v_i16m8(size_t tss, vint16m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v_u16m8(size_t tss, vuint16m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v_f16m8(size_t tss, vfloat16m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v_bf16m8(size_t tss, vbf16m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v_i32m8(size_t tss, vint32m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v_u32m8(size_t tss, vuint32m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v_f32m8(size_t tss, vfloat32m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v_i64m8(size_t tss, vint64m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v_u64m8(size_t tss, vuint64m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v_f64m8(size_t tss, vfloat64m8_t src, size_t vl);
```

Overload

```
void __riscv_sf_vtmv_t_v(size_t tss, vint8m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v(size_t tss, vuint8m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v(size_t tss, vfloat8e4m3m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v(size_t tss, vfloat8e5m2m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v(size_t tss, vint16m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v(size_t tss, vuint16m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v(size_t tss, vfloat16m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v(size_t tss, vbf16m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v(size_t tss, vint32m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v(size_t tss, vuint32m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v(size_t tss, vfloat32m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v(size_t tss, vint64m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v(size_t tss, vuint64m8_t src, size_t vl);
void __riscv_sf_vtmv_t_v(size_t tss, vfloat64m8_t src, size_t vl);
```

1.11.4 Matrix Arithmetic Instructions

Non-overload

```
void __riscv_sf_mm_u_u_w4_u8m8_u8m8(size_t mtd, vuint8m8_t vs2, vuint8m8_t vs1, size_t tm,
size_t tn, size_t tk);
void __riscv_sf_mm_s_u_w4_i8m8_u8m8(size_t mtd, vint8m8_t vs2, vuint8m8_t vs1, size_t tm,
size_t tn, size_t tk);
void __riscv_sf_mm_u_s_w4_u8m8_i8m8(size_t mtd, vuint8m8_t vs2, vint8m8_t vs1, size_t tm,
size_t tn, size_t tk);
void __riscv_sf_mm_s_s_w4_i8m8_i8m8(size_t mtd, vint8m8_t vs2, vint8m8_t vs1, size_t tm, size_t
tn, size_t tk);

void __riscv_sf_mm_e5m2_e5m2_w4_u8m8_u8m8(size_t mtd, vuint8m8_t vs2, vuint8m8_t vs1, size_t
tm, size_t tn, size_t tk);
void __riscv_sf_mm_e5m2_e4m3_w4_u8m8_u8m8(size_t mtd, vuint8m8_t vs2, vuint8m8_t vs1, size_t
tm, size_t tn, size_t tk);
void __riscv_sf_mm_e4m3_e5m2_w4_u8m8_u8m8(size_t mtd, vuint8m8_t vs2, vuint8m8_t vs1, size_t
tm, size_t tn, size_t tk);
void __riscv_sf_mm_e4m3_e4m3_w4_u8m8_u8m8(size_t mtd, vuint8m8_t vs2, vuint8m8_t vs1, size_t
tm, size_t tn, size_t tk);

void __riscv_sf_mm_f_f_w2_f16m8(size_t mtd, vfloat16m8_t vs2, vfloat16m8_t vs1, size_t tm,
size_t tn, size_t tk);

void __riscv_sf_mm_f_f_w2_bf16m8(size_t mtd, vbf16m8_t vs2, vbf16m8_t vs1, size_t tm,
size_t tn, size_t tk);

void __riscv_sf_mm_f_f_w1_f32m8(size_t mtd, vfloat32m8_t vs2, vfloat32m8_t vs1, size_t tm,
size_t tn, size_t tk);

void __riscv_sf_mm_f_f_w1_f64m8(size_t mtd, vfloat64m8_t vs2, vfloat64m8_t vs1, size_t tm,
size_t tn, size_t tk);
```

Overload

```
void __riscv_sf_mm_w4(size_t mtd, vuint8m8_t vs2, vuint8m8_t vs1, size_t tm, size_t tn, size_t tk);
void __riscv_sf_mm_w4(size_t mtd, vint8m8_t vs2, vuint8m8_t vs1, size_t tm, size_t tn, size_t tk);
void __riscv_sf_mm_w4(size_t mtd, vuint8m8_t vs2, vint8m8_t vs1, size_t tm, size_t tn, size_t tk);
void __riscv_sf_mm_w4(size_t mtd, vint8m8_t vs2, vint8m8_t vs1, size_t tm, size_t tn, size_t tk);

void __riscv_sf_mm_w4(size_t mtd, vfloat8e5m2m8_t vs2, vfloat8e5m2m8_t vs1, size_t tm, size_t tn, size_t tk);
void __riscv_sf_mm_w4(size_t mtd, vfloat8e5m2m8_t vs2, vfloat8e4m3m8_t vs1, size_t tm, size_t tn, size_t tk);
void __riscv_sf_mm_w4(size_t mtd, vfloat8e4m3m8_t vs2, vfloat8e5m2m8_t vs1, size_t tm, size_t tn, size_t tk);
void __riscv_sf_mm_w4(size_t mtd, vfloat8e4m3m8_t vs2, vfloat8e4m3m8_t vs1, size_t tm, size_t tn, size_t tk);

void __riscv_sf_mm_w2(size_t mtd, vfloat16m8_t vs2, vfloat16m8_t vs1, size_t tm, size_t tn, size_t tk);

void __riscv_sf_mm_w2(size_t mtd, vbffloat16m8_t vs2, vbffloat16m8_t vs1, size_t tm, size_t tn, size_t tk);

void __riscv_sf_mm_w1(size_t mtd, vfloat32m8_t vs2, vfloat32m8_t vs1, size_t tm, size_t tn, size_t tk);

void __riscv_sf_mm_w1(size_t mtd, vfloat64m8_t vs2, vfloat64m8_t vs1, size_t tm, size_t tn, size_t tk);
```

1.11.5 Matrix Tile-Zero Instruction

```
void __riscv_sf_vtzero_t_e8w1(size_t mtd, size_t tm, size_t tn);
void __riscv_sf_vtzero_t_e8w2(size_t mtd, size_t tm, size_t tn);
void __riscv_sf_vtzero_t_e8w4(size_t mtd, size_t tm, size_t tn);
void __riscv_sf_vtzero_t_e16w1(size_t mtd, size_t tm, size_t tn);
void __riscv_sf_vtzero_t_e16w2(size_t mtd, size_t tm, size_t tn);
void __riscv_sf_vtzero_t_e16w4(size_t mtd, size_t tm, size_t tn);
void __riscv_sf_vtzero_t_e32w1(size_t mtd, size_t tm, size_t tn);
void __riscv_sf_vtzero_t_e32w2(size_t mtd, size_t tm, size_t tn);
void __riscv_sf_vtzero_t_e64w1(size_t mtd, size_t tm, size_t tn);
```

1.11.6 Context Status

```
void __riscv_sf_vtdiscard();
```

ABI Attributes

The RISC-V Matrix Extension introduces a set of function attributes that enforce compile-time safety for matrix tile state management and allow functions to communicate data in tile registers. These attributes create a type system for tile state, preventing common programming errors that could lead to undefined behavior, data corruption, or incorrect computation results when working with matrix operations.

ABI attributes have their own rules:

1. Outside of an attributed function, the matrix state is presumed to be inactive. A function marked `__riscv_new("xsfmm")` must be called from an inactive state, and creates an active state in its body, i.e. it is not allowed to have any attributed subroutine call that is also `__riscv_new("xsfmm")` which means it doesn't allow nested tile states to be created.
2. All attributes other than `__riscv_new("xsfmm")` must be called from inside an active state.
3. If a function is marked as `__riscv_in("xsfmm")`, it is only allowed to call functions with attribute `__riscv_in("xsfmm")` or `__riscv_preserves("xsfmm")`.
4. If a function is marked as `__riscv_preserves("xsfmm")`, it is only allowed to call functions with attribute `__riscv_preserves("xsfmm")`.

The compiler should implement semantic checks for function calls and emit an error at the front end when the ABI is violated. Builtin library calls (e.g. `libc`, `libm`, etc.) are not necessarily expanded to native instructions and are considered as normal calls, they also need to be checked. Other builtin (intrinsic) calls are considered as normal instructions and don't need to be checked.

This specification does not impose any rules on how data is arranged in tile state between function calls or which tile registers are used to pass data. Each caller-callee function pair must agree on their own convention.

Table 1: Function Calling Convention Matrix

Row calls Col?	None	new	out	in	inout	preserves
None	Y	Y				
new			Y	Y	Y	Y
out			Y	Y		Y
in				Y		Y
inout			Y	Y	Y	Y

Table 1: Function Calling Convention Matrix

Row calls Col?	None	new	out	in	inout	preserves
preserves						Y

2.1 __riscv_in("xsfmm")

The `__riscv_in("xsfmm")` keyword applies to prototyped function types and specifies that the function shares the tile with its caller. For `__riscv_in("xsfmm")`, the function takes the tile as input and returns with the tile unchanged.

Note

`__riscv_in("xsfmm")`, `__riscv_out("xsfmm")`, `__riscv_inout("xsfmm")`, `__riscv_new("xsfmm")` and `__riscv_preserves("xsfmm")` are mutually exclusive.

Example

```
void callee() __riscv_in("xsfmm");
void test() __riscv_in("xsfmm") {
    callee(); // read tile state
    callee(); // read tile state
}
```

2.2 __riscv_out("xsfmm")

The `__riscv_out("xsfmm")` keyword applies to prototyped function types and specifies that the function shares the tile with its caller. For `__riscv_out("xsfmm")`, the function ignores the incoming tile state and returns a new tile state.

Note

`__riscv_in("xsfmm")`, `__riscv_out("xsfmm")`, `__riscv_inout("xsfmm")`, `__riscv_new("xsfmm")` and `__riscv_preserves("xsfmm")` are mutually exclusive.

Example

```
void callee1() __riscv_out("xsfmm");
void callee2() __riscv_in("xsfmm");
void test() __riscv_new("xsfmm") {
```

```
callee1(); // new tile state generated
callee2(); // read the tile state
}
```

2.3 `__riscv_inout("xsfmm")`

The `__riscv_inout("xsfmm")` keyword applies to prototyped function types and specifies that the function shares the tile with its caller. For `__riscv_inout("xsfmm")`, the function takes the tile as input and returns a new tile state.

Note

`__riscv_in("xsfmm")`, `__riscv_out("xsfmm")`, `__riscv_inout("xsfmm")`, `__riscv_new("xsfmm")` and `__riscv_preserves("xsfmm")` are mutually exclusive.

Example

```
void callee1() __riscv_inout("xsfmm");
void callee2() __riscv_in("xsfmm");
void test() __riscv_new("xsfmm") {
    callee1(); // read the tile state and produce the new one.
    callee2(); // read the tile state.
}
```

2.4 `__riscv_new("xsfmm")`

The `__riscv_new("xsfmm")` keyword applies to prototyped function types and specifies that the function will create a new scope for tile state.

Example

```
void callee() __riscv_inout("xsfmm");
void test() __riscv_new("xsfmm") {
    callee();
    callee();
}
```

Note

`__riscv_in("xsfmm")`, `__riscv_out("xsfmm")`, `__riscv_inout("xsfmm")`, `__riscv_new("xsfmm")` and `__riscv_preserves("xsfmm")` are mutually exclusive.

2.5 `__riscv_preserves("xsfmm")`

The `__riscv_preserves("xsfmm")` keyword applies to prototyped function types and specifies that the function does not read or write the tile state and returns with the tile state unchanged. If a function is marked as `__riscv_preserves("xsfmm")`, it is not allowed to have any attributed subroutine call that might read or write the tile state.

The `[[__riscv_preserves("xsfmm")]]` keyword applies to the statement which is used for the call where the callee is not defined as an attributed function, e.g. libc calls such as `printf`. Ref: <https://clang.llvm.org/docs/AttributeReference.html#noinline>. Note that the use of `[[__riscv_preserves]]` with functions like `printf` is not technically correct and should not be used in production code, but is intended as a temporary hack to support debugging.

Note

`__riscv_in("xsfmm")`, `__riscv_out("xsfmm")`, `__riscv_inout("xsfmm")`, `__riscv_new("xsfmm")` and `__riscv_preserves("xsfmm")` are mutually exclusive.

Example

```
void callee() __riscv_inout("xsfmm");
void test() __riscv_preserves("xsfmm") {
    [[__riscv_preserves]] printf("foo");
}
```

2.6 Example Scenario

```
void function();
void setup_tile() __riscv_out("xsfmm");
void use_tile() __riscv_in("xsfmm");
void shares_tile() __riscv_inout("xsfmm");

void test() __riscv_new("xsfmm") {
    setup_tile();    // tile is live after this call

    shares_tile();  // tile is live before and after this call, no save necessary
    [[__riscv_preserves]] function();    // This statement attribute tells compiler that tile
    state is not modified across the call.

    use_tile();     // tile is no longer live after this call
}
```

2.7 Analogy Using General Program

```
int generate_int();
```

Xsfrm Family of Attached Matrix Extensions

```
void read_int(const int&);
int read_write_int(const int&);
void test() {
    int a = 0;           // zero the tile
    a = generate_int();   // write only: __tile_out
    read_int(a);          // read only: __tile_in
    a = read_write_int(a); // read and write: __tile_inout
    read_int(a);          // read only: __tile_in
    return;
}
```