900, 900/L, 900/H CPU Core Different Points

There are 3 type CPU core: 1. 900, 2. 900/L, 3. 900/H in TLCS-900 series and they are different from following points.

In this chapter, only the 2. 900/L CPU are explained.

CPU / Different Point	1. 900	2. 900/L	3. 900/H
CPU Operating mode	System/Normal mode	System mode	System mode
CPU Register mode	After reset, the mode is set to MIN (minimum) mode. The MAX mode is set by a 'MAX' instruction.	After reset, the mode is set to MAX (maximum) mode. (Note1)	After reset, the mode is set to MAX (maximum) mode.
Interrupt vector formula	Restart formula	Vector formula	Vector formula
Normal Stack Pointer XNSP	exist	not exist	not exist
Interrupt Nesting Counter INTNEST	not exist	exist	exist
Operating Voltage	5 V ± 10%	2.7 to 5.5 V (Note2)	5 V ± 10%

CPU Different Points

Note1: TMP93CS32, TMP93PW32, TMP93CS44/S45, TMP93CU44, TMP93CW44, TMP93PS44, TMP93PW44A, TMP93CS20, TMP93PW20A, TMP93CT/PT75, TMP93CW/PW76 and TMP93C071 can't set the MIN mode.

Note2: Operating voltage of TMP93CS42A, TMP93PS42A and TMP93PW44A is 5 V ± 10%.

1. Outline

The TLCS-900 series is the Toshiba proprietary high-performance 16-bit CPU. With the CPU and various I/O function blocks (such as timers, serial I/Os, ADs) on broad it fits in various applications.
Though the TLCS-900 CPU is 16-bit CPU, it has 32-bit register bank configuration, and it is suitable as an embedded controller.
The TLCS-900 CPU features are as follows:

(1) Upward compatible with TLCS-90
- Upward compatibility on mnemonic and register set levels

(2) General-purpose registers
- All 8 registers can be used as an accumulator

(3) Register bank system
One of the follwing mode can be selected.
- Minimum mode: eight 16-bit register banks
- Maximum mode: four 32-bit register banks

(4) 16 Mbyte linear address space; 9 types of addressing modes

(5) Dynamic bus sizing system
- Can consist 8-/16-bit external data bus together

(6) High reliability
- Supporting system mode and normal mode (900)
- Supporting only system mode (900/L, 900/H)

(7) Orthogonal instruction sets
- 8-/16-/32-bit data transfer/arithmetic instructions
- 16-bit multiplication/division
  16 × 16 to 32-bits (signed/unsigned)
  32 ÷ 16 to 16-bits (unsigned/signed)
- Bit processing including bit arithmetic
- Supporting instruction for C compiler
- Filter calculations: multiplication-addition arithmetic, module increment instruction

(8) High-speed processing
- Pipeline system with 4-byte instruction queue buffer
- 16-bit ALU

2. CPU Operating Modes

The 900/L has one type of operating: system mode. In system mode, there are no restrictions on using instructions or registers.

The CPU resources effective in system mode are as follows:

(1) General-purpose registers
- Four 16-bit general-purpose registers × 8 banks (minimum mode)
  or
  Four 32-bit general-purpose registers × 4 banks (maximum mode)
- Four 32-bit general-purpose registers (including system stack pointer: XSP)

(2) Status register (SR): including system mode flag.

(3) Program counter (PC): 32 bits for maximum mode. 16 bits for minimum mode.

(4) Control register: parameter register for micro DMA, etc.

(5) All CPU instructions

(6) All built-in I/O registers

(7) All built-in memories

Note: TMP93CS32, TMP93PW32, TMP93CS44/S45, TMP93CU44, TMP93CW44, TMP93PS44, TMP93PW44A, TMP93CS20, TMP93PW20A, TMP93CT/PT75, TMP93CW/PW76 and TMP93C071 can't set the MIN mode.

3. Registers

3.1 Register Structure

Figure 3.1.1, Figure 3.1.2 and Figure 3.1.3 illustrates the format of registers. The TLCS-900/L has two register modes.

(1) Minimum mode: 64-Kbyte program area/16-Mbyte data area

Four 16-bit general-purpose registers × 8 banks

+

Four 32-bit general-purpose registers

+

16-bit program counter

(2) Maximum mode: 16-Mbyte program area/16-Mbyte data area

Four 32-bit general-purpose registers × 4 banks

+

Four 32-bit general-purpose registers

+

32-bit program counter

Register mode changing

The <MAX> bit in status register (SR) is initialized to "1" and set to Maximum mode by resetting. The "MIN" instruction changes to Minimum mode.

Note: TMP93CS32, TMP93PW32, TMP93CS44/S45, TMP93CU44, TMP93CW44, TMP93PS44, TMP93PW44A, TMP93CS20, TMP93PW20A, TMP93CT/PT75, TMP93CW/PW76 and TMP93C071 can't set the MIN mode.

Stack Pointer

The stack pointer is provided for only System mode (XSP). The System stack pointer (XSP) is set to 100H by resetting.

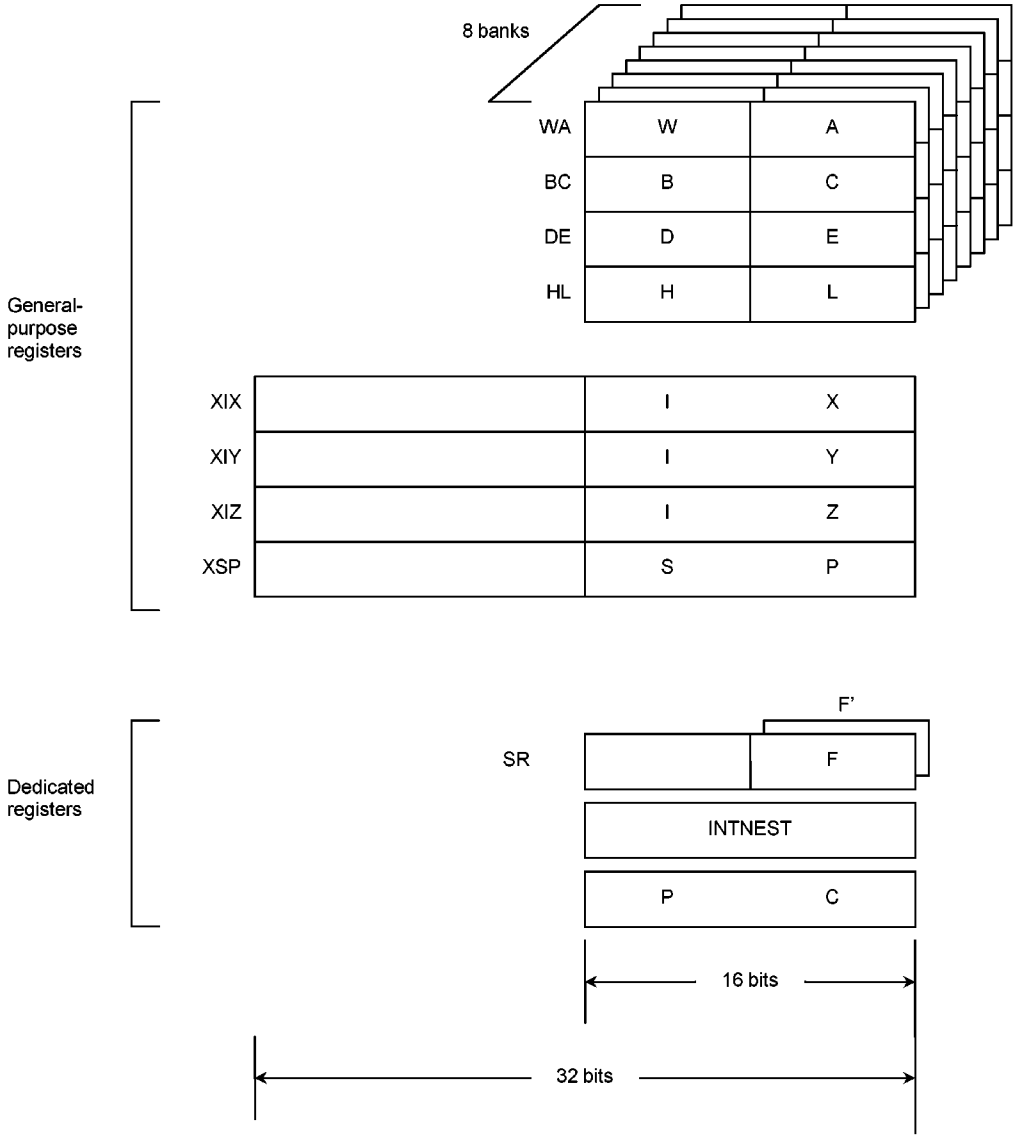


Figure 3.1.1 Register Format (minimum mode: 64-Kbyte program area)

Note: The data memory area is 16 Mbyte.
The whole 16-Mbyte area can be accessed by using the registers (XIX, XIY, XIZ, XSP) or absolute addressing mode.

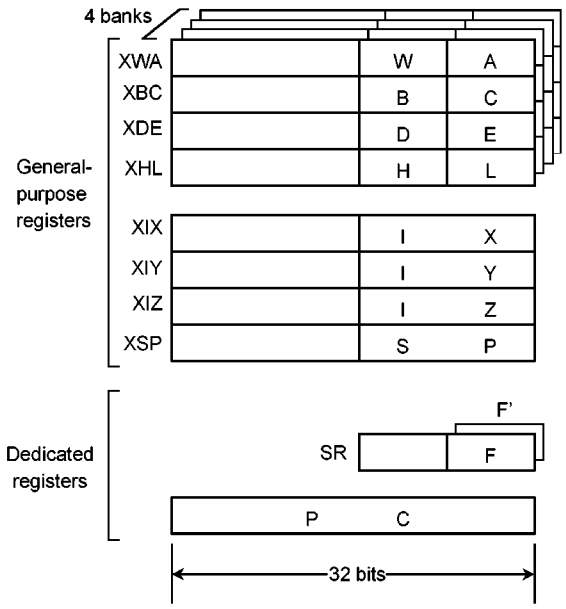


Figure 3.1.2 Register Format (maximum mode: 16-Mbyte program area)

To change from maximum to minimum mode or from minimum to maximum mode or from minimum to maximum mode, there is no dedicated instruction; instead, the RETI or POP SR instruction changes the <MAX> bit of the status register.

When the mode changes from minimum to maximum, the 16-bit general-purpose registers (WA, BC, DE and HL) are extended to 32-bit general-purpose registers (XWA, XBC, XDE and XHL). The value of the upper 16 bits (that is, bit 16 to bit 31) are undefined. Those registers need to be initialized before use. Changing the mode from minimum to maximum also extends the program counter to 32 bits which automatically writes "0" to the upper 16 bits. It keeps the program continuity.

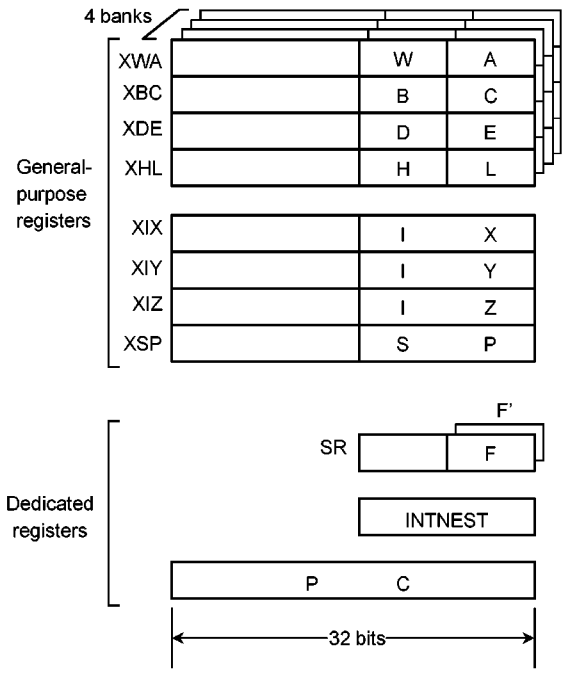


Figure 3.1.3 Register Format (maximum mode: 16-Mbyte program area)

To change from maximum to minimum mode or from minimum to maximum mode or from minimum to maximum mode, there is no dedicated instruction; instead, the RETI or POP SR instruction changes the <MAX> bit of the status register.

When the mode changes from minimum to maximum, the 16-bit general-purpose registers (WA, BC, DE and HL) are extended to 32-bit general-purpose registers (XWA, XBC, XDE and XHL). The value of the upper 16 bits (that is, bit 16 to bit 31) are undefined. Those registers need to be initialized before use. Changing the mode from minimum to maximum also extends the program counter to 32 bits which automatically writes "0" to the upper 16 bits. It keeps the program continuity.

3.2 Register Details

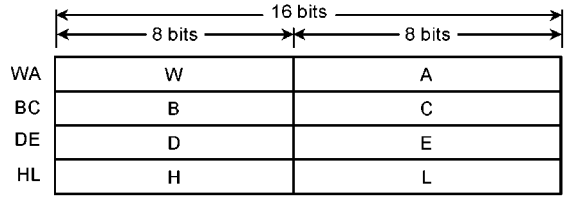
3.2.1 General-purpose bank registers

As explained in the previous section, the TLCS-900 series has two register formats. Which of the register formats is used depends on whether the mode is minimum or maximum. In either way, the register sets and registers in each bank are used exactly the same.

(1) General-purpose Bank Registers in Minimum Mode

In minimum mode , the following four 16-bit general-purpose registers consisting of 8 banks can be used. The register format in a bank is shown below.

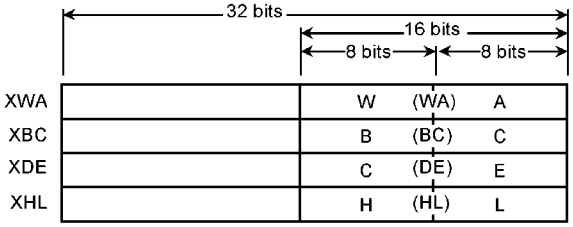
Four 16-bit registers (WA, BC, DE and HL) are general-purpose registers and can be used as accumulators, index registers, and displacement registers. They can also be used as 8-bit registers (W, A, B, C, D, E, H and L) to function for example as accumulators.



(2) General-purpose Bank Registers in Maximum Mode

In maximum mode, the following four 32-bit general-purpose registers consisting of 4 banks can be used. The register format in a bank is shown below.

Four 32-bit registers (XWA, XBC, XDE and XHL) are general-purpose registers and can be used as an accumulators and index registers. They can also be used as 16-bit registers (WA, BC, DE and HL), in which case, the lower 16 bits of the 32-bit registers are assigned.



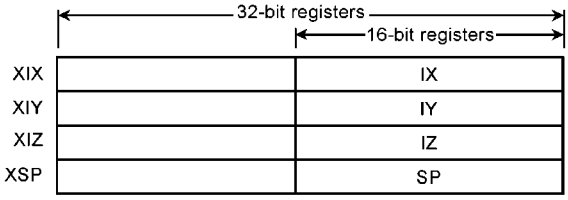
Note: Round brackets ( ) signify 16-bit registers.

16-bit registers can be used as accumulators, index registers in index addressing mode, and displacement registers. They can also be used as two 8-bit general-purpose registers (W, A, B, C, D, E, H and L) to function for example as accumulators.

3.2.2 32-bit General-purpose Registers

The TLCS-900 series has four 32-bit general-purpose registers (XIX, XIY, XIZ and XSP). They are fixed, independent of maximum or minimum mode. The register format is shown below.

These registers can also be used as accumulators, index registers, and displacement registers. They can be used either as 16-bit, or 8-bit registers. Names when registers are used as 8-bit registers are listed later.



Stack Pointer

The XSP register is utilized for stack pointer. It is used when the interrupt is occured or "CALL", "RET" instruction are executed. The stack pointer (XSP) is set to 100H by resetting.

3.2.3 Status Register (SR)

The status register contains flags indicating the status (operating mode, register format, etc.) of the CPU and operation results. This register consists of two parts. The upper byte of the status register (bits 8 to 15) indicates the CPU status. The lower byte (bits 0 to 7) are referred to as the flag register (F). This indicates the status of the operation result. The TLCS-900 series has two flag registers (F and F'). They can be switched using the EX instruction.

(1) Upper Byte of Status Register

15	14	13	12	11	10	9	8
SYSM	IFF2	IFF1	IFF0	MAX	RFP2	RFP1	RFP0

1. SYSM (System Mode)

Indicates the CPU operating mode, system or normal. 900/L has only system mode.

It is initialized to 1 (system mode) when reset.

0	Normal mode
1	System mode (900/L has only this mode.)

2. IFF2 to IFF0 (Interrupt mask Flip-Flop2 to 0)

Mask registers with interrupt levels from 1 to 7. Level 7 has the highest priority.

They are initialized to 111 when reset.

000	Enables interrupts with level 1 or higher.	Same
001	Enables interrupts with level 1 or higher.	
010	Enables interrupts with level 2 or higher.	
011	Enables interrupts with level 3 or higher.	
100	Enables interrupts with level 4 or higher.	
101	Enables interrupts with level 5 or higher.	
110	Enables interrupts with level 6 or higher.	
111	Enables interrupts with level 7 only (non-maskable interrupt).	

Any value can be set using the EI instruction.

When an interrupt is received, the mask register sets a value higher by 1 than the interrupt level received. When an interrupt with level 7 is received, 111 is set. Unlike with the TLCS-90 series, the EI instruction becomes effective immediately after execution.

3. MAX ($\overline{\text{MIN}}$imum/ MAXimum)

Bit used to specify the register mode which determines the sizes of the register banks and the program counter.

0	Minimum mode
1	Maximum mode (900/L1 has only this mode.)

If the program size exceeds 64 K bytes, use the "MAX" instruction to set this register to "1" so that register mode becomes maximum mode.

Initialized to "1" (maximum mode) for 900/L by reset.

Note: TMP93CS32, TMP93PW32, TMP93CS44/S45, TMP93CU44, TMP93CW44, TMP93PS44, TMP93PW44A, TMP93CS20, TMP93PW20A, TMP93CT/PT75, TMP93CW/PW76 and TMP93C071 can't set the MIN mode.

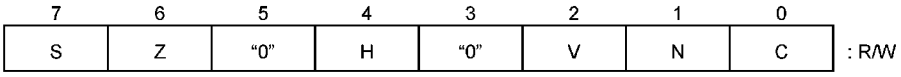
4. RFP2 to RFP0 (Register File Pointer2 to 0)

   Indicates the number of register file (register bank) currently being used. Initialized to 000 by reset.

   The values in these registers can be operated on using the following three instructions. RFP2 is fixed to 0 in maximum mode. It remains 0 even if an attempt to change it to 1 using following instructions.

   - LDF imm ; RFP ← imm (0 to 7)
   - INCF ; RFP ← RFP + 1
   - DECF ; RFP ← RFP – 1

(2) Flag Register, F



1. S (Sign flag)

   "1" is set when the operation result is negative, "0" when positive.

   (The value of the most significant bit of the operation result is copied.)

2. Z (Zero flag)

   "1" is set when the operation result is zero, otherwise "0".

3. H (Half carry flag)

   "1" is set when a carry or borrow from bit 3 to bit 4 occurs as a result of the operation, otherwise "0". With a 32-bit operation instruction, an undefined value is set.

4. V (Parity/over-flow flag)

   Indicates either parity or overflow, depending on the operation type.

   Parity (P): "0" is set when the number of bits set to 1 is odd, "1" when even. An undefined value is set with a 32-bit operation instruction.

   Overflow (V): "0" is set if no overflow, if overflow "1".

5. N (Negative)

   ADD/SUB flag

   "0" is set after an addition instruction such as ADD is executed, "1" after a subtraction instruction such as SUB.

   Used when executing the DAA (decimal addition adjust accumulator) instruction.

6. C (Carry flag)

   "1" is set when a carry or borrow occurs, otherwise "0".

Read and write process of status register

Read from bits 0 to 15	1. PUSH SR POP dst
Write to bits 0 to 15	1. POP SR
Only bit 15 <SYSM>	"1" is always set, because 900/L1 CPU has only system mode.
Only bits 14 to 12 <IFF2:0>	1. EI num A value of "num" is written.
Only bit 11 <MAX>	1. MIN "0" is written.
Only bits 10 to 8 <RFP2:0>	1. LDF imm 2. INCF 3. DECF
Only bits 7 to 0	1. PUSH F/POP F 2. EX F, F' 3. A flag is set indirectly by executing arithmetic instructions etc.

3.2.4 Program Counter (PC)

The program counter is a pointer indicating the memory address to be executed next.

The program counter bit length depends on whether the register format is in minimum or maximum mode.

In minimum mode, the program counter consists of 16 bits, and a maximum program area of 64 Kbytes (from addresses 000000H to 00FFFFH) can be accessed.

In maximum mode, the program counter consists of 32 bits. The size of the program area depends on the number of the address pins that the product has. With 24 address pins (A0 to A23), a maximum program area of 16 M bytes can be accessed as a linear address space. In this case, the upper 8 bits of the program counter (bits 24 to 31) are ignored.

When the register format changes from minimum to maximum mode, the upper word of the program counter (bits 16 to 31) is extended so that the program counter becomes 32 bits long. This automaticaly writes "0" to the upper word of the program counter. So doing ensures program continuity.

PC after reset

The 900/L reads a value of a reset vector from a vector base address by reset and sets the value into a program counter. Then, program after the vector specified by the program counter are executed.

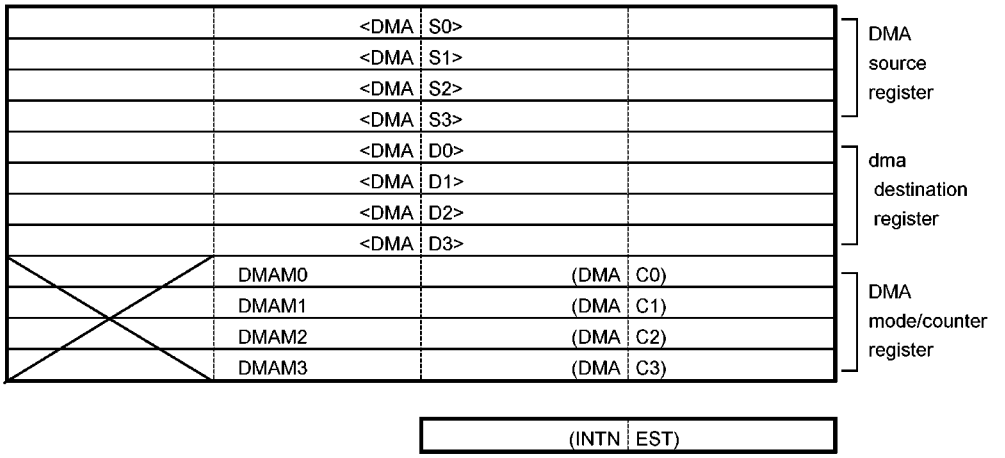
The vector base address is depending on products. They are given below.

Type No.	Vector Base Address	PC setting sequence after reset	Notes
93CM40 93CS40/CS41 93PS40 93CW40/CW41 93PW40 93CS42A/PS42A 93CW46A/PW46A	008000H	PC (7:0) ← address 8000H PC (15:8) ← address 8001H PC (23:16) ← address 8002H	P27 to 20/A23 to 16 pins are input ports with pull-down due to reset. The logic data is "00H". When Port 2 is used as A23 to 16 pins to access the program ROM, set PC (23 to 16) to "00H" and the reset vector to "0000H to FFFFH". (for mainly products without ROM)
93CS44/CS45 93CU44 93CW44 93PS44 93PW44A 93CS32/PW32 93CS20 93PW20A	0FFFF00H	PC (7:0) ← address FFFF00H PC (15:8) ← address FFFF01H PC (23:16) ← address FFFF02H	P27 to 20/A23 to 16 pins are input ports with pull-up due to reset. The logic data is "FFH". When Port 2 is used as A23 to 16 pins to access the program ROM, set PC (23 to 16) to "FFH" and the reset vector to "FF0000H to FFFFFFH". (for mainly products without ROM)

3.2.5 Control registers (CR)

The control registers consist of registers used to control micro DMA operation and an interrupt nesting counter. Control registers can be accessed by using the LDC instruction.

Control registers are illustrated below.



( ) : Word register name (16 bits)
< > : Long word register name (32 bits)

(1) For high-speed micro DMA, refer to "Part 2 TLCS-900 LSI Devices".

(2) Interrupt nesting counter: INTNEST

The INTNEST is incremented by acceptance of interrupts and decremented by execution of the RETI instruction.

However the INTNEST is not incremented by acceptance of software interrupts and illegal (undefined) instruction execution interrupts which is generated by the CPU.

The INTNEST is undefined by reset, so set the initial-value of INTNEST prior to using. The following shows an example of setting.

```
Example: LD   WA, 0
         LDC  INTNEST, WA
```

3.3 Register Bank Switching

Register banks are classified into the following three types.

Current bank registers

Previous bank registers

Absolute bank registers

The current bank is indicated by the register file pointer, <RFP>, (status register bits 8 to 10). The registers in the current bank are used as general-purpose registers, as described in the previous section. By changing the contents of the <RFP>, another register bank becomes the current register bank.

The previous bank is indicated by the value obtained by subtracting 1 from the <RFP>. For example, if the current bank is bank 3, bank 2 is the previous bank. The names of registers in the previous bank are indicated with a dash (WA', BC', DE', HL'). The EX instruction (EX A,A') is used to switch between current and previous banks.

All bank registers, including the current and previous ones, have a numerical value (absolute bank number) to indicate the bank. With a register name which includes a numerical value such as RW0, RA0, etc., all bank registers can be used. These registers (that is, all registers) are called absolute bank registers.

The TLCS-900/L CPU is designed to perform optimally when the current bank registers are operated as the working registers. In other words, if the CPU uses other bank registers, its performance degrades somewhat. In order to obtain maximum CPU efficiency, the TLCS-900/L has a function which easily switches register banks.

The bank switching function provides the following advantages:

- Optimum CPU operating efficiency
- Reduced programming size (Object codes)
- Higher response speed and reduced programming size when used as a context switch for an interrupt service routine.

Bank switching is performed by the instructions listed below.

LDF imm : Sets the contents of the immediate value in <RFP>. imm: 0 to 7

INCF : Increments <RFP> by 1.

DECF : Decrements <RFP> by 1.

In minimum mode, the immediate values used by the LDF instruction are from 0 to 7, in maximum mode 0 to 3. If a carry or borrow occurs when the INCF or DECF instruction is executed, it is ignored. The value of the <RFP> rotates. For example, if the INCF instruction is executed with bank 7, the result is bank 0. If the DECF instruction is executed with bank 0, the result is bank 7. Note that careless execution of the INCF or DECF instruction may destroy the contents of the register bank.

- Example of Register Bank Usage

The TLCS-900/L registers are formatted in banks. Banks can be used for processing objectives or interrupt levels. Two examples are given below.

<Example 1> When assigning register banks to interrupt processing routines.

Register bank 0 = Used for the main program and interrupt processing other than that shown below.

Register bank 1 = Used for processing INT0 .

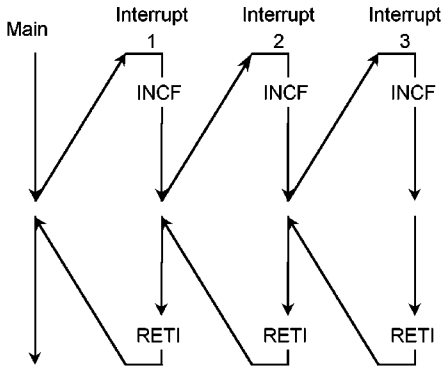
Register bank 2 = Used for processing timer 0.

Register bank 3 = Used for processing timer 1.

For example, if a timer 1 interrupt occurs during main program execution, processing jumps to a subroutine as follows. PUSH/POP processing for the register is unnecessary.

```
LDF 3     ; Sets register bank to 3.
  :
  :
RETI      ; Returns to previous status including <RFP>.
```

<Example 2> When assigning register banks to their appropriate interrupt level nesting.



Note 1: In the above example, when interrupt nesting exceeds the number of register banks (4), the <RFP> becomes 000 and the contents of register bank 0 are destroyed.

Note 2: The INCF instruction is used to execute <RFP> ← <RFP> + 1.

3.4 Accessing General-purpose Registers

The register access code is formatted in a varied code length on byte basis. The current bank registers can be accessed by the shortest code length. All general-purpose registers can be accessed by an instruction code which is 1 byte longer. General-purpose registers are as follows.

1. General-purpose registers in current bank

(Minimum mode)

					W	(W	A)	A
					B	(B	C)	C
					D	(D	E)	E
					H	(H	L)	L

(Maximum mode)

QW	(Q	WA)	QA	<X	WA >	W	(W	A)	A
QB	(Q	BC)	QC	<X	BC >	B	(B	C)	C
QD	(Q	DE)	QE	<X	DE >	D	(D	E)	E
QH	(Q	HL)	QL	<X	HL >	H	(H	L)	L

( ) : Word register name (16 bits)

< > : Long word register name (32 bits)

2. General-purpose registers in previous bank

(Minimum mode)

					W'	(W	A')	A'
					B'	(B	C')	C'
					D'	(D	E')	E'
					H'	(H	L')	L'

(Maximum mode)

QW'	(Q	WA')	QA'	<X	WA' >	W'	(W	A')	A'
QB'	(Q	BC')	QC'	<X	BC' >	B'	(B	C')	C'
QD'	(Q	DE')	QE'	<X	DE' >	D'	(D	E')	E'
QH'	(Q	HL')	QL'	<X	HL' >	H'	(H	L')	L'

3. 32-bit general-purpose registers

(Both minimum and maximum modes)

QIXH	(Q	IX)	QIXL	<X	IX >	IZH	(I	X)	IXL
QIYH	(Q	IY)	QIYL	<X	IY >	IYH	(I	Y)	IYL
QIZH	(Q	IZ)	QIZL	<X	IZ >	IZH	(I	Z)	IZL
QSPH	(Q	SP)	QSPL	<X	SP >	SPH	(S	P)	SPL

4. Absolute bank registers

(Minimum mode)

RW0	(RWA	0)	RA0			Bank0
RB0	(RBC	0)	RC0			
RD0	(RDE	0)	RE0			
RH0	(RHL	0)	RL0			
RW1	(RWA	1)	RA1			Bank1
RB1	(RBC	1)	RC1			
RD1	(RDE	1)	RE1			
RH1	(RHL	1)	RL1			
RW2	(RWA	2)	RA2			Bank2
RB2	(RBC	2)	RC2			
RD2	(RDE	2)	RE2			
RH2	(RHL	2)	RL2			
RW3	(RWA	3)	RA3			Bank3
RB3	(RBC	3)	RC3			
RD3	(RDE	3)	RE3			
RH3	(RHL	3)	RL3			
RW4	(RWA	4)	RA4			Bank4
RB4	(RBC	4)	RC4			
RD4	(RDE	4)	RE4			
RH4	(RHL	4)	RL4			
RW5	(RWA	5)	RA5			Bank5
RB5	(RBC	5)	RC5			
RD5	(RDE	5)	RE5			
RH5	(RHL	5)	RL5			
RW6	(RWA	6)	RA6			Bank6
RB6	(RBC	6)	RC6			
RD6	(RDE	6)	RE6			
RH6	(RHL	6)	RL6			
RW7	(RWA	7)	RA7			Bank7
RB7	(RBC	7)	RC7			
RD7	(RDE	7)	RE7			
RH7	(RHL	7)	RL7			

(Maximum mode)

QW0	(QWA	0)	QA0	<XWA	0>	RW0	(RWA	0)	RA0	Bank0
QB0	(QBC	0)	QC0	<XBC	0>	RB0	(RBC	0)	RC0	
QD0	(QDE	0)	QE0	<XDE	0>	RD0	(RDE	0)	RE0	
QH0	(QHL	0)	QL0	<XHL	0>	RH0	(RHL	0)	RL0	
QW1	(QWA	1)	QA1	<XWA	1>	RW1	(RWA	1)	RA1	Bank1
QB1	(QBC	1)	QC1	<XBC	1>	RB1	(RBC	1)	RC1	
QD1	(QDE	1)	QE1	<XDE	1>	RD1	(RDE	1)	RE1	
QH1	(QHL	1)	QL1	<XHL	1>	RH1	(RHL	1)	RL1	
QW2	(QWA	2)	QA2	<XWA	2>	RW2	(RWA	2)	RA2	Bank2
QB2	(QBC	2)	QC2	<XBC	2>	RB2	(RBC	2)	RC2	
QD2	(QDE	2)	QE2	<XDE	2>	RD2	(RDE	2)	RE2	
QH2	(QHL	2)	QL2	<XHL	2>	RH2	(RHL	2)	RL2	
QW3	(QWA	3)	QA3	<XWA	3>	RW3	(RWA	3)	RA3	Bank3
QB3	(QBC	3)	QC3	<XBC	3>	RB3	(RBC	3)	RC3	
QD3	(QDE	3)	QE3	<XDE	3>	RD3	(RDE	3)	RE3	
QH3	(QHL	3)	QL3	<XHL	3>	RH3	(RHL	3)	RL3	

( ) : Word register name (16 bits)

< > : Long word register name (32 bits)

4. Addressing Modes

The TLCS-900/L has nine addressing modes. These are combined with most instructions to improve CPU processing capabilities.

TLCS-900 series addressing modes are listed below. They cover the entire TLCS-90 addressing modes.

No.	Addressing mode	Description
1.	Register	reg8 reg16 reg32
2.	Immediate	n8 n16 n32
3.	Register indirect	(reg)
4.	Register indirect pre-decrement	(–reg)
5.	Register indirect post-increment	(reg+)
6.	Index	(reg + d8) (reg + d16)
7.	Register index	(reg + reg8) (reg + reg16)
8.	Absolute	(n8) (n16) (n24)
9.	Relative	(PC + d8) (PC + d16)

reg 8: All 8-bit registers such as W, A, B, C, D, E, H, L, etc.

reg 16: All 16-bit registers such as WA, BC, DE, HL, IX, IY, IZ, SP, etc.

reg 32: All 32-bit registers such as XWA, WBC, XDE, XHL, XIX, XIY, XIZ, XSP, etc.

reg: All 32-bit registers such as XWA, WBC, XDE, XHL, XIX, XIY, XIZ, XSP, etc. (Maximum mode)
All 16-bit bank registers such as WA, BC, DE, HL, etc. and XIX, XIY, XIZ and XSP. (Minimum mode)

d8: 8-bit displacement (–80H to + 7FH)

d16: 16-bit displacement (–8000H to + 7FFFH)

n8: 8-bit constant (00H to FFH)

n16: 16-bit constant (0000H to FFFFH)

n32: 32-bit constant (00000000H to FFFFFFFFH)

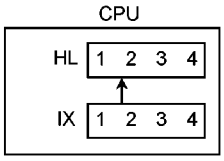
Note 1: Relative addressing mode can only be used with the following instructions:
LDAR, JR, JRL, DJNZ, and CALR

Note 2: In minimum mode, register bank blocks (current bank registers and previous bank registers, and bank 0 to 7 registers) consist of 16 bits. When these 16-bit registers are used for addressing, the CPU extends bits 16 to 31 to 0000H for address calculations.

(1) Register Addressing Mode

In this mode, the operand is the specified register.

Example: LD HL, IX

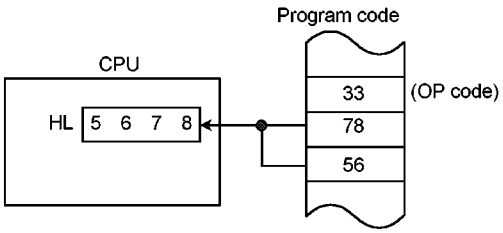


The IX register contents, 1234H, are loaded to the HL register.

(2) Immediate Addressing Mode

In this mode, the operand is in the instruction code.

Example: LD HL, 5678H

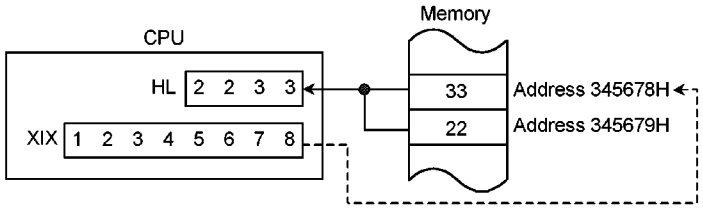


The immediate data, 5678H, is loaded to the HL register.

(3) Register Indirect Addressing Mode

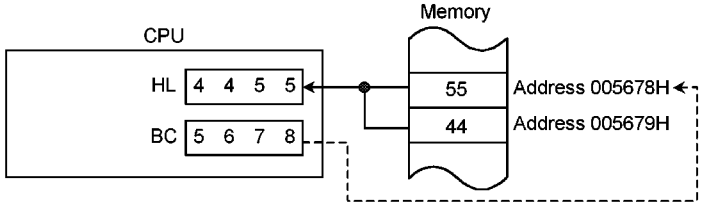
In this mode, the operand is the memory address specified by the contents of the register.

Example 1: LD, HL, (XIX)…in both minimum and maximum modes



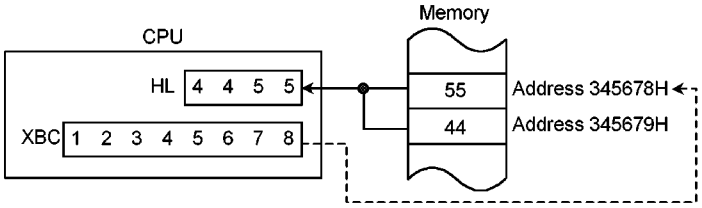
Memory data, 2233H, at address 345678H is loaded to the HL register.

Example 2: LD HL, (BC)…in minimum mode



In minimum mode, if a bank register (WA, BC, DE or HL) is used for addressing, address bits 16 to 23 are set to 00H.

Example 3: LD, HL, (XBC)…in maximum mode

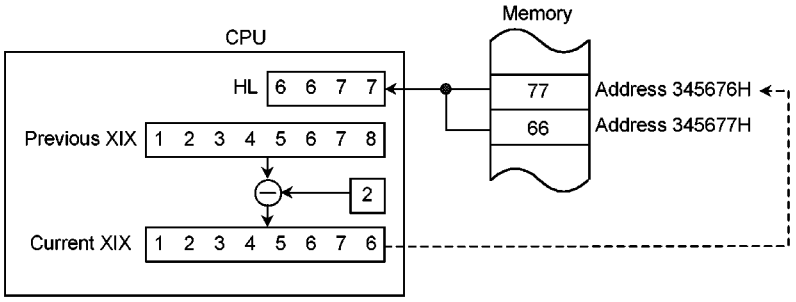


In maximum mode, if a bank register (XWA, XBC, XDE or XHL) is used for addressing, the values of bits 0 to 23 are output to the address bus.

(4) Register Indirect Pre-decrement Addressing Mode

In this mode, the contents of the register is decremented by the pre-decrement values. In this case, the operand is the memory address specified by the decremented register.

Example 1: LD HL, (–XIX)...in both minimum and maximum modes



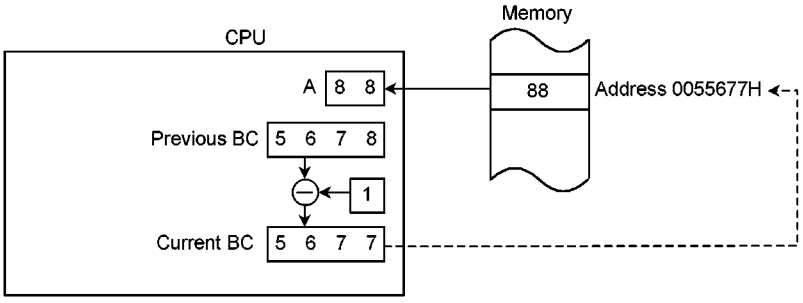
The pre-decrement values are as follows:

When the size of the operand is one byte (8 bits): –1

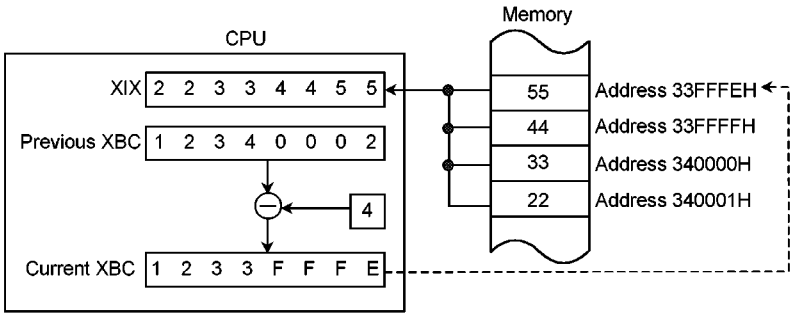
When the size of the operand is one word (16 bits): –2

When the size of the operand is one long word (32 bits): –4

Example 2: LD A, (–BC)...in minimum mode



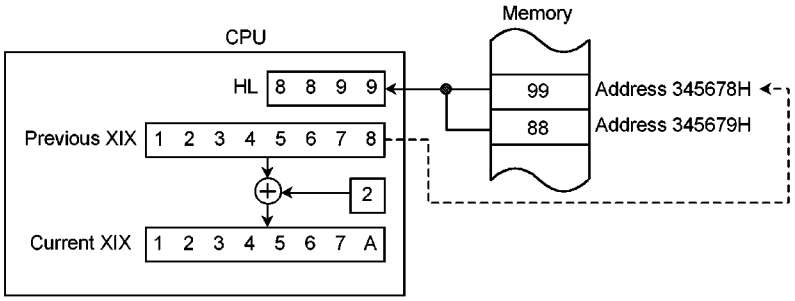
Example 3: LD XIX, (–XBC)...in maximum mode



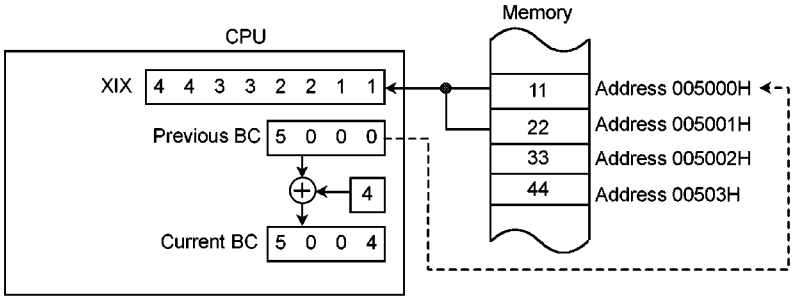
(5) Register Indirect Post-increment Addressing Mode

In this mode, the operand is the memory address specified by the contents of the register. After the operation, the contents of the register are incremented by the size of the operand.

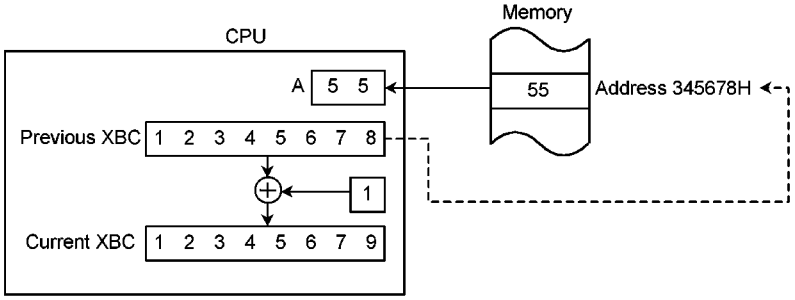
Example 1: LD HL, (XIX+)...in both minimum and maximum modes



Example 2: LD XIX, (BC+)...in minimum mode



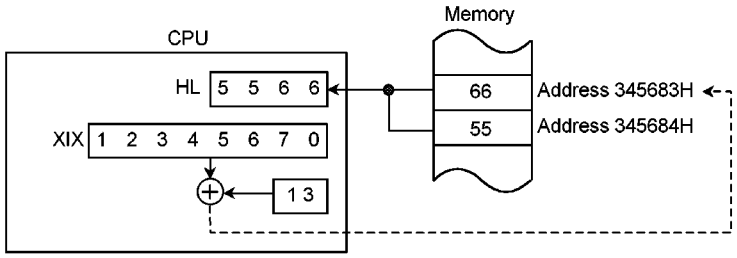
Example 3: LD A, (XBC+)...in maximum mode



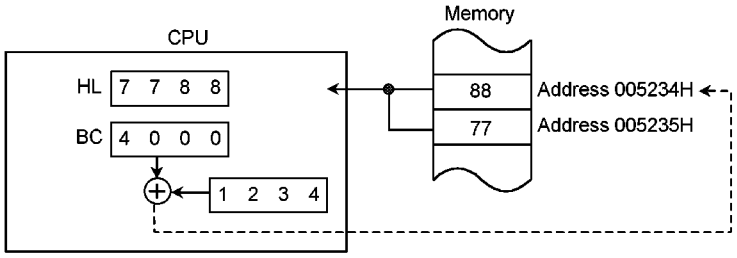
(6) Index Addressing Mode

In this mode, the operand is the memory address obtained by adding the contents of the specified register to the 8- or 16-bit displacement value in the instruction code.

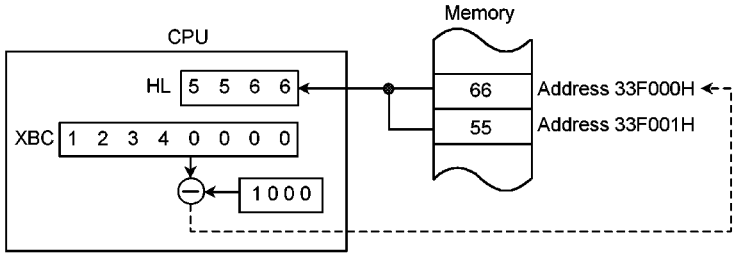
Example 1: LD HL, (XIX + 13H)...in both minimum and maximum modes



Example 2: LD HL, (BC + 1234H)...in minimum mode



Example 3: LD HL, (XBC – 1000H)...in maximum mode

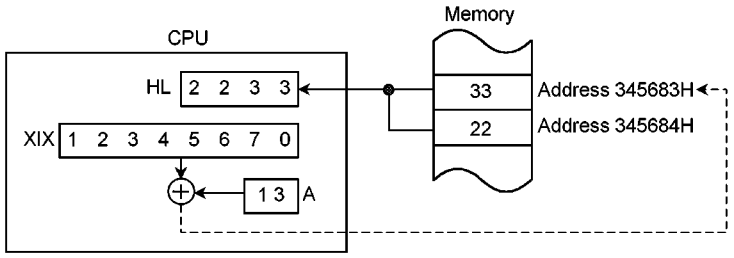


The displacement values range from –8000H to +7FFFH.

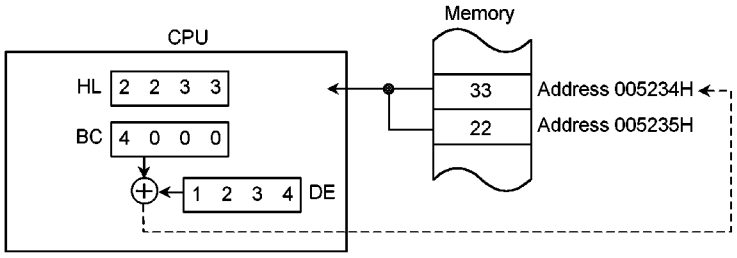
(7) Register Index Addressing Mode

In this mode, the operand is the memory address obtained by adding the contents of the register specified as the base to the register specified as the 8- or 16-bit displacement.

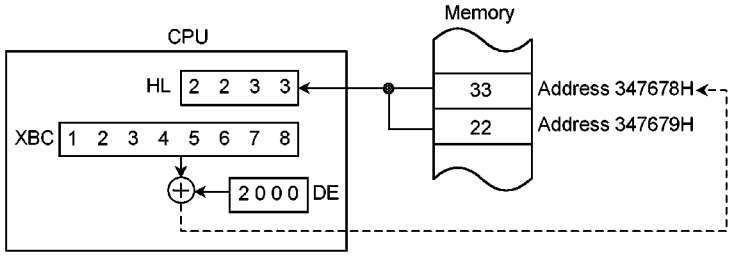
Example 1: LD HL, (XIX + A)...in both minimum and maximum modes



Example 2: LD HL,(BC + DE)...in minimum mode



Example 3: LD HL, (XBC + DE)...in maximum mode

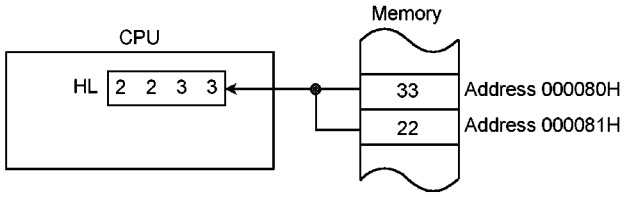


The range of displacement is: –80H to + 7FH in case of 8 bit and –8000H to + 7FFFH in case of 16 bit.

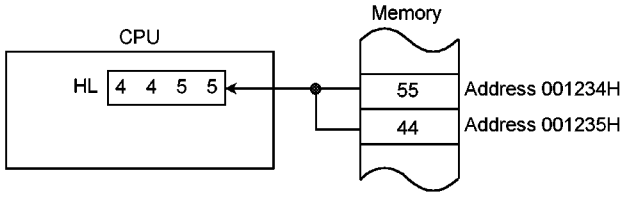
(8) Absolute Addressing Mode

In this mode, the operand is the memory address specified by 1 to 3 bytes in the instruction code. Addresses 000000H to 0000FFH can be specified by 1 byte. Addresses 000000H to 00FFFFH can be specified by 2 bytes. Addresses 000000H to FFFFFFH can be specified by 3 bytes.

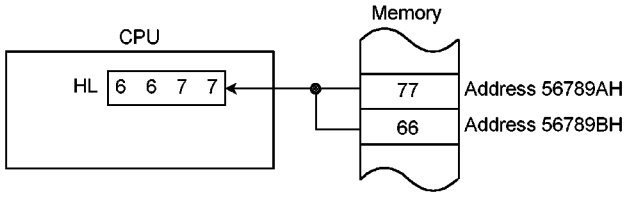
Example 1: LD HL, (80H)



Example 2: LD HL, (1234H)



Example 3: LD HL, (56789AH)



(9) Relative Addressing Mode

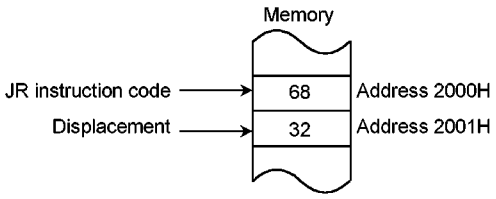
In this mode, the operand is the memory address obtained by adding the 8- or 16-bit displacement value to the address where the instruction code being executed is located.

In this mode, only the following five instructions can be used.

LDAR R, $ + 4 + d16
JR cc, $ + 2 + d8
JRL cc, $ + 3 + d16
CALR $ + 3 + d16
DJNZ r, $ + 3 + d8 ($: start address of instruction code)

In calculating the displacement object code value, the adjustment value (+2 to +4) depends on the instruction type.

Example 1: JR 2034H



In the above example, the displacement object code value is:

2034H – (2000H + 2) = 32H.

5. Instructions

In addition to its various addressing modes, the TLCS-900/L also has a powerful instruction set. The basic instructions are classified into the following nine groups:

- Load instructions (8/16/32 bits)
- Exchange instructions (8/16 bits)
- Block transfer & Block search instructions (8/16 bits)
- Arithmetic operation instructions (8/16/32 bits)
- Logical operation instructions (8/16/32 bits)
- Bit operation instructions (1 bit)
- Special operations, CPU control instructions
- Rotate and Shift instructions (8/16/32 bits)
- Jump, Call, and Return instructions

Table 5 lists the basic instructions of the TLCS-900/L. For details of instructions, see Appendix A; for the instruction list, Appendix B; for the instruction code map, Appendix C; and for the differences between the TLCS-90 and TLCS-900 series, Appendix D.

Table 5.1 TLCS-900/L Basic Instructions

LD	dst, src	Load dst ← src		
PUSH	src	Push src data to stack.		
		SP ← SP – size: (SP) ← src		
POP	dst	Pop data from stack to dst.		
		dst ← (SP): SP ← SP + size		
LDA	dst, src	Load address: set src effective address in dst.		
LDAR	dst, PC + dd	Load address relative:		
		set program counter relative address value in dst. dst ← PC + dd		
EX	dst1, dst2	Exchange dst1 and dst2 data.		
MIRR	dst	Mirror-invert dst bit pattern.		
LDI		Load increment		
LDIR		Load increment repeat		
LDD		Load decrement		
LDDR		Load decrement repeat		
CPI		Compare increment		
CPIR		Compare increment repeat		
CPD		Compare decrement		
CPDR		Compare decrement repeat		
ADD	dst, src	Add	dst ← dst + src	
ADC	dst, src	Add with carry	dst ← dst + src + CY	
SUB	dst, src	Subtract	dst ← dst – src	
SBC	dst, src	Subtract with carry	dst ← dst – src – CY	
CP	dst, src	Compare	dst – src	
AND	dst, src	And	dst ← dst	AND src
OR	dst, src	Or	dst ← dst	OR src
XOR	dst, src	Exclusive-or	dst ← dst	XOR src
INC	imm, dst	Increment	dst ← dst + imm	
DEC	imm, dst	Decrement	dst ← dst – imm	
MUL	dst, src	Multiply unsigned	dst ← dst (low) × src	
MULS	dst, src	Multiply signed	dst ← dst (low) × src	
DIV	dst, src	Divide unsigned		
		dst (low) ← dst ÷ src		
		dst (high) ← remainder		
		V flag set due to division by 0 or overflow.		
DIVS	dst, src	Divide signed		
		dst (low) ← dst ÷ src		
		dst (high) ← remainder: sign is same as that of dividend.		
		V flag set due to division by 0 or overflow.		

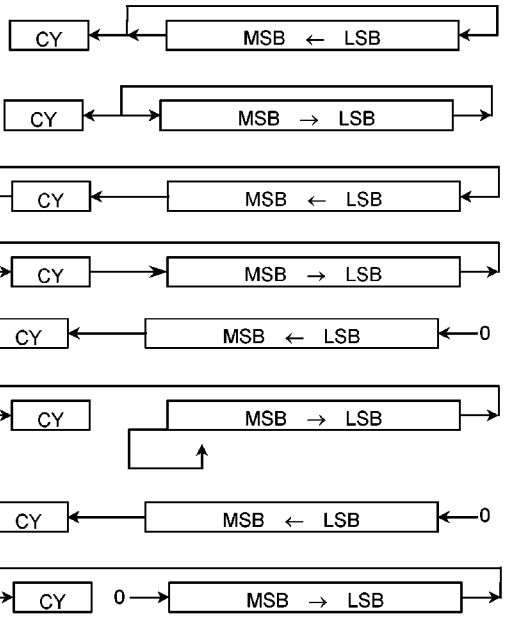
MULA	dst	Multiply and add dst ← dst + (XDE) × (XHL–) 32 bit 32 bit 16 bit 16 bit
MINC1	num, dst	Modulo increment 1
MINC2	num, dst	Modulo increment 2
MINC4	num, dst	Modulo increment 4
MDEC1	num, dst	Modulo decrement 1
MDEC2	num, dst	Modulo decrement 2
MDEC4	num, dst	Modulo decrement 4
NEG	dst	Negate dst ← 0 – dst (Twos complement)
CPL	dst	Complement dst ← not dst (Ones complement)
EXTZ	dst	Extend zero: set upper data of dst to 0.
EXTS	dst	Extend signed: copy the MSB of the lower data of dst to upper data.
DAA	dst	Decimal adjustment accumulator
PAA	dst	Pointer adjustment accumulator: when dst is odd, increment dst by 1 to make it even. if dst (0) = 1 then dst ← dst + 1.
LDCF	bit, src	Load carry flag: copy src<bit> value to C flag.
STCF	bit, dst	Store carry flag: copy C flag value to dst<bit>.
ANDCF	bit, src	And carry flag: and src<bit> value and C flag, then load the result to C flag.
ORCF	bit, src	Or carry flag: or src<bit> and C flag, then load result to C flag.
XORCF	bit, src	Exclusive-or carry flag: exclusive-or src<bit> value and C flag, then load result to C flag.
RCF		Reset carry flag: reset C flag to 0.
SCF		Set carry flag: set C flag to 1.
CCF		Complement carry flag: invert C flag value.
ZCF		Zero flag to carry flag: copy inverted value of Z flag to C flag.
BIT	bit, src	Bit test: Z flag ← not src<bit>
RES	bit, dst	Bit reset
SET	bit, dst	Bit set
CHG	bit, dst	Bit change dst<bit> ← not dst<bit>
TSET	bit, dst	Bit test and set: Z flag ← not dst<bit> dst<bit> ← 1

BS1F	A, dst	Bit search 1 forward: search dst for the first bit set to 1 starting from the LSB, then set the bit number in the A register.
BS1B	A,dst	Bit search 1 backward: search dst for the first bit set to 1 starting fom the MSB, then set the bit number in the A register.
NOP		No operation
MIN		Set CPU to minimum mode
EI	imm	Enable interrupt. IFF ← imm
DI		Disable maskable interrupt. IFF ← 7
PUSH	SR	Push status registers.
POP	SR	Pop status registers.
SWI	imm	Software interrupt PUSH PC&SR: JP 8000H + 10H × imm
HALT		Halt CPU.
LDC	CTRL-REG, reg	Load control: copy the register contents to control register of CPU.
LDC	reg, CTRL-REG	Load control: copy the control register contents to register.
LDX	dst, src	Load extract. dst ← src

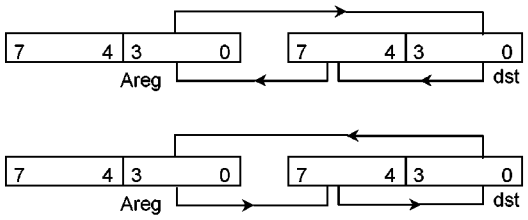
LINK	reg, dd	Link: generate stack frame.
		PUSH reg
		LD reg, XSP
		ADD XSP, dd

UNLK	reg	Unlink: delete stack frame. LD XSP, reg POP reg
LDF	imm	Load register file pointer: specify register bank. RFP ← imm
INCF		Increment register file pointer: move to new register bank. RFP ← RFP + 1
DECF		Decrement register file pointer: return to previous register bank. RFP ← RFP – 1
SCC	cc, dst	Set dst with condition codes. if cc then dst ←1 else dst ← 0.

RLC	num, dst	Rotate left without carry
RRC	num, dst	Rotate right without carry
RL	num, dst	Rotate left
RR	num, dst	Rotate right
SLA	num, dst	Shift left arithmetic
SRA	num, dst	Shift right arithmetic
SLL	num, dst	Shift left logical
SRL	num, dst	Shift right logical



RLD	dst	Rotate left digit
RRD	dst	Rotate right digit



JR	cc, PC+d	Jump relative (8-bit displacement) if cc then PC ← PC + d.
JRL	cc, PC + dd	Jump relative long (16-bit displacement) if cc then PC ← PC + dd.
JP	cc, dst	Jump if cc then PC ← dst.
CALR	RC + dd	Relative call (16-bit displacement) PUSH PC: PC ← PC + dd.
CALL	cc, dst	Call relative if cc then PUSH PC: PC ← dst.
DJNZ	dst, PC + d	Decrement and jump if non-zero dst ← dst – 1 if dst ≠ 0 then PC ← PC + d.
RET	cc	Return if cc then POP PC.
RETD	dd	Return and deallocate RET XSP ← XSP + dd
RETI		Return from interrupt POP SR&PC

Table 5.2 Instruction List

BWL	LD	reg, reg	BWL	INC	imm3, reg	---	NOP	
BWL	LD	reg, imm		DEC	imm3, mem.B/W			
BWL	LD	reg, mem				---	*MIN	
BWL	LD	mem, reg				---	EI	[imm3]
BW-	LD	mem, imm				---	DI	
BW-	LD	(nn), mem	BW-	MUL	reg, reg	-W-	*PUSH	SR
BW-	LD	mem, (nn)		*MULS	reg, imm	-W-	*POP	SR
				DIV	reg, mem	---	SWI	[imm3]
				*DIVS		---	HALT	
BWL	PUSH	reg/F				BWL	*LDC	CTRL – R, reg
BW-	PUSH	imm	-W-	*MULA	reg	BWL	*LDC	reg, CTRL – R
BW-	PUSH	mem				B--	*LDX	(n), n
			-W-	*MINC1	imm, reg			
BWL	POP	reg/F	-W-	*MINC2	imm, reg	--L	*LINK	reg, dd
BW-	POP	mem	-W-	*MINC4	imm, reg	--L	*UNLK	reg
			-W-	*MDEC1	imm, reg	---	*LDF	imm3
			-W-	*MDEC2	imm, reg	---	*INCF	
-WL	LDA	reg, mem	-W-	*MDEC4	imm, reg	---	*DECF	
-WL	LDAR	reg, PC+dd				BW-	*SCC	cc, reg
			BW-	NEG	reg			
			BW-	CPL	reg	BWL	RLC	imm, reg
			-WL	*EXTZ	reg		RRC	A, reg
B--	EX	F, F'	-WL	*EXTS	reg		RL	mem. B/W
BW-	EX	reg, reg	B--	DAA	reg		RR	
BW-	EX	mem, reg	-WL	*PAA	reg		SLA	
							SRA	
							SLL	
-W-	*MIRR	reg	BW-	*LDCF	imm, reg		SRL	
				*STCF	A, reg			
				*ANDCF	imm, mem.B	B--	RLD	[A,] mem
				*ORCF	A, mem.B	B--	RRD	[A,] mem
BW-	LDI			*XORCF				
BW-	LDIR							
BW-	LDD		---	RCF		---	JR	[cc,] PC + d
BW-	LDDR		---	SCF		---	JRL	[cc,] PC + dd
			---	CCF		---	JP	[cc,] mem
			---	*ZCF		---	CALR	PC + dd
BW-	CPI					---	CALL	[cc,] mem
BW-	CPIR		BW-	BIT	imm, reg			
BW-	CPD			RES	imm, mem.B	BW-	DJNZ	[reg], PC + d
BW-	CPDR			SET				
				*CHG		---	RET	[cc]
				TSET		---	*RETD	dd
BWL	ADD	reg, reg				---	RETI	
	ADC	reg, imm	-W-	*BS1F	A, reg			
	SUB	reg, mem		*BS1B				
	SBC	mem, reg						
	CP	mem, imm.B/W						
	AND							
	OR							
	XOR							

B = Byte (8 bit), W = Word (16 bit), L = Long – Word (32 bit).

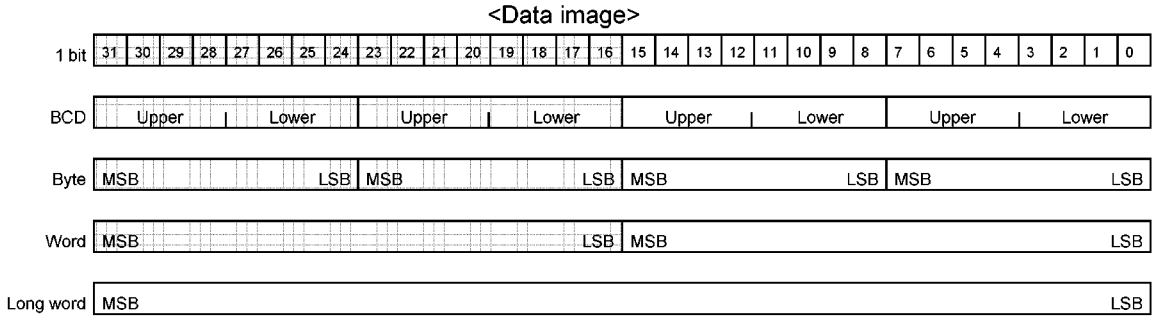
* : Indicates instruction added to the TLCS-90 series.

[ ] : Indicates can be omitted.

6. Data Formats

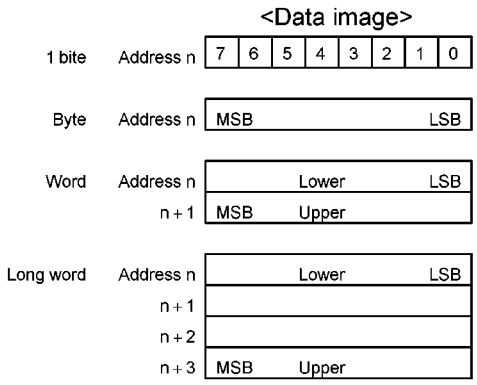
The TLCS-900/L can handle 1/4/8/16/32-bit data.

(1) Register Data Format



Note 1: To access the parts indicated by ▭, the instruction code is one byte longer than when accessing the other parts.

(2) Memory Data Format



Note 2: There are no restrictions on the location of word or long word data in memory. They can be located from even or odd numbered address.

Note 3: When the PUSH instruction is used to save data to the stack area, the stack pointer is decremented, then the data is saved.

Example: PUSH HL; XSP ← XSP – 2
(XSP) ← L
(XSP + 1) ← H

This is the same in register indirect pre-decrement mode. The order is reversed in the TLCS-90 series: data is saved first, then the stack pointer is decremented.

Example: PUSH HL; (XSP – 1) ← H
(XSP – 2) ← L
XSP ← XSP – 2

(3) Dynamic Bus Sizing

The TLCS-900/L can switch between 8- and 16-bit data buses dynamically during each bus cycle. This is called dynamic bus sizing. The function enables external memory extension using both 8- and 16-bit data bus memories. Products with a built-in chip select/wait controller can control external data bus size for each address area.

Table 6.1 Dynamic Bus Sizing

Operand data size	Operand start address	Data size at memory side	CPU address	CPU data	
				D15 to D8	D7 to D0
8 bits	2n + 0 (even)	8 bits	2n + 0	xxxxx	b7 to b0
		16 bits	2n + 0	xxxxx	b7 to b0
	2n + 1 (odd)	8 bits	2n + 1	xxxxx	b7 to b0
		16 bits	2n + 1	b7 to b0	xxxxx
16 bits	2n + 0 (even)	8 bits	2n + 0	xxxxx	b7 to b0
			2n + 1	xxxxx	b15 to b8
		16 bits	2n + 0	b15 to b8	b7 to b0
	2n + 1 (odd)	8 bits	2n + 1	xxxxx	b7 to b0
			2n + 2	xxxxx	b15 to b8
		16 bits	2n + 1	b7 to b0	xxxx
			2n + 2	xxxxx	b15 to b8
32 bits	2n + 0 (even)	8 bits	2n + 0	xxxxx	b7 to b0
			2n + 1	xxxxx	b15 to b8
			2n + 2	xxxxx	b23 to b16
			2n + 3	xxxxx	b31 to b24
		16 bits	2n + 0	b15 to b8	b7 to b0
			2n + 2	b31 to b24	b23 to b16
	2n + 1 (odd)	8 bits	2n + 1	xxxxx	b7 to b0
			2n + 2	xxxxx	b15 to b8
			2n + 3	xxxxx	b23 to b16
			2n + 4	xxxxx	b31 to b24
		16 bits	2n + 1	b7 to b0	xxxxx
			2n + 2	b23 to b16	b15 to b8
			2n + 4	xxxx	b31 to b24

xxxxx: During read, indicates the data input to the bus are ignored. During write, indicates the bus is at high impedance and the write strobe signal is non-active.

(4) Internal Data Bus Format

With the TLCS-900/L, the CPU and the internal memory (built-in ROM or RAM) are connected via a 16-bit internal data bus. The internal memory operates with 0 wait. The CPU and the built-in I/Os are connected using an 8-bit internal data bus. This is because the built-in I/O access speed has little influence on the overall system operation speed.

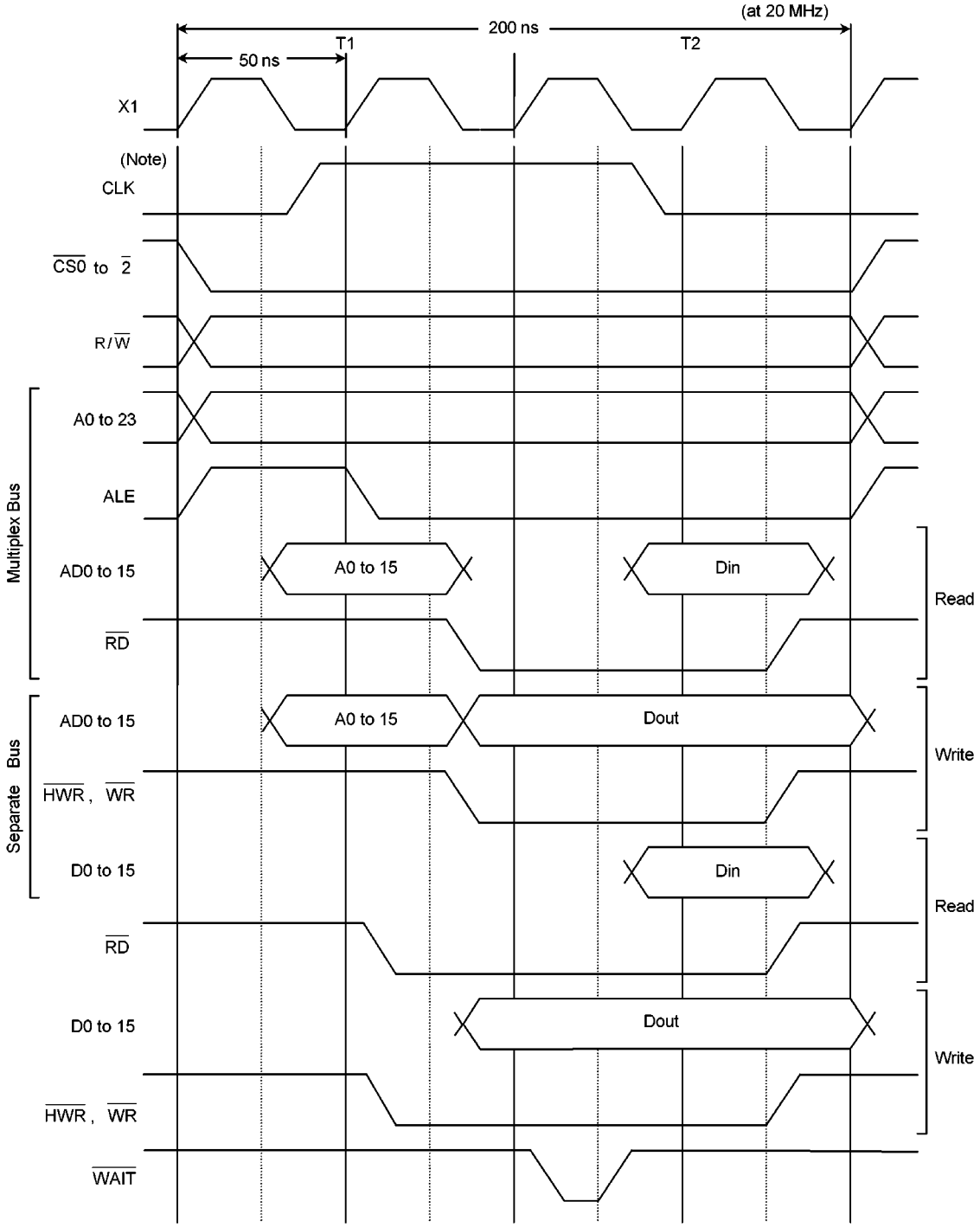
Overall system operation speed depends largely on the speed of program memory access. The built-in I/O operates in sync with the signal phase of the CLK pin. It is synchronized so that the CLK rises (⌐¬) in the middle of the bus cycle. (Figure7.1 shouws signal phases.) If the CLK is "1" when the ALE signal rises, 1 wait is inserted automatically for synchronization.

7. Basic Timings

The TLCS-900/L runs the following basic timings.

- Read cycle
- Write cycle
- Dummy cycle
- Interrupt receive timing
- Reset

Figure 7.1 to Figure 7.8 show the basic timings.



Note: CLK outputs are not always the same as the above phases.

Figure 7.1 0 WAIT Read/Write Cycle

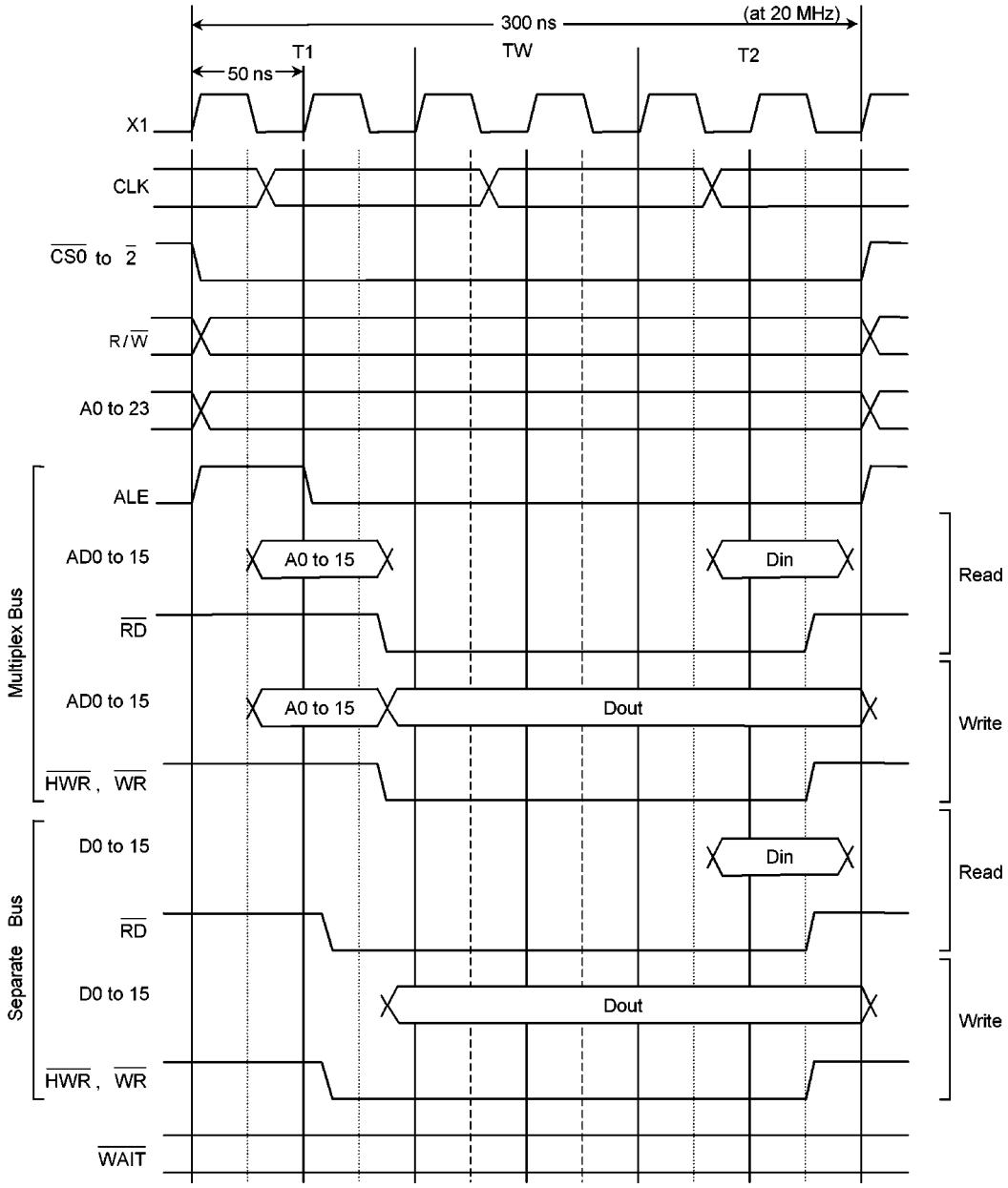


Figure 7.2 1 WAIT Read/Write Cycle

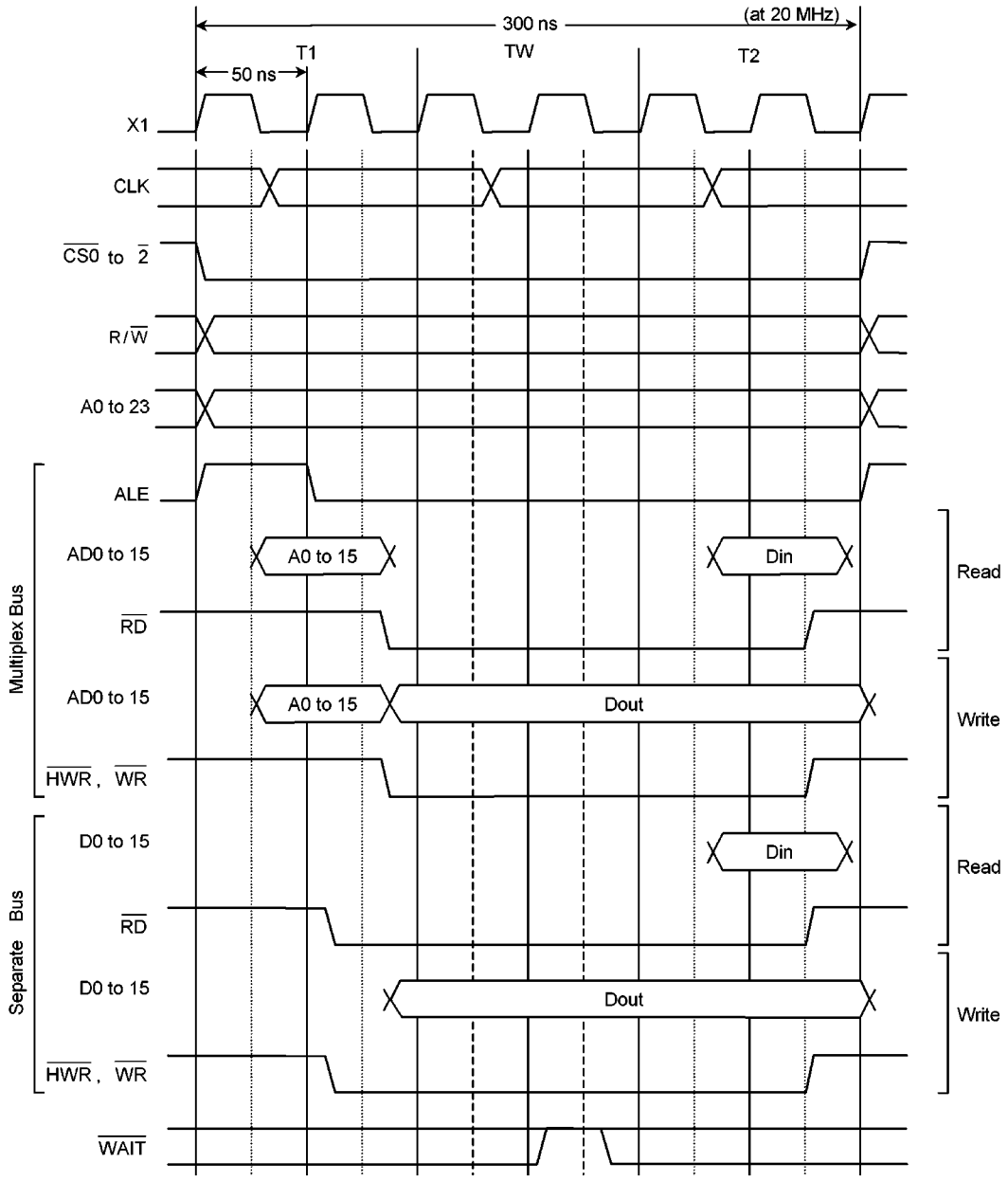


Figure 7.3 1 WAIT + n Read/Write Cycle (n = 0)

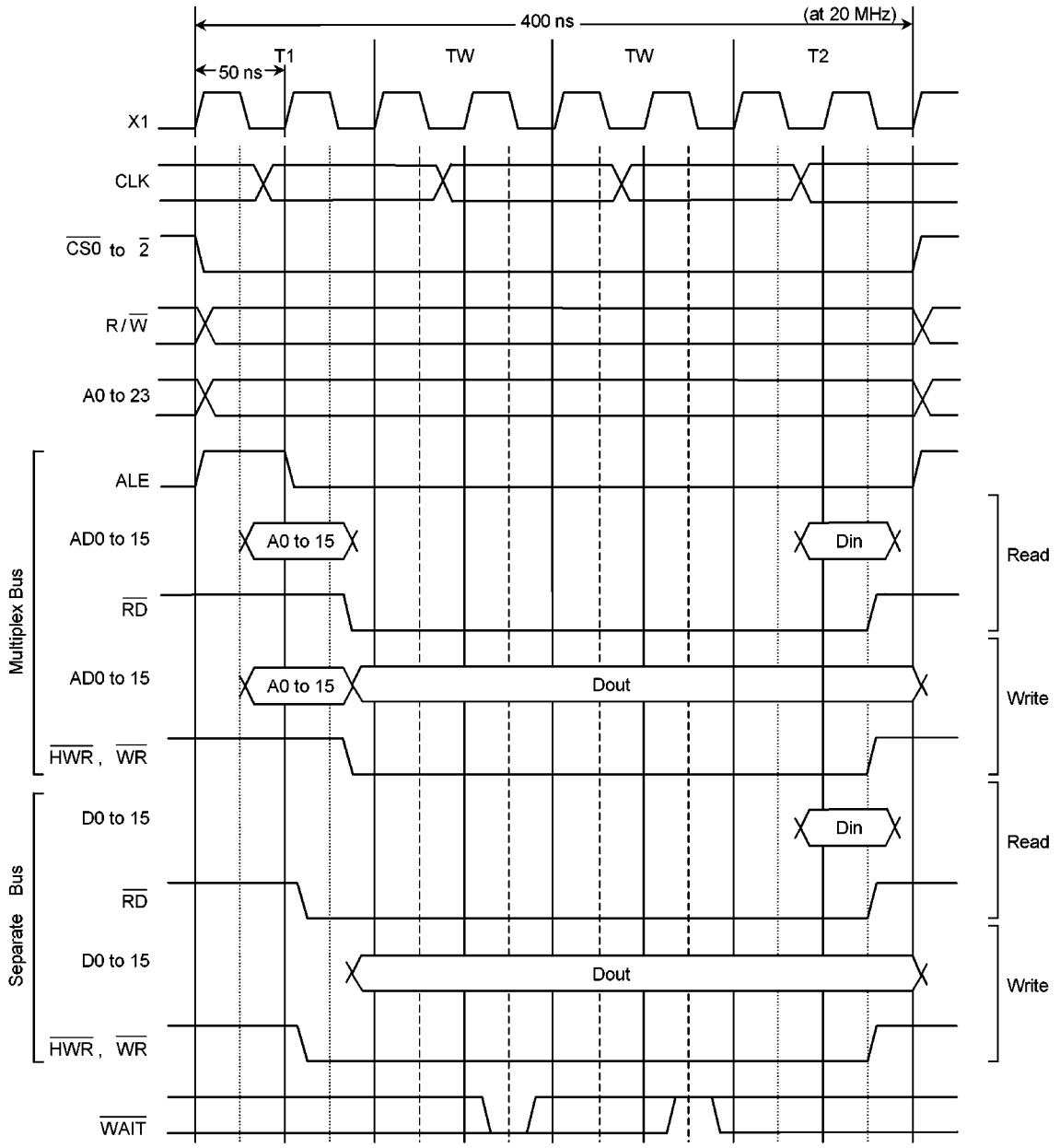


Figure 7.4 1 WAIT + n Read/Write Cycle (n = 1)

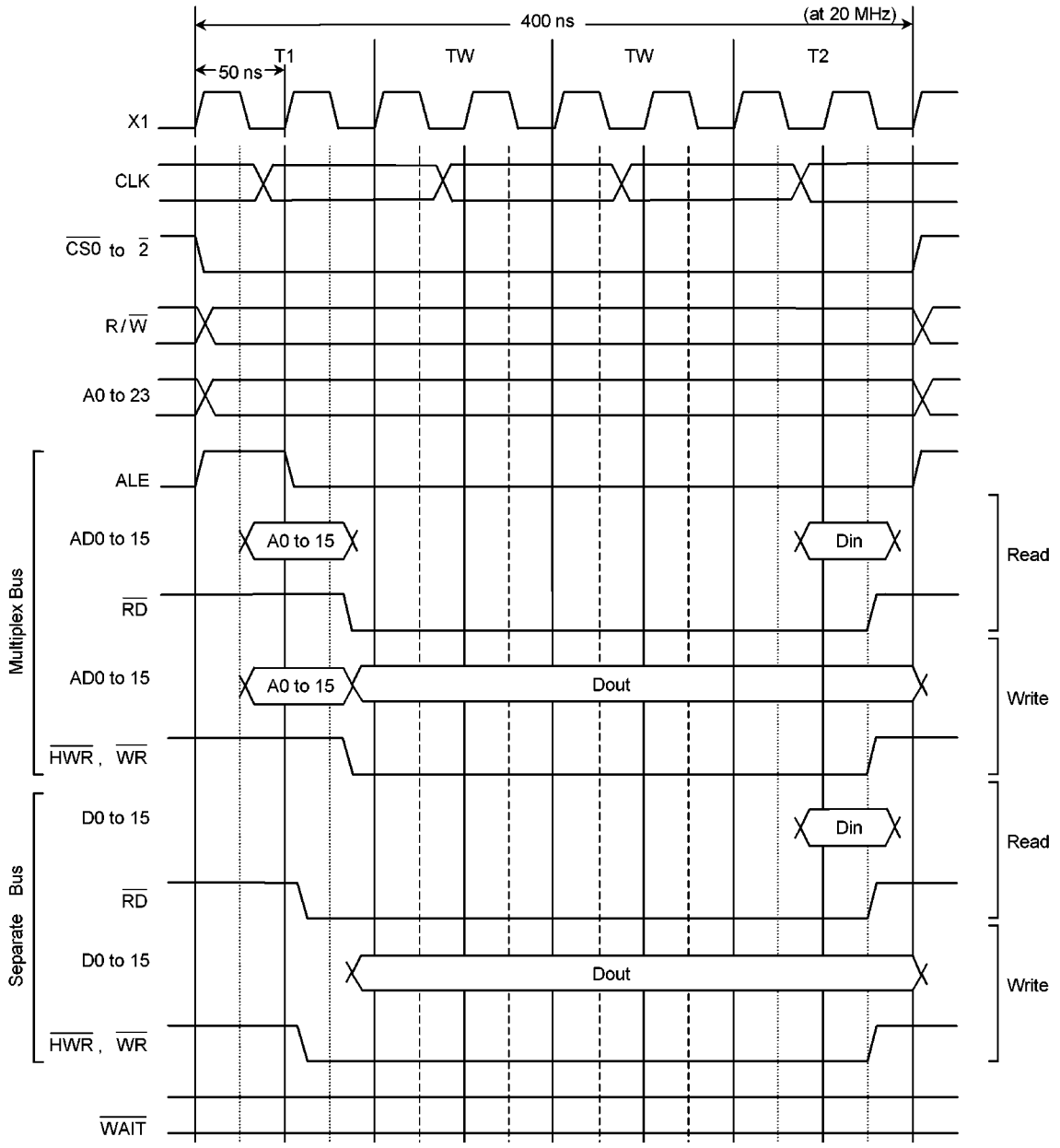


Figure 7.5 2 WAIT Read/Write Cycle

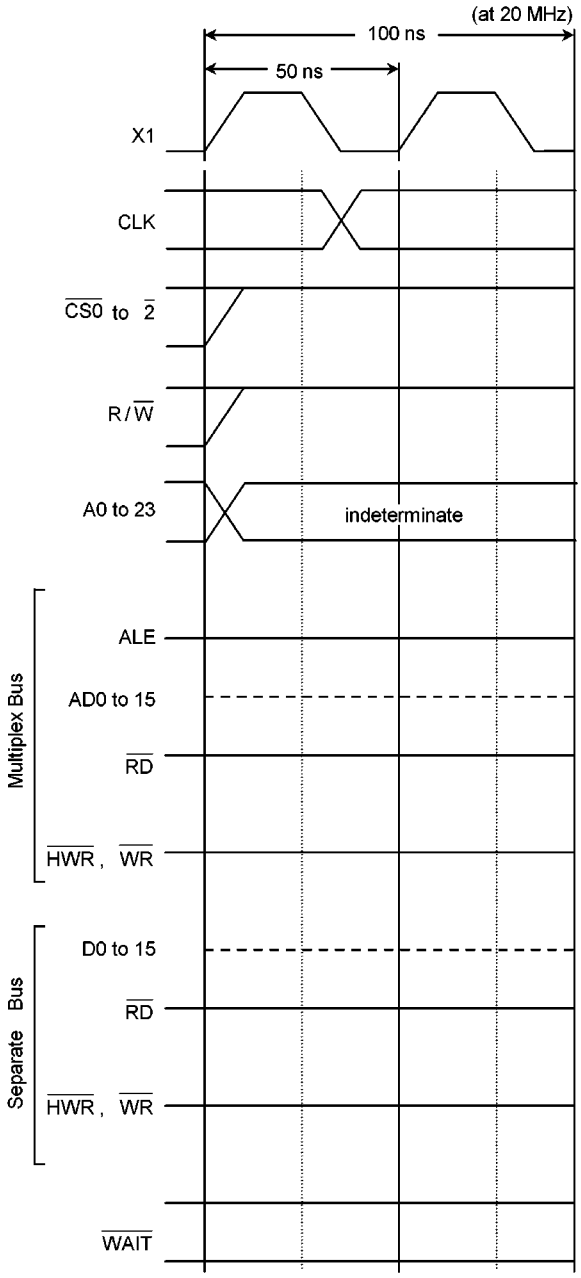
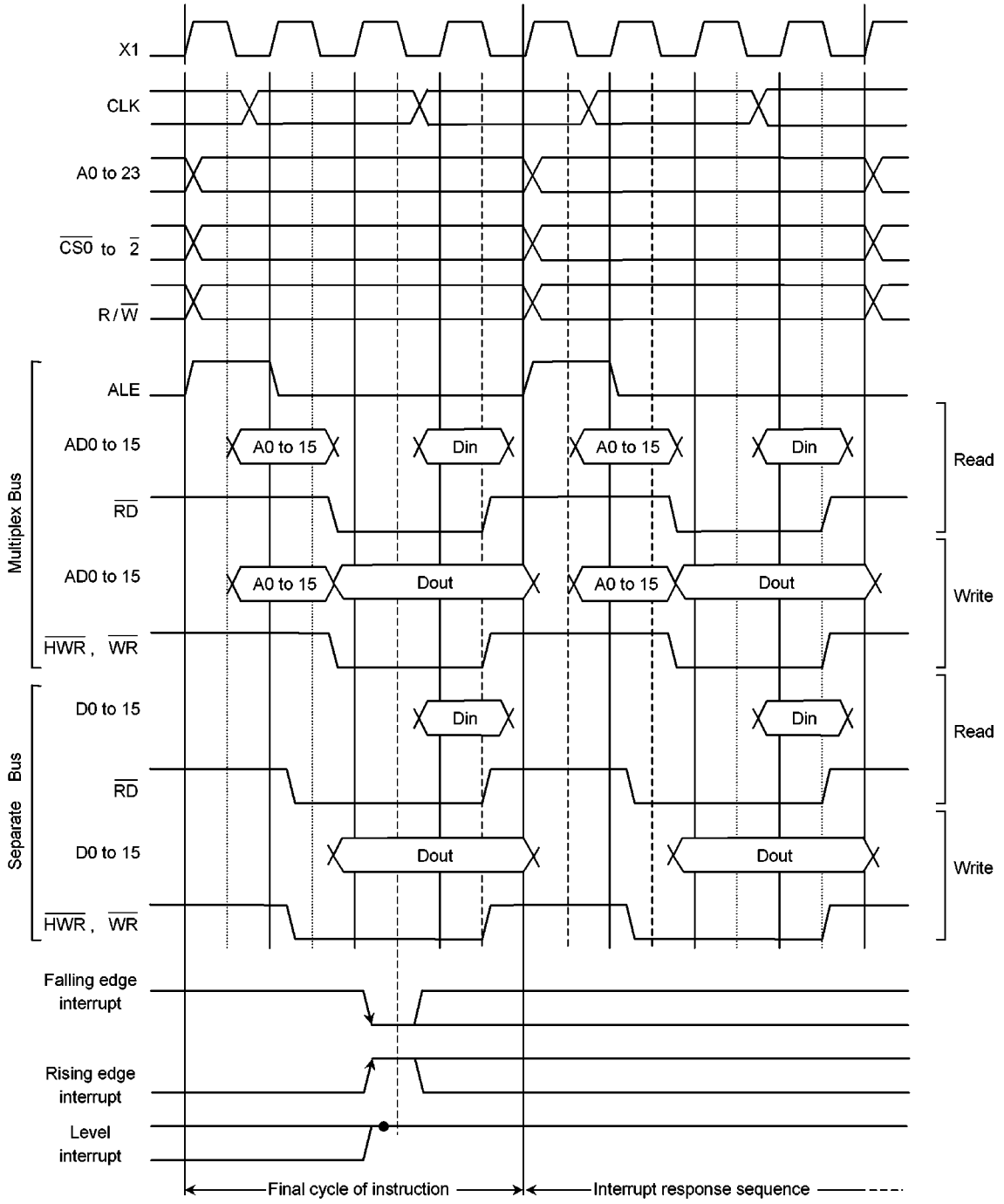


Figure 7.6 1 State Dummy Cycle



Note: This timing chart is a theoretical example. In practice, due to the operation of the bus interface unit in the CPU, external bus and internal interrupt receive timings do not correspond one to one.

Figure 7.7 Interrupt Receive Timing

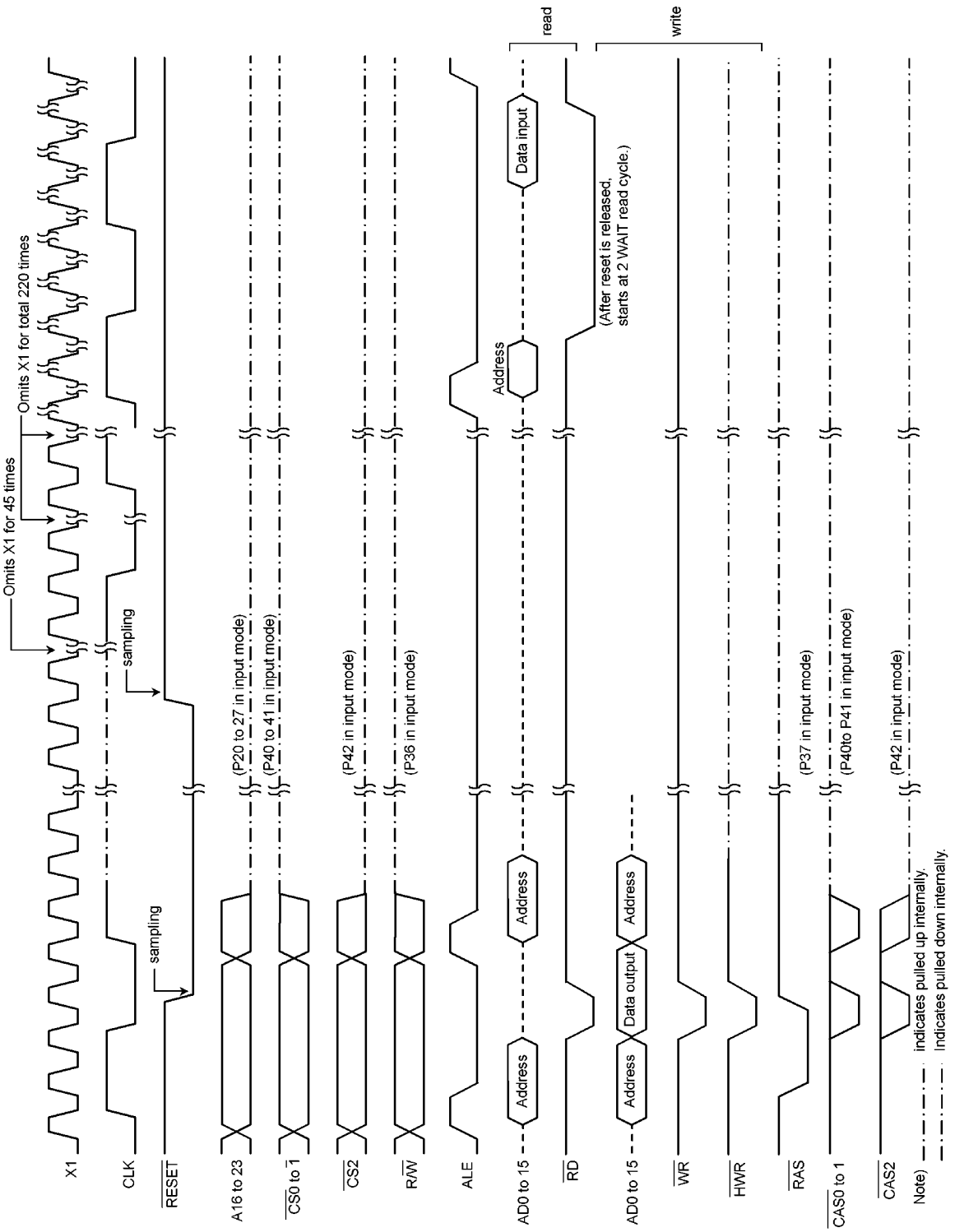


Figure 7.8 Reset Timings (external ROM operation: TMP95C061)

Appendix A: Details of Instructions

- Instruction List

1. Load

LD	PUSH	POP	LDA	LDAR

2. Exchange

EX	MIRR

3. Load Increment/Decrement & Compare Increment/Decrement

LDI	LDIR	LDD	LDDR	CPI	CPIR	CPD	CPDR

4. Arithmetic operations

ADD	ADC	SUB	SBC	CP	INC	DE	NEG
EXTZ	EXTS	DAA	PAA	MUL	MULS	DIV	DIVS
MULA	MINC	MDEC					

5. Logical operations

AND	OR	XOR	CPL

6. Bit operations

LDCF	STCF	ANDCF	ORCF	XORCF	RCF	SCF	CCF
ZCF	BIT	RES	SET	CHG	TSET	BS1	

7. Special operations and CPU control

NOP	MIN	EI	DI	PUSH–SR	POP–SR	SWI	HALT
LDC	LDX	LINK	UNLK	LDF	INCF	DECF	SCC

8. Rotate and shift

RLC	RRC	RL	RR	SLA	SRA	SLL	SRL
RLD	RRD						

9. Jump, call, and return

JP	JR	JRL	CALL	CALR	DJNZ	RET	RETD
RETI							

- Explanations of symbols used in this document

dst	Destination: destination of data transfer or operation result load.
src	Source: source of data transfer or operation data read.
num	Number: numerical value.
condition	Condition: based on flag status.
R	Eight general–purpose registers including 8/16/32-bit current bank registers. 8-bit registers : W, A, B, C, D, E, H, L (only eight registers) 16-bit registers : WA, BC, DE, HL, IX, IY, IZ, SP (only eight registers) 32-bit registers : XWA, XBC, XDE, XHL, XIX, XIY, XIZ, XSP (only eight registers)
r	8/16/32-bit general-purpose registers } (Please refer to " Register map" on page CPU900-54, 55.)
r16	16-bit general-purpose registers
r32	32-bit general-purpose registers
cr	All 8/16/32-bit CPU control registers DMAS0 to 3, DMAD0 to 3, DMAC0 to 3, DMAM0 to 3, INTNEST
A	A register (8 bits)
F	Flag registers (8 bits)
F'	Inverse flag registers (8 bits)
SR	Status registers (16 bits)
PC	Program counter (in minimum mode, 16 bits; in maximum mode, 32 bits)
(mem)	8/16/32-bit memory data
mem	Effective address value
<W>	When the operand size is a word, W must be specified.
[ ]	Operands enclosed in square brackets can be omitted.
#	8/16/32-bit immediate data.
#3	3-bit immediate data : 0 to 7 or 1 to 8 ... for abbreviated codes.
#4	4-bit immediate data : 0 to 15 or 1 to 16
d8	8-bit displacement : –80H to + 7FH
d16	16-bit displacement : –8000H to + 7FFFH
cc	Condition code
CY	Carry flag
Z	Zero flag
(#8)	Direct addressing: (00H) to (0FFH) ... 256–byte area
(#16)	64K–byte area addressing: (0000H) to (0FFFFH)
(–r32)	Pre–decrement addressing
(r32+)	Post–increment addressing
$	Start address of instruction

- Explanations of symbols in object codes

z
zz
zzz
s

Operand size specify code

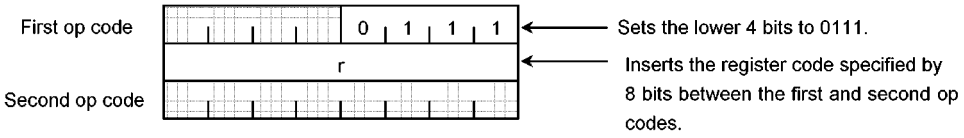
	Byte	Word	Long word
z	0	1	–
zz	00	01	10
zzz	010	011	100
s	–	0	1

R
r

Register specify code

Code	Byte	Word	Long word
000	W	WA	XWA
001	A	BC	XBC
010	B	DE	XDE
011	C	HL	XHL
100	D	IX	XIX
101	E	IY	XIY
110	H	IZ	XIZ
111	L	SP	XSP

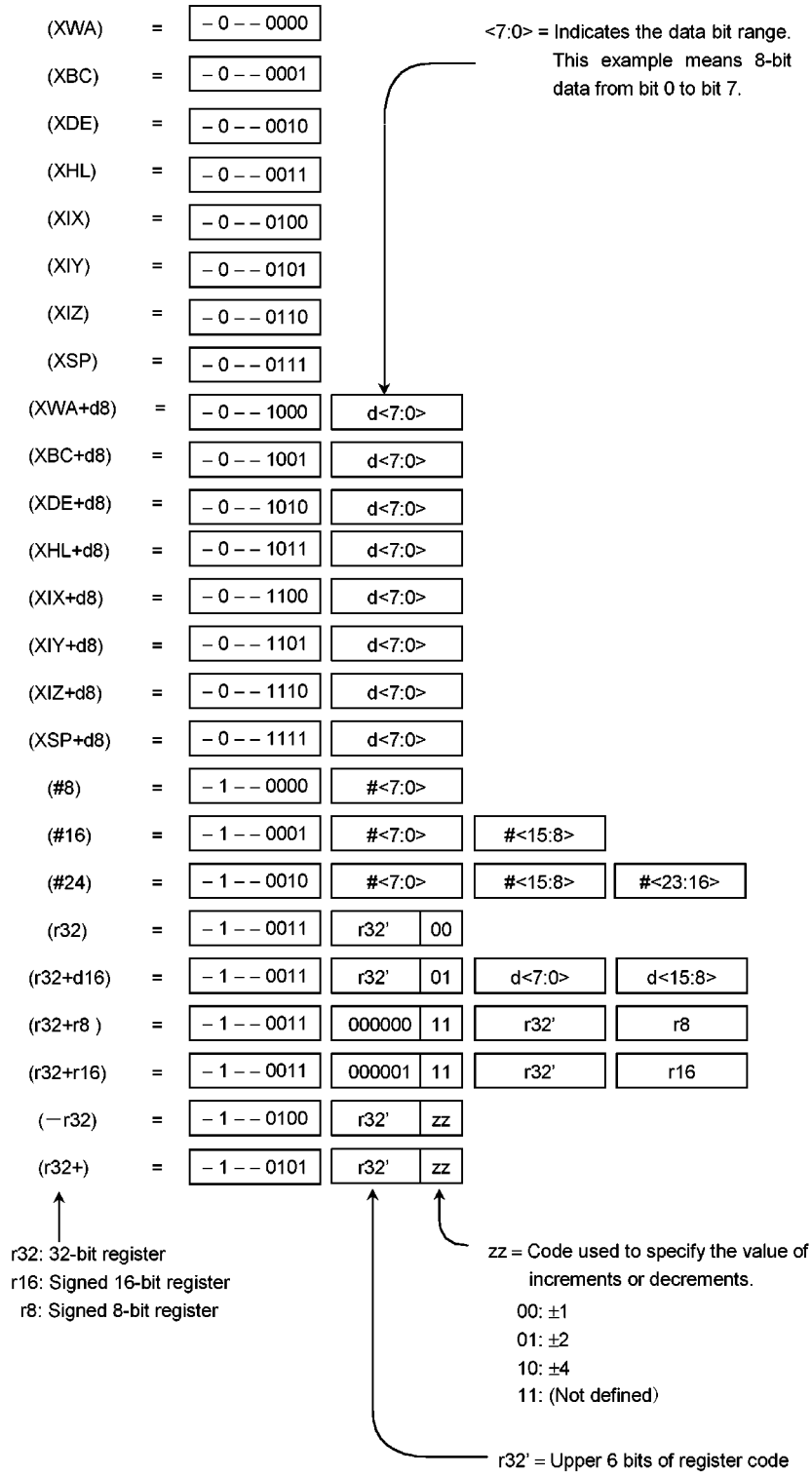
Note: In addition to the above, all registers can be specified by "r" using extension codes. In this case, the number of execution states increases by 1. The format is shown below.



The code value in "r" must be:
Multiple of 2, if accessed as a word register.
Multiple of 4, if accessed as a long word.
For registers specified by 8 bits, see Register Maps.

mem

Memory addressing mode specify code



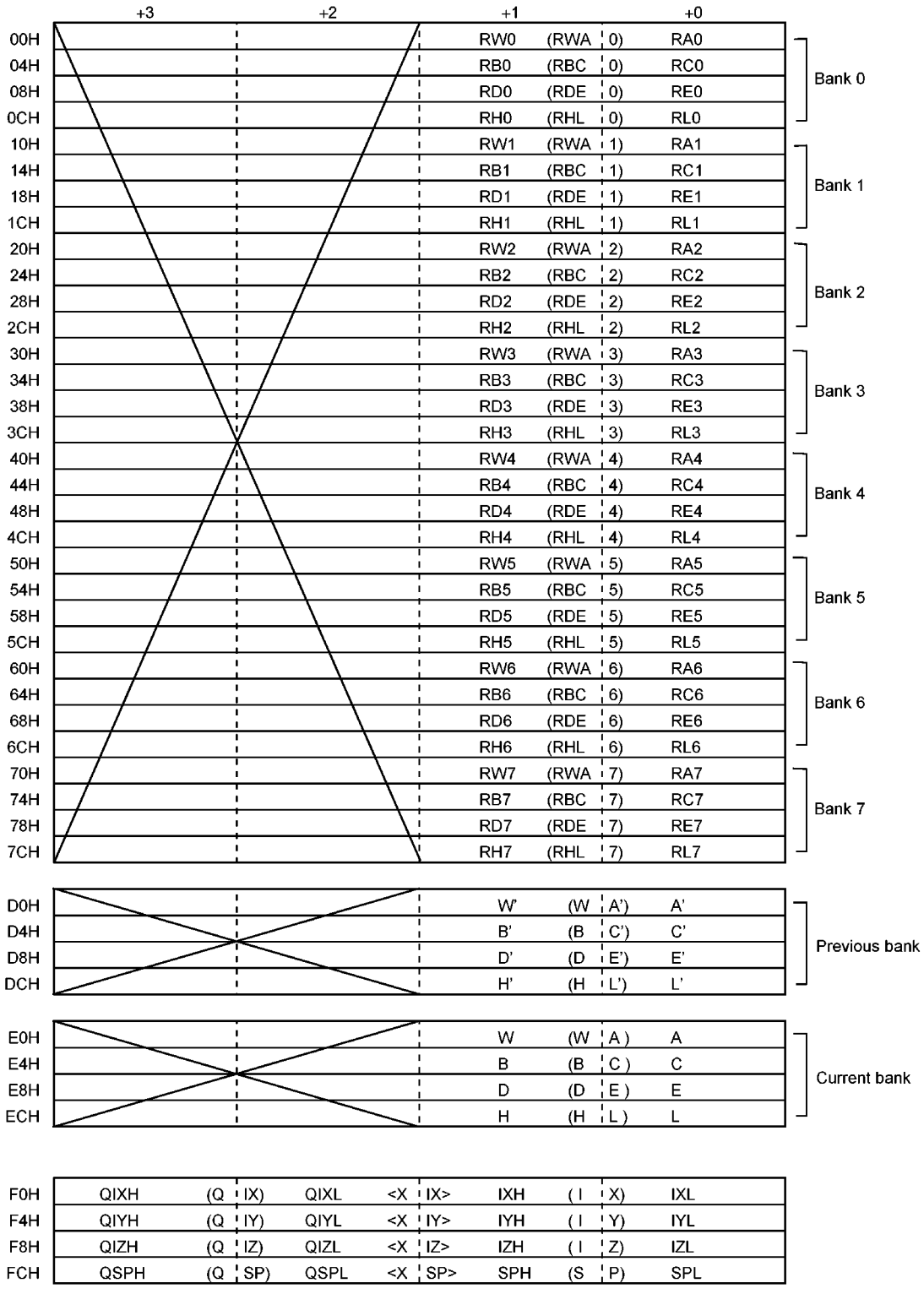
cc — Condition codes

Code	Symbol	Description	Conditional expression
0000	F	always False	–
1000	(none)	always True	–
0110	Z	Zero	Z = 1
1110	NZ	Not Zero	Z = 0
0111	C	Carry	C = 1
1111	NC	Not Carry	C = 0
1101	PL or P	PLus	S = 0
0101	MI or M	MInus	S = 1
1110	NE	Not Equal	Z = 0
0110	EQ	EQual	Z = 1
0100	OV	OVerflow	P/V = 1
1100	NOV	No OVerflow	P/V = 0
0100	PE	Parity is Even	P/V = 1
1100	PO	Parity is Odd	P/V = 0
1001	GE	Greater than or Equal (signed)	(S xor P/V) = 0
0001	LT	Less Than (signed)	(S xor P/V) = 1
1010	GT	Greater Than (signed)	[Z or (S xor P/V)] = 0
0010	LE	Less than or Equal (signed)	[Z or (S xor P/V)] = 1
1111	UGE	Unsigned Greater than or Equal	C = 0
0111	ULT	Unsigned Less Than	C = 1
1011	UGT	Unsigned Greater Than	(C or Z) = 0
0011	ULE	Unsigned Less than or Equal	(C or Z) = 1

- Flag changes

0	Reset to “0”.
1	Set to “1”.
–	No change.
*	“0” or “1” depending on the result of the calculation.
X	Indeterminate value.
P	Parity result is set.
V	Overflow result is set.

- Register map "r" (minimum mode)



	+3	+2	+1			+0	
00H			RW0	(RWA	0)	RA0	Bank 0
04H			RB0	(RBC	0)	RC0	
08H			RD0	(RDE	0)	RE0	
0CH			RH0	(RHL	0)	RL0	
10H			RW1	(RWA	1)	RA1	Bank 1
14H			RB1	(RBC	1)	RC1	
18H			RD1	(RDE	1)	RE1	
1CH			RH1	(RHL	1)	RL1	
20H			RW2	(RWA	2)	RA2	Bank 2
24H			RB2	(RBC	2)	RC2	
28H			RD2	(RDE	2)	RE2	
2CH			RH2	(RHL	2)	RL2	
30H			RW3	(RWA	3)	RA3	Bank 3
34H			RB3	(RBC	3)	RC3	
38H			RD3	(RDE	3)	RE3	
3CH			RH3	(RHL	3)	RL3	
40H			RW4	(RWA	4)	RA4	Bank 4
44H			RB4	(RBC	4)	RC4	
48H			RD4	(RDE	4)	RE4	
4CH			RH4	(RHL	4)	RL4	
50H			RW5	(RWA	5)	RA5	Bank 5
54H			RB5	(RBC	5)	RC5	
58H			RD5	(RDE	5)	RE5	
5CH			RH5	(RHL	5)	RL5	
60H			RW6	(RWA	6)	RA6	Bank 6
64H			RB6	(RBC	6)	RC6	
68H			RD6	(RDE	6)	RE6	
6CH			RH6	(RHL	6)	RL6	
70H			RW7	(RWA	7)	RA7	Bank 7
74H			RB7	(RBC	7)	RC7	
78H			RD7	(RDE	7)	RE7	
7CH			RH7	(RHL	7)	RL7	
D0H			W'	(W	A')	A'	Previous bank
D4H			B'	(B	C')	C'	
D8H			D'	(D	E')	E'	
DCH			H'	(H	L')	L'	
E0H			W	(W	A)	A	Current bank
E4H			B	(B	C)	C	
E8H			D	(D	E)	E	
ECH			H	(H	L)	L	

F0H	QIXH	(Q	IX)	QIXL	<X	IX>	IXH	(I	X)	IXL
F4H	QIYH	(Q	IY)	QIYL	<X	IY>	IYH	(I	Y)	IYL
F8H	QIZH	(Q	IZ)	QIZL	<X	IZ>	IZH	(I	Z)	IZL
FCH	QSPH	(Q	SP)	QSPL	<X	SP>	SPH	(S	P)	SPL

( ): Word register name (16 bits)

< >: Long word register name (32 bits)

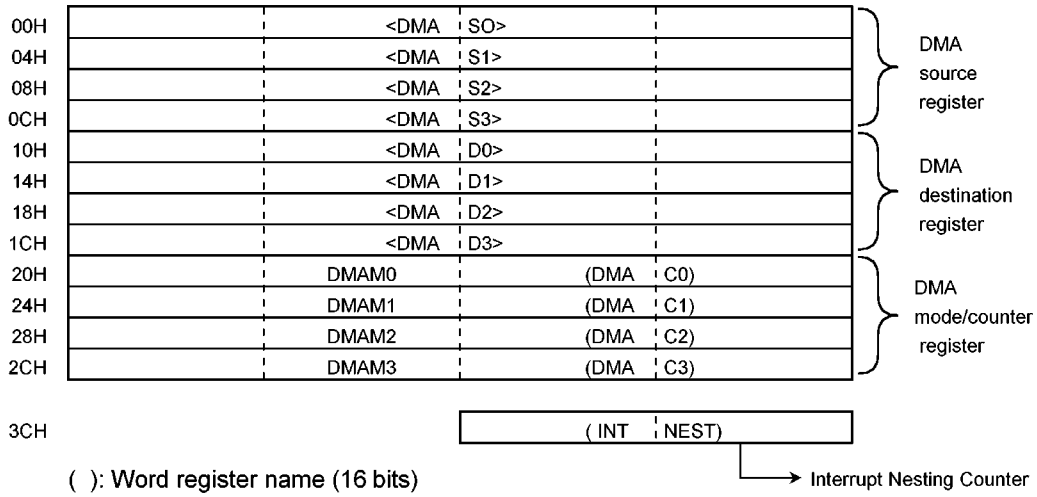
- Register map "r" (maximum mode)

	+3			+2			+1			+0		
00H	QWO	(QWA	0)	QAO	<XWA	0>	RW0	(RWA	0)	RA0		Bank 0
04H	QBO	(QBC	0)	QCO	<XBC	0>	RB0	(RBC	0)	RC0		
08H	QDO	(QDE	0)	QEO	<XDE	0>	RD0	(RDE	0)	RE0		
0CH	QHO	(QHL	0)	QLO	<XHL	0>	RH0	(RHL	0)	RL0		
10H	QW1	(QWA	1)	QA1	<XWA	1>	RW1	(RWA	1)	RA1		Bank 1
14H	QB1	(QBC	1)	QC1	<XBC	1>	RB1	(RBC	1)	RC1		
18H	QD1	(QDE	1)	QE1	<XDE	1>	RD1	(RDE	1)	RE1		
1CH	QH1	(QHL	1)	QL1	<XHL	1>	RH1	(RHL	1)	RL1		
20H	QW2	(QWA	2)	QA2	<XWA	2>	RW2	(RWA	2)	RA2		Bank 2
24H	QB2	(QBC	2)	QC2	<XBC	2>	RB2	(RBC	2)	RC2		
28H	QD2	(QDE	2)	QE2	<XDE	2>	RD2	(RDE	2)	RE2		
2CH	QH2	(QHL	2)	QL2	<XHL	2>	RH2	(RHL	2)	RL2		
30H	QW3	(QWA	3)	QA3	<XWA	3>	RW3	(RWA	3)	RA3		Bank 3
34H	QB3	(QBC	3)	QC3	<XBC	3>	RB3	(RBC	3)	RC3		
38H	QD3	(QDE	3)	QE3	<XDE	3>	RD3	(RDE	3)	RE3		
3CH	QH3	(QHL	3)	QL3	<XHL	3>	RH3	(RHL	3)	RL3		
40H												
44H												
48H												
4CH												
50H												
54H												
58H												
5CH												
60H												
64H												
68H												
6CH												
70H												
74H												
78H												
7CH												
D0H	QW'	(Q	WA')	QA'	<X	WA'>	W'	(W	A')	A'		Previous bank
D4H	QB'	(Q	BC')	QC'	<X	BC'>	B'	(B	C')	C'		
D8H	QD'	(Q	DE')	QE'	<X	DE'>	D'	(D	E')	E'		
DCH	QH'	(Q	HL')	QL'	<X	HL'>	H'	(H	L')	L'		
E0H	QW	(Q	WA)	QA	<X	WA >	W	(W	A)	A		Current bank
E4H	QB	(Q	BC)	QC	<X	BC >	B	(B	C)	C		
E8H	QD	(Q	DE)	QE	<X	DE >	D	(D	E)	E		
ECH	QH	(Q	HL)	QL	<X	HL >	H	(H	L)	L		
F0H	QIXH	(Q	IX)	QIXL	<X	IX>	IXH	(I	X)	IXL		
F4H	QIYH	(Q	IY)	QIYL	<X	IY>	IYH	(I	Y)	IYL		
F8H	QIZH	(Q	IZ)	QIZL	<X	IZ>	IZH	(I	Z)	IZL		
FCH	QSPH	(Q	SP)	QSPL	<X	SP>	SPH	(S	P)	SPL		

( ): Word register name (16 bits)

< >: Long word register name (32 bits)

- Control register map “cr”



( ): Word register name (16 bits)

< >: Long word register name (32 bits)

ADC dst, src

<Add with Carry>

Operation: dst ← dst + src + CY

Description: Adds the contents of dst, src, and carry flag, and transfers the result to dst.

Details:

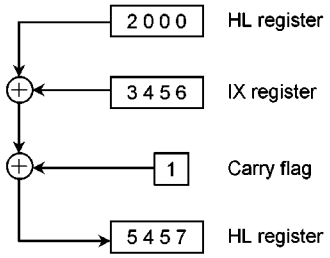
Size			Mnemonic		Code
Byte	Word	Long word			
O	O	O	ADC	R, r	1 1 z z 1 r 1 0 0 1 0 R
O	O	O	ADC	r, #	1 1 z z 1 r 1 1 0 0 1 0 0 1 #<7:0> #<15:8> #<23:16> #<31:24>
O	O	O	ADC	R, (mem)	1 m z z m m m m 1 0 0 1 0 R
O	O	O	ADC	(mem), R	1 m z z m m m m 1 0 0 1 1 R
O	O	×	ADC<W>	(mem), #	1 m 0 z m m m m 0 0 1 1 1 0 0 1 #<7:0> #<15:8>

Flags:

S	Z	H	V	N	C
*	*	*	*	0	*

S = MSB value of the result is set.
Z = 1 is set if the result is 0, otherwise 0.
H = 1 is set if a carry from bit 3 to bit 4 occurs as a result of the operation; otherwise, 0. If the operand is 32-bit, an undefined value is set.
V = 1 is set if an overflow occurs as a result of the operation; otherwise, 0.
N = Cleared to zero.
C = 1 is set if a carry occurs from the MSB, otherwise 0.

Execution example: ADC HL,IX
When the HL register = 2000H, the IX register = 3456H, and the carry flag = 1, execution sets the HL register to 5457H.



ADD dst, src

<Add>

Operation: dst ← dst + src

Description: Adds the contents of dst to those of src and transfers the result to dst.

Details:

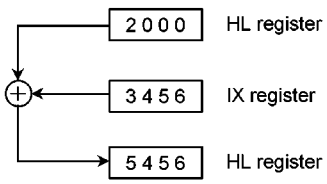
Size			Mnemonic		Code
Byte	Word	Long word			
O	O	O	ADD	R, r	1 1 z z 1 r 1 0 0 0 0 R
O	O	O	ADD	r, #	1 1 z z 1 r 1 1 0 0 1 0 0 0 #<7:0> #<15:8> #<23:16> #<31:24>
O	O	O	ADD	R, (mem)	1 m z z m m m m 1 0 0 0 0 R
O	O	O	ADD	(mem), R	1 m z z m m m m 1 0 0 0 1 R
O	O	×	ADD<W>	(mem), #	1 m 0 z m m m m 0 0 1 1 1 0 0 0 #<7:0> #<15:8>

Flags:

S	Z	H	V	N	C
*	*	*	*	0	*

S = MSB value of the result is set.
Z = 1 is set if the result is 0, otherwise 0.
H = 1 is set if a carry from bit 3 to bit 4 occurs as a result of the operation, otherwise 0. If the operand is 32-bit, an undefined value is set.
V = 1 is set if an overflow occurs as a result of the operation, otherwise 0.
N = Cleared to zero.
C = 1 is set if a carry occurs from the MSB, otherwise 0.

Execution example: ADD HL,IX
When the HL register = 2000H and the IX register = 3456H, execution sets the HL register to 5456H.



AND dst, src

<And>

Operation: dst ← dst AND src

Description: Ands the contents of dst and src, then transfers the result to dst.

(Truth table)

A	B	A and B
0	0	0
0	1	0
1	0	0
1	1	1

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	O	O	AND	R, r	1 1 z z 1 r 1 1 0 0 0 R
O	O	O	AND	r, #	1 1 z z 1 r 1 1 0 0 1 1 0 0 #<7:0> #<15:8> #<23:16> #<31:24>
O	O	O	AND	R, (mem)	1 m z z m m m m 1 1 0 0 0 R
O	O	O	AND	(mem), R	1 m z z m m m m 1 1 0 0 1 R
O	O	×	AND<W>	(mem), #	1 m 0 z m m m m 0 0 1 1 1 1 0 0 #<7:0> #<15:8>

Flags:

S	Z	H	V	N	C
*	*	1	*	0	0

S = MSB value of the result is set.
Z = 1 is set if the result is 0, otherwise 0.
H = 1 is set.
V = 1 is set if a parity of the result is even, 0 if odd. If the operand is 32 bits, an undefined value is set.
N = Cleared to zero.
C = Cleared to zero.

Execution example: AND HL,IX
When the HL register = 7350H and the IX register = 3456H, execution sets the HL register to 3050H.

```
        0111  0011  0101  0000 ← HL register (before execution)
AND)    0011  0100  0101  0110 ← IX register  (before execution)
        ---------------------
        0011  0000  0101  0000 ← HL register (after execution)
```

ANDCF num, src

<And Carry Flag>

Operation: CY ← CY and src<num>

Description: Ands the contents of the carry flag and bit num of src, and transfers the result to the carry flag.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	ANDCF	#4, r	1 1 0 z 1 r 0 0 1 0 0 0 0 0 0 0 0 0 # 4
○	○	×	ANDCF	A, r	1 1 0 z 1 r 0 0 1 0 1 0 0 0
○	×	×	ANDCF	#3, (mem)	1 m 1 1 m m m m 1 0 0 0 0 #3
○	×	×	ANDCF	A, (mem)	1 m 1 1 m m m m 0 0 1 0 1 0 0 0

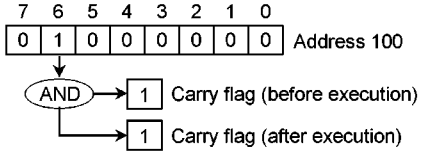
Notes: When bit num is specified by the A register, the value of the lower 4 bits of the A register is used as bit num. When the operand is a byte and the value of the lower 4 bits of bit num is from 8 to 15, the result is undefined.

Flags:

S	Z	H	V	N	C
–	–	–	–	–	*

S = No change
Z = No change
H = No change
V = No change
N = No change
C = The value obtained by anding the contents of the carry flag and the bit num of src is set.

Execution example: ANDCF 6,(100H)
When the contents of memory address 100 = 01000000B (binary) and the carry flag = 1, execution sets the carry flag to 1.



BIT num, src

<Bit test>

Operation: Z flag ← inverted value of src<num>

Description: Transfers the inverted value of the bit num of src to the Z flag.

Details:

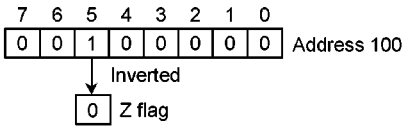
Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	BIT	#4, r	1 1 0 z 1 r 0 0 1 1 0 0 1 1 0 0 0 0 # 4
○	×	×	BIT	#3, (mem)	1 m 1 1 m m m m 1 1 0 0 1 #3

Flags:

S	Z	H	V	N	C
×	*	1	×	0	–

S = An undefined value is set.
Z = The inverted value of src <num> is set.
H = 1 is set.
V = An undefined value is set.
N = Reset to 0.
C = No change

Execution example: BIT 5, (100H)
When the contents of memory address 100 = 00100000B (binary), execution sets the Z flag to 0.



BS1B dst, src

<Bit Search 1 Backward>

Operation: dst ← src backward searched value

Description: Searches the src bit pattern backward (from MSB to LSB) for the first bit set to 1 and transfers the bit number to dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	×	BS1B	A, r	1 1 0 1 1 r 0 0 0 0 1 1 1 1

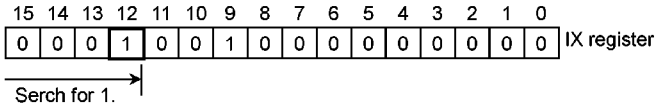
Note: dst in the operand must be the A register; src must be the register in words. If no bit set to 1 is found in the searched bit pattern, sets the A register to an undefined value and the V flag to 1.

Flags:

S	Z	H	V	N	C
–	–	–	*	–	–

S = No change
Z = No change
H = No change
V = 1 is set if the contents of src are all 0s (no bit is set to 1), otherwise 0.
N = No change
C = No change

Execution example: BS1B A,IX
When the IX register = 1200H, execution sets the A register to 0CH.



BS1F dst, src

<Bit Search 1 Forward>

Operation: dst ← src forward searched result

Description: Searches the src bit pattern forward (from LSB to MSB) for the first bit set to 1 and transfers the bit number to dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	×	BS1F	A, r	1 1 0 1 1 r 0 0 0 0 1 1 1 0

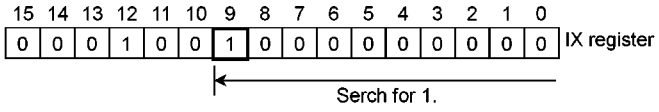
Note: dst in the operand must be the A register; src must be a register in words. If no bit set to 1 is found in the searched bit pattern, sets the A register to an undefined value and the V flag to 1.

Flags:

S	Z	H	V	N	C
–	–	–	*	–	–

S = No change
Z = No change
H = No change
V = 1 is set if the contents of src are all 0s (no bit is set to 1), otherwise 0.
N = No change
C = No change

Execution example: BS1F A,IX
When the IX register = 1200H, execution sets the A register to 09H.



CALL condition, dst

<Call subroutine>

Operation: In minimum mode, if cc is true, then XSP ← XSP – 2,(XSP) ← 16-bit PC, PC ← dst.
In maximum mode, if cc is true, then XSP ← XSP – 4,(XSP) ← 32-bit PC, PC ← dst.

Description: If the operand condition is true, saves the contents of the program counter to the stack area and jumps to the program address specified by dst.

Details:

Mnemonic	Code
CALL #16	0 0 0 1 1 1 0 0 #<7:0> #<15:8>
CALL #24	0 0 0 1 1 1 0 1 #<7:0> #<15:8> #<23:16>
CALL[cc,] mem	1 m 1 1 m m m m 1 1 1 0 c c

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: CALL 9000H
When the stack pointer XSP is 100H, executing this instruction at memory address 8000H writes the return address 8003H (long word data) to memory address 0FCH, sets the stack pointer XSP to 0FCH, and jumps to address 9000H.

CALR dst

<Call Relative>

Operation: In minimum mode, XSP ← XSP – 2,(XSP) ← 16-bit PC, PC ← dst.
In maximum mode, XSP ← XSP – 4,(XSP) ← 32-bit PC, PC ← dst.

Description: Saves the contents of the program counter to the stack area and makes a relative jump to the program address specified by dst.

Details:

Mnemonic	Code
CALR $ + 3 + d16	0 0 0 1 1 1 1 0 d<7:0> d<15:8>

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

CCF

<Complement Carry Flag>

Operation: CY ← inverted value of CY

Description: Inverts the contents of the carry flag.

Details:

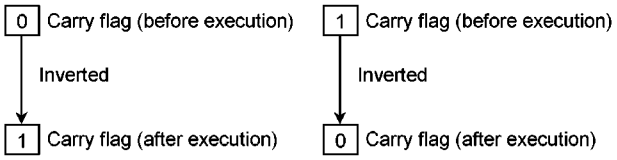
Mnemonic	Code
CCF	0 0 0 1 0 0 1 0

Flags:

S	Z	H	V	N	C
–	–	×	–	0	*

S = No change
Z = No change
H = An undefined value is set.
V = No change
N = Reset to 0.
C = Inverted value of itself is set.

Execution example: When the carry flag = 0, executing CCF sets the carry flag to 1; executing CCF again sets the carry flag to 0.



CHG num, dst

<Change>

Operation: dst<num> ← Inverted value of dst<num>

Description: Inverts the value of bit num of dst.

Details:

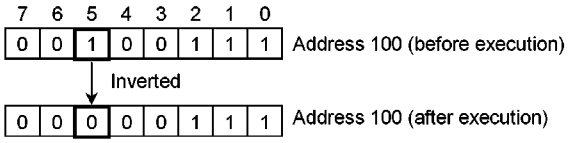
Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	CHG	#4, r	1 1 0 z 1 r 0 0 1 1 0 0 1 0 0 0 0 0 # 4
○	×	×	CHG	#3, (mem)	1 m 1 1 m m m m 1 1 0 0 0 #3

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: CHG 5, (100H)
When the contents of memory address 100 = 00100111B (binary), execution sets the contents to 00000111B (binary).



CP src1, src2

<Compare>

Operation: src1 – src2

Description: Compares the contents of src1 with those of src2 and indicates the results in flag register F.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	CP	R, r	1 1 z z 1 r 1 1 1 1 0 R
○	○	×	CP	r, #3	1 1 0 z 1 r 1 1 0 1 1 #3
○	○	○	CP	r, #	1 1 z z 1 r 1 1 0 0 1 1 1 1 #<7:0> #<15:8> #<23:16> #<31:24>
○	○	○	CP	R, (mem)	1 m z z m m m m 1 1 1 1 0 R
○	○	○	CP	(mem), R	1 m z z m m m m 1 1 1 1 1 R
○	○	×	CP <W>	(mem), #	1 m 0 z m m m m 0 0 1 1 1 1 1 1 #<7:0> #<15:8>

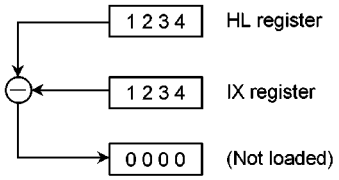
Note: #3 in operands indicates from 0 to 7

Flags:

S	Z	H	V	N	C
*	*	*	*	1	*

S = MSB value of the result is set.
Z = 1 is set if the result is 0, otherwise 0.
H = 1 is set if a borrow from bit 3 to bit 4 occurs as a result of the operation, otherwise 0.
If the operand is 32 bits, an undefined value is set.
V = 1 is set if an overflow occurs as a result of the operation, otherwise 0.
N = 1 is set.
C = 1 is set if a borrow occurs from the MSB bit as a result of the operation, otherwise 0.

Execution example: CP HL,IX
When the HL register = 1234H and the IX register = 1234H, execution sets the Z and N flags to 1 and clears the S, H, V, and C flags to zero.



CPD src1, src2

<Compare Decrement>

Operation: src1 – src2, BC ← BC – 1

Description: Compares the contents of src1 with those of src2, then decrements the contents of the BC register by 1. src1 must be the A or WA register. src2 must be in post–decrement register indirect addressing mode.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	O	×	CPD	[A/WA, (R–)]	1 0 0 z 0 R 0 0 0 1 0 1 1 0

Note: Omitting operands in square brackets [ ] specifies A,(XHL–).

Flags:

S	Z	H	V	N	C
*	*	*	*	1	–

S = MSB value of the result of src1 – src2 is set.
Z = 1 is set if the result of src1 – src2 is 0, otherwise 0.
H = 1 is set if a borrow from bit 3 to bit 4 occurs as a result of src1 – src2, otherwise 0.
V = 0 is set if the BC register value is 0 after execution, otherwise 1.
N = 1 is set.
C = No change

Execution example: CPDA, (XIX–)
When the XIX register = 00123456H and the BC register = 0200H, execution compares the contents of the A register with those of memory address 123456H, then sets the XIX register to 00123455H, the BC register to 01FFH.

CPDR src1, src2

<Compare Decrement Repeat>

Operation: src1 – src2, BC ← BC – 1, Repeat until src1 = src2 or BC = 0

Description: Compares the contents of src1 with those of src2. Then decrements the contents of the BC register by 1. Repeats until src1 = src2 or BC = 0. src1 must be the A or WA register. src2 must be in post-decrement register indirect addressing mode.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	CPDR	[A/WA, (R–)]	1 0 0 z 0 R 0 0 0 1 0 1 1 1

Note: Omitting operands in square brackets [ ] specifies A,(XHL–).

Flags:

S	Z	H	V	N	C
*	*	*	*	1	–

S = MSB value of the result of src1 – src2 is set.
Z = 1 is set if the result of src1 – src2 is 0, otherwise 0.
H = 1 is set if a borrow from bit 3 to bit 4 occurs as a result of src1 – src2, otherwise 0.
V = 0 is set if the BC register value is 0 after execution, otherwise 1.
N = 1 is set.
C = No change

Execution example: CPDR A,(XIX–)

Under the following conditions, execution reads the contents of memory addresses 123456H, 123455H, and 123454H. The instruction ends with condition BC = 0 and sets the XIX register to 00123453H and the BC register to 0000H.

Conditions: A register = 55H
XIX register = 00123456H
BC register = 0003H
Memory address 123456H = 11H
Memory address 123455H = 22H
Memory address 123454H = 33H

CPI src1, src2

<Compare Increment>

Operation: src1 – src2, BC ← BC – 1

Description: Compares the contents of src1 with those of src2, then decrements the contents of the BC register by 1. src1 must be the A or WA register. src2 must be in post-increment register indirect addressing mode.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	O	×	CPI	[A/WA, (R+)]	0 0 0 z 0 R 0 0 0 1 0 1 0 0

Note: Omitting operands enclosed in square brackets [ ] specifies A,(XHL+).

Flags:

S	Z	H	V	N	C
*	*	*	*	1	–

S = MSB value of the result of src1 – src2 is set.
Z = 1 is set if the result of src1 – src2 is 0, otherwise 0.
H = 1 is set if a borrow from bit 3 to bit 4 occurs as a result of src1 – src2, otherwise 0.
V = 0 is set if the BC register value is 0 after execution, otherwise 1.
N = 1 is set.
C = No change

Execution example: CPI A, (XIX+)
When the XIX register = 00123456H and the BC register = 0200H, execution compares the contents of the A register with those of memory address 123456H, and sets the XIX register to 00123457H and the BC register to 01FFH.

CPIR src1, src2

<Compare Increment Repeat>

Operation: src1 – src2, BC ← BC – 1, repeat until src1 = src2 or BC = 0

Description: Compares the contents of src1 with those of src2. Then decrements the contents of the BC register by 1. Repeats until src1 = src2 or BC = 0. src1 must be the A or WA register. src2 must be in post-increment register indirect addressing mode.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	CPIR	[A/WA, (R+)]	1 0 0 z 0 R 0 0 0 1 0 1 0 1

Note: Omitting operands in square brackets [ ] specifies A,(XHL+).

Flags:

S	Z	H	V	N	C
*	*	*	*	1	–

S = MSB value of the result of src1 – src2 is set.
Z = 1 is set if the result of src1 – src2 is 0, otherwise 0.
H = 1 is set if a borrow from bit 3 to bit 4 occurs as a result of src1 – src2, otherwise 0.
V = 0 is set if the BC register value is 0 after execution, otherwise 1.
N = 1 is set.
C = No change

Execution example: CPIR A, (XIX+)

Under the following conditions, execution reads memory addresses 123456H, 123457H, and 123458H. The instruction ends with condition src1 = src2, sets the XIX register to 00123459H and the BC register to 01FDH.

Conditions: A register = 33H
XIX register = 00123456
HBC register = 0200H
Memory address 123456H = 11H
Memory address 123457H = 22H
Memory address 123458H = 33H

CPL dst

<Complement>

Operation: dst ← Ones complement of dst

Description: Transfers the value of ones complement (inverted bit of 0/1) of dst to dst.

Details:

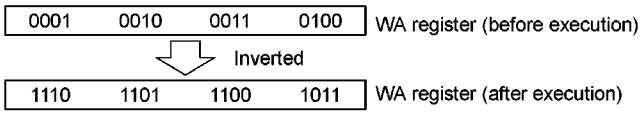
Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	CPL	r	1 1 0 z 1 r 0 0 0 0 0 1 1 0

Flags:

S	Z	H	V	N	C
–	–	1	–	1	–

S = No change
Z = No change
H = 1 is set.
V = No change
N = 1 is set.
C = No change

Execution example: CPL WA
When the WA register = 1234H, execution sets the WA register to EDCBH.



DAA dst

<Decimal Adjust Accumulator>

Operation: dst ← decimal adjustment of dst

Description: Decimal adjusts the contents of dst depending on the states of the C, H, and N flags. Used to adjust the execution result of the add or subtract instruction as binary-coded decimal (BCD).

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	×	×	DDA	r	1 1 0 0 1 r 0 0 0 1 0 0 0 0

Operation	N flag before DAA instruction execution	C flag before DAA instruction execution	Upper 4 bits of dst	H flag before DAA instruction execution	Lower 4 bits of dst	Added value	C flag after DAA instruction execution
	0	0	0 to 9	0	0 to 9	00	0
	0	0	0 to 8	0	A to F	06	0
ADD	0	0	0 to 9	1	0 to 3	06	0
	0	0	A to F	0	0 to 9	60	1
ADC	0	0	9 to F	0	A to F	66	1
	0	0	A to F	1	0 to 3	66	1
	0	1	0 to 2	0	0 to 9	60	1
	0	1	0 to 2	0	A to F	66	1
	0	1	0 to 3	1	0 to 3	66	1
SUB	1	0	0 to 9	0	0 to 9	00	0
SBC	1	0	0 to 8	1	6 to F	FA	0
NEG	1	1	7 to F	0	0 to 9	A0	1
	1	1	6 to F	1	6 to F	9A	1

Note: Decimal adjustment cannot be performed for the INC or DEC instruction. This is because the C flag does not change.

Flags:

S	Z	H	V	N	C
*	*	*	*	–	*

S = MSB value of the result is set.
Z = 1 is set if the result is 0, otherwise 0.
H = 1 is set if a carry from bit 3 to bit 4 occurs as a result of the operation, otherwise 0.
V = 1 is set if the parity (number of 1s) of the result is even, otherwise 0.
N = No change
C = 1 is set if a carry occurs from the MSB as a result of the operation or a carry was 1 before operation, otherwise 0.

Execution example:
```
ADD A,B
DAA A
```
When the A register = 59H and the B register = 13H, execution sets the A register to 72H.

DEC num, dst

<Decrement>

Operation: dst ← dst – num

Description: Decrements dst by the contents of num and transfers the result to dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	DEC	#3, r	1 1 z z 1 r 0 1 1 0 1 #3
○	○	×	DEC<W>	#3, (mem)	1 m 0 z m m m m 0 1 1 0 1 #3

Note: #3 in operands indicates from 1 to 8; object codes correspond from 1 to 7,0.

Flags:

S	Z	H	V	N	C
*	*	*	*	1	–

S = MSB value of the result is set.
Z = 1 is set if the result is 0, otherwise 0.
H = 1 is set if a borrow from bit 3 to bit 4 occurs as a result of the operation, otherwise 0.
V = 1 is set if an overflow occurs as a result of the operation, otherwise 0.
N = 1 is set.
C = No change

Note: With the DEC #3, r instruction, if the operand is a word or a long word, no flags change.

Execution example: DEC 4, HL
When the HL register = 5678H, execution sets the HL register to 5674H.

DECF

<Decrement Register File Pointer>

Operation: RFP<2:0> ← RFP<2:0> – 1

Description: Decrements the contents of register file pointer RFP <2:0> in the status register by 1. RFP2 is fixed to 0.

Details:

Mnemonic	Code
DECF	0 0 0 0 1 1 0 1

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: DECF
When the contents of RFP<2:0> = 2, execution sets the contents of RFP<2:0> to 1.

DI

<Disable Interrupt>

Operation: IFF<2:0> ← 7

Description: Sets the contents of the interrupt enable flag (IFF) <2:0> in status register to 7. After execution, only non-maskable interrupts (interrupt level 7) can be received.

Details:

Mnemonic	Code
DI	0 0 0 0 0 1 1 0 0 0 0 0 0 1 1 1

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

DIV dst, src

<Divide>

Operation: dst<lower half> ← dst ÷ src,dst<upper half> ← remainder (unsigned)

Description: Divides unsigned the contents of dst by those of src and transfers the quotient to the lower half of dst, the remainder to the upper half of dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	DIV	RR, r	1 1 0 z 1 r 0 1 0 1 0 R
○	○	×	DIV	rr, #	1 1 0 z 1 r 0 0 0 0 1 0 1 0 #<7:0> #<15:8>
○	○	×	DIV	RR, (mem)	1 m 0 z m m m m 0 1 0 1 0 R

*For RR, see the following page.

Notes: When the operation is in bytes, dst (lower byte) ← dst (word) ÷ src (byte), dst (upper byte) ← remainder.
When the operation is in words, dst (lower word) ← dst (long word) ÷ src (word), dst (upper word) ← remainder. Match coding of the operand dst with the size of the dividend.

Flags:

S	Z	H	V	N	C
–	–	–	V	–	–

S = No change
Z = No change
H = No change
V = 1 is set when divided by 0 or the quotient exceeds the numerals which can be expressed in bits of dst for load; otherwise, 0 is set.
N = No change
C = No change

Execution example: DIV XIX,IY

When the XIX register = 12345678H and the IY register = 89ABH, execution results in a quotient of 21DAH and a remainder of 0FDAH, and sets the XIX register to 0FDA21DAH.

Note: "RR" of the DIV RR,r and DIV RR,(mem) instruction is as listed below.

Operation size in bytes
(8 bits ← 16 bits ÷ 8 bits)

RR	Code "R"
WA	001
BC	011
DE	101
HL	111
IX	Specification not possible!
IY	
IZ	
SP	

Operation size in words
(16 bits ← 32 bits ÷ 16 bits)

RR	Code "R"
XWA	000
XBC	001
XDE	010
XHL	011
XIX	100
XIY	101
XIZ	110
XSP	111

*1 When the CPU is in minimum mode, XWA, XBC, XDE, and XHL cannot be used.

"rr" of the DIV rr,# instruction is as listed below.

Operation size in bytes
(8 bits ← 16 bits ÷ 8 bits)

rr	Code "r"
WA	001
BC	011
DE	101
HL	111
IX	C7H : F0H
IY	C7H : F4H
IZ	C7H : F8H
SP	C7H : FCH
	1st byte 2nd byte

*2 Any other word registers can be specified in the same extension coding as IX to SP.

Operation size in words
(16 bits ← 32 bits ÷ 16 bits)

rr	Code "r"
XWA	000
XBC	001
XDE	010
XHL	011
XIX	100
XIY	101
XIZ	110
XSP	111

*3 When the CPU is in minimum mode, XWA, XBC, XDE, and XHL cannot be used.

*4 Any other long word registers can be specified in the extension coding.

DIVS dst, src

<Divide Signed>

Operation: dst<lower half> ← dst ÷ src,dst<upper half> ← remainder (signed)

Description: Divides signed the contents of dst by those of src and transfers the quotient to the lower half of dst, the remainder to the upper half of dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	O	×	DIVS	RR, r	1 1 0 z 1 r 0 1 0 1 1 R
O	O	×	DIVS	rr, #	1 1 0 z 1 r 0 0 0 0 1 0 1 1 #<7:0> #<15:8>
O	O	×	DIVS	RR, (mem)	1 m 0 z m m m m 0 1 0 1 1 R

*For RR, see the following page.

Notes: When the operation is in bytes, dst (lower byte) ← dst (word) ÷ src (byte), dst (upper byte) ← remainder.
When the operation is in words,dst (lower word) ← dst (long word) ÷ src (word), dst (upper word) ← remainder.
Match coding of the operand dst with the size of the dividend. The sign of the remainder is the same as that of the dividend.

Flags:

S	Z	H	V	N	C
–	–	–	*	–	–

S = No change
Z = No change
H = No change
V = 1 is set when divided by 0, or the quotient exceeds the value which can be expressed in bits of the dst used for loading, otherwise 0.
N = No change
C = No change

Execution example: DIVS XIX,IY
When the XIX register = 12345678H and the IY register = 89ABH, execution results in the quotient as 16EEH and the remainder as D89EH, and sets the XIX register to 16EED89EH.

Note: "RR" of the DIVS RR,r and DIVS RR, (mem) instruction is as listed below.

Operation size in bytes
(8 bits ← 16 bits ÷ 8 bits)

RR	Code "R"
WA	001
BC	011
DE	101
HL	111
IX	Specification not possible!
IY	
IZ	
SP	

Operation size in words
(16 bits ← 32 bits ÷ 16 bits)

RR	Code "R"
XWA	000
XBC	001
XDE	010
XHL	011
XIX	100
XIY	101
XIZ	110
XSP	111

*1 When the CPU is in minimum mode, XWA, XBC, XDE, or XHL cannot be used.

"rr" of the DIVS rr,# instruction is as listed below.

Operation size in bytes
(8 bits ← 16 bits ÷ 8 bits)

rr	Code "r"
WA	001
BC	011
DE	101
HL	111
IX	C7H : F0H
IY	C7H : F4H
IZ	C7H : F8H
SP	C7H : ECH
	1st byte 2nd byte

*2 Any other word registers can be specified in the same extension coding as those for IX to SP.

Operation size in words
(16 bits ← 32 bits ÷ 16 bits)

rr	Code "r"
XWA	000
XBC	001
XDE	010
XHL	011
XIX	100
XIY	101
XIZ	110
XSP	111

*3 When the CPU is in minimum mode, XWA, XBC, XDE, or XHL cannot be used.

*4 Any other long word registers can be specified in the extension coding.

DJNZ dst1, dst2

<Decrement and Jump if Non Zero>

Operation: dst1 ← dst1 – 1. if dst1 ≠ 0, then PC ← dst2.

Description: Decrements the contents of dst1 by 1. Makes a relative jump to the program address specified by dst2 if the result is other than 0.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	DJNZ	[r,]$ + 3/4 + d8	1 1 0 z 1 r 0 0 0 1 1 1 0 0 d<7:0>

(Note) $ + 4 + d8 ("r" is specified using extension codes.)
$ + 3 + d8 (otherwise)

Note: Omitting "r" of the operand in square brackets [ ] is regarded as specifying the B register.

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example:
```
LOOP: ADD  A, A
      DJNZ W, LOOP
```
When the A register = 12H and the W register = 03H, execution loops three times and sets the A register to 24H → 48 → 90H and the W register to 02H → 01H → 00H.

EI num

<Enable Interrupt>

Operation: IFF <2:0> ← num

Description: Sets the contents of the IFF<2:0> in the status register to num. After execution, the CPU interrupt receive level becomes num.

Details:

Mnemonic		Code
EI	[#3]	0 0 0 0 0 1 1 0 0 0 0 0 0 #3

Note: A value from 0 to 7 can be specified as the operand value. If the operand is omitted, the default value is "0" (EI 0).

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

EX dst, src

<Exchange>

Operation: dst ↔ src

Description: Exchanges the contents of dst and src.

Details:

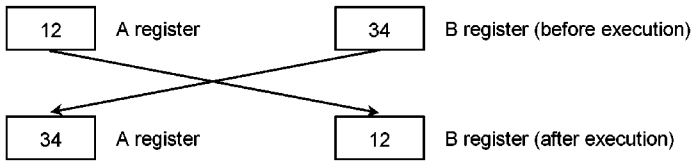
Size			Mnemonic		Code
Byte	Word	Long word			
O	×	×	EX	F, F'	0 0 0 1 0 1 1 0
O	O	×	EX	R, r	1 1 z z 1 r 1 0 1 1 1 R
O	O	×	EX	(mem), r	1 m z z m m m m 0 0 1 1 0 R

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change
* Executing EX F,F' changes all flags.

Execution example: EX A,B
When the A register = 12H and the B register = 34H, execution sets the A register to 34H and the B register to 12H.



EXTS dst

<Extend Sign>

Operation: dst <upper half> ← signed bit of dst <lower half>

Description: Transfers (copies) the signed bit (bit 7 when the operand size is a word, bit 15 when a long word) of the lower half of dst to all bits of the upper half of dst.

Details:

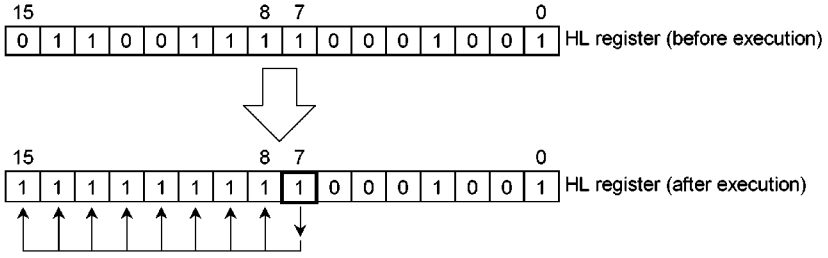
Size			Mnemonic		Code
Byte	Word	Long word			
×	O	O	EXTS	r	1 1 z z 1 r / 0 0 0 1 0 0 1 1

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: EXTS HL
When the HL register = 6789H, execution sets the HL register to FF89H.



EXTZ dst

<Extend Zero>

Operation: dst<upper half> ← 0

Description: Clears the upper half of dst to zero. Used for making the operand sizes the same when they are different.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	○	○	EXTZ	r	1 1 z z 1 r / 0 0 0 1 0 0 1 0

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: EXTZ HL
When the HL register = 6789H, execution sets the HL register to 0089H.

EXTZ XIX
When the XIX register = 12345678H, execution sets the XIX register to 00005678H.

HALT

<Halt CPU>

Operation: CPU halt

Description: Halts the instruction execution. To resume, an interrupt must de received.

Details:

Mnemonic	Code
HALT	0 0 0 0 0 1 0 1

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

INC num, dst

<Increment>

Operation: dst ← dst + num

Description: Adds the contents of dst and num and transfers the result to dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	INC	#3, r	1 1 z z 1 r 0 1 1 0 0 #3
○	○	×	INC<W>	#3, (mem)	1 m 0 z m m m m 0 1 1 0 0 #3

Note: #3 in operands indicates from 1 to 8 and object codes correspond from 1 to 7,0.

Flags:

S	Z	H	V	N	C
*	*	*	*	0	–

S = MSB value of the result is set.
Z = 1 is set if the result is 0, otherwise 0.
H = 1 is set if a carry occurs from bit 3 to bit 4 as a result of the operation, otherwise 0.
V = 1 is set if an overflow occurs as a result of the operation, otherwise 0.
N = Cleared to zero.
C = No change

Note: With the INC #3,r instruction, if the operand is a word or a long word, no flags change.

Execution example: INC 5,WA
When the WA register = 1234H, execution sets the WA register to 1239H.

INCF

<Increment Register File Pointer>

Operation: RFP<2:0> ← RFP<2:0> + 1

Description: Increments the contents of RFP<2:0> in the status register by 1. RFP2 is fixed to 0.

Details:

Mnemonic	Code
INCF	0 0 0 0 1 1 0 0

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: INCF
When the contents of RFP<2:0> = 2, execution sets the contents of RFP<2:0> to 3.

JP condition, dst

<Jump>

Operation: If cc is true, then PC ← dst.

Description: If the operand condition is true, jumps to the program address specified by dst.

Details:

Size	Mnemonic	Code
	JP #16	0 0 0 1 1 0 1 0 #<7:0> #<15:8>
	JP #24	0 0 0 1 1 0 1 1 #<7:0> #<15:8> #<23:16>
	JP [cc,] mem	1 m 1 1 m m m m 1 1 0 1 c c

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: JP 2000H
Execution jumps unconditionally to address 2000H.

JR condition, dst

<Jump Relative>

Operation: If cc is true, then PC ← dst.

Description: If the operand condition is true, makes a relative jump to the program address specified by dst.

Details:

Mnemonic		Code
JR	[cc,] $ + 2 + d8	0 1 1 0 c c / d<7:0>
JRL	[cc,] $ + 3 + d16	0 1 1 1 c c / #<7:0> / #<15:8>

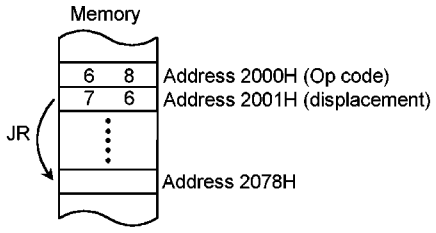
Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: JR 2078H

When this instruction is executed at memory address 2000H, execution relative jumps unconditionally to address 2078H. The object code of the instruction is 68H : 76H.



LD dst, src

<Load>

Operation: dst ← src

Description: Loads the contents of src to dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	LD	R, r	1 1 z z 1 r 1 0 0 0 1 R
○	○	○	LD	r, R	1 1 z z 1 r 1 0 0 1 1 R
○	○	○	LD	r, #3	1 1 z z 1 r 1 0 1 0 1 #3
○	○	○	LD	R, #	0 z z z 0 R #<7:0> #<15:8> #<23:16> #<31:24>
○	○	○	LD	r, #	1 1 z z 1 r 0 0 0 0 0 0 1 1 #<7:0> #<15:8> #<23:16> #<31:24>
○	○	○	LD	R, (mem)	1 m z z m m m m 0 0 1 0 0 R
○	○	○	LD	(mem), R	1 m 1 1 m m m m 0 1 z z 0 R
○	○	×	LD<W>	(#8), #	0 0 0 0 1 0 z 0 #8 #<7:0> #<15:8>

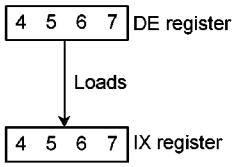
Size			Mnemonic		Code
Byte	Word	Long word			
O	O	×	LD<W>	(mem), #	1 m 1 1 m m m m 0 0 0 0 0 0 z 0 #<7:0> #<15:8>
O	O	×	LD<W>	(#16), (mem)	1 m 0 z m m m m 0 0 0 1 1 0 0 1 #16<7:0> #16<15:8>
O	O	×	LD<W>	(mem), (#16)	1 m 1 1 m m m m 0 0 0 1 0 1 z 0 #16<7:0> #16<15:8>

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: LD IX, DE
When the DE register = 4567H, execution sets the IX register to 4567H.



LDA dst, src

<Load Address>

Operation: dst ← src effective address value

Description: Loads the src effective address value to dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	O	LDA	R, mem	1 m 1 1 m m m m / 0 0 1 s 0 R

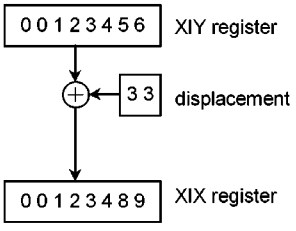
Note: This instruction operates much like the ADD instruction; the difference is that dst is specified independently from src. Mainly used for handling the pointer with the C compiler.

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: LDA XIX, XIY + 33H
When the XIY register = 00123456H, execution sets the XIX register to 00123489H.



LDAR dst, src

<Load Address Relative>

Operation: dst ← src relative address value

Description: Loads the relative address value specified in src to dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	O	LDAR	R, $ + 4 + d16	1 1 1 1 0 0 1 1 0 0 0 1 0 0 1 1 d<7:0> d<15:8> 0 0 1 s 0 R

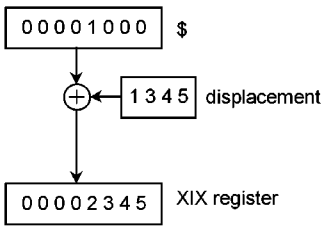
Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: LDAR XIX, $ + 1345H

When this instruction is executed at memory address 1000H, execution sets the XIX register to 00002345H. $ indicates the start address of the instruction. The instruction's object codes are: F3H:13H:41H:13H:34H.



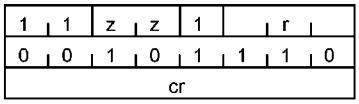
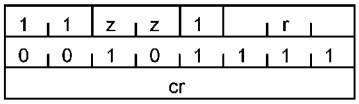
LDC dst, src

<Load Control Register>

Operation: dst ← src

Description: Loads the contents of src to dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	O	O	LDC	cr, r	
O	O	O	LDC	r, cr	

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: LDC DMAC0, WA
When the WA register = 1234H, execution sets control register DMAC0 to 1234H.

LDCF num, src

<Load Carry Flag>

Operation: CY ← src<num>

Description: Loads the contents of bit num of src to the carry flag.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	O	×	LDCF	#4, r	1 1 0 z 1 r 0 0 1 0 0 0 1 1 0 0 0 0 # 4
O	O	×	LDCF	A, r	1 1 0 z 1 r 0 0 1 0 1 0 1 1
O	×	×	LDCF	#3, (mem)	1 m 1 1 m m m m 1 0 0 1 1 #3
O	×	×	LDCF	A, (mem)	1 m 1 1 m m m m 0 0 1 0 1 0 1 1

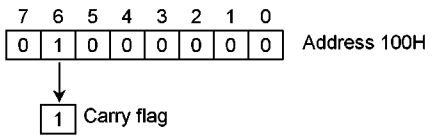
Notes: When bit num is specified by the A register, the value of the lower 4 bits of the A register is used as bit num. When the operand is a byte and the value of the lower 4 bits of bit num is from 8 to 15, the value of the carry flag is undefined.

Flags:

S	Z	H	V	N	C
–	–	–	–	–	*

S = No change
Z = No change
H = No change
V = No change
N = No change
C = Contents of bit num of src is set.

Execution example: LDCF 6, (100H)
When the contents of memoryad address 100 = 01000000B (binary), execution sets the carry flag to 1.



LDD dst, src

<Load Decrement>

Operation: dst ← src, BC ← BC – 1

Description: Loads the contents of src to dst, then decrements the contents of the BC register by 1. src and dst must be in post-decrement register indirect addressing mode.

Details:

Size			Mnemonic	Code
Byte	Word	Long word		
○	○	×	LDD<W>[(XDE–), (XHL–)]	1 0 0 z 0 0 1 1 0 0 0 1 0 0 1 0
			LDD<W> (XIX–), (XIY–)	1 0 0 z 0 1 0 1 0 0 0 1 0 0 1 0

*Coding in square brackets [ ] can be omitted.

Flags:

S	Z	H	V	N	C
–	–	0	*	0	–

S = No change
Z = No change
H = Cleared to 0.
V = 0 is set if the BC register value is 0 after execution, otherwise 1.
N = Cleared to zero.
C = No change

Execution example: LDD (XIX–), (XIY–)
When the XIX register = 00123456H, the XIY register = 00335577H, and the BC register = 0700H, execution loads the contents at address 335577 to address 123456H and sets the XIX register to 123455H, the XIY register to 00335576H, and the BC register to 06FFH.

LDDR dst, src

<Load Decrement Repeat>

Operation: dst ← src, BC ← BC – 1, Repeat until BC = 0

Description: Loads the contents of src to dst, then decrements the contents of the BC register by 1. If the result is other than 0, the operation is repeated. src and dst must be in post-decrement register indirect addressing mode.

Details:

Size			Mnemonic	Code
Byte	Word	Long word		
O	O	×	LDDR<W>[(XDE–), (XHL–)]	1 0 0 z 0 0 1 1 0 0 0 1 0 0 1 1
O	O	×	LDDR<W> (XIX–), (XIY–)	1 0 0 z 0 1 0 1 0 0 0 1 0 0 1 1

* Coding in square brackets [ ] can be omitted.

Flags:

S	Z	H	V	N	C
–	–	0	0	0	–

S = No change
Z = No change
H = Cleared to zero.
V = Cleared to zero.
N = Cleared to zero.
C = No change

Execution example: LDDR (XIX–), (XIY–)
When the XIX register = 00123456H, the XIY register = 00335577H, and the BC register = 0003H, the results of the execution are as follows:

Loads the contents of address 335577H to 123456H.
Loads the contents of address 335576H to 123455H.
Loads the contents of address 335575H to 123454H.
Sets the XIX register to 00123453H.
Sets the XIY register to 00335574H.
Sets the BC register to 0000H.

LDF num

<Load Register File Pointer>

Operation: RFP<2:0> ← num

Description: Loads the num value to the register file pointer RFP<2:0> in status register. Since RFP2 is fixed to 0 in maximam mode, when the num value is from 4 to 7, REP is set to from 0 to 3.

Details:

Mnemonic		Code
LDF	#3	0 0 0 1 0 1 1 1 0 0 0 0 0 #3

Note: In minimum mode, the operand value can be specified from 0 to 7; in maximum mode, from 0 to 3.

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

LDI dst, src

<Load Increment>

Operation: dst ← src, BC ← BC – 1

Description: Loads the contents of src to dst, then decrements the contents of the BC register by 1. src and dst must be in post-increment register indirect addressing mode.

Details:

Size			Mnemonic	Code
Byte	Word	Long word		
O	O	×	LDI<W>[(XDE+), (XHL+)]	1 0 0 z 0 0 1 1 0 0 0 1 0 0 0 0
O	O	×	LDI<W> (XIX+), (XIY+)	1 0 0 z 0 1 0 1 0 0 0 1 0 0 0 0

* Coding in square brackets [ ] can be omitted.

Flags:

S	Z	H	V	N	C
–	–	0	*	0	–

S = No change
Z = No change
H = Cleared to zero.
V = 0 is set when the BC register value is 0 after execution, otherwise 1.
N = Cleared to zero.
C = No change

Execution example: LDI (XIX+), (XIY+)
When the XIX register = 00123456H, the XIY register = 00335577H, and the BC register = 0700H, execution loads the contents of address 335577H to 123456H and sets the XIX register to 00123457H, the XIY register to 00335578H, and the BC register to 06FFH.

LDIR dst, src

<Load Increment Repeat>

Operation: dst ← src, BC ← BC – 1, Repeat until BC = 0

Description: Loads the contents of src to dst, then decrements the contents of the BC register by 1. If the result is other than 0, the operation is repeated. src and dst must be in post-increment register indirect addressing mode.

Details:

Size			Mnemonic	Code
Byte	Word	Long word		
O	O	×	LDIR<W>[(XDE+), (XHL+)]	1 0 0 z 0 0 1 1 0 0 0 1 0 0 0 1
O	O	×	LDIR<W> (XIX+), (XIY+)	1 0 0 z 0 1 0 1 0 0 0 1 0 0 0 1

* Coding in square brackets [ ] can be omitted.

Note: Interrupt requests are sampled every time 1 item of data is loaded.

Flags:

S	Z	H	V	N	C
–	–	0	0	0	–

S = No change
Z = No change
H = Cleared to zero.
V = Cleared to zero.
N = Cleared to zero.
C = No change

Execution example: LDIR (XIX+), (XIY+)
When the XIX register = 00123456H, the XIY register = 00335577H, and the BC register = 0003H, execution results as follows:

Loads the contents of address 335577H to 123456H.
Loads the contents of address 335578H to 123457H.
Loads the contents of address 335579H to 123458H.
Sets the XIX register to 00123459H.
Sets the XIY register to 0033557AH.
Sets the BC register to 0000H.

LDX dst, src

<Load eXtract>

Operation: dst ← src

Description: Loads the contents of src to dst. The effective code is assigned to this instruction every other byte. Used to fetch the code from 8-bit data bus memory in 16-bit data bus mode.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	×	×	LDX	(#8), #	1 1 1 1 0 1 1 1 0 0 0 0 0 0 0 0 #8 0 0 0 0 0 0 0 0 # 0 0 0 0 0 0 0 0

Note: Even if the second, fourth, or sixth instruction code value is not 00H, the instruction operates correctly.

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

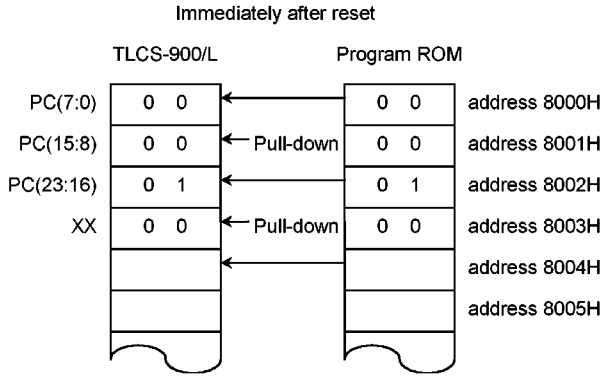
When the CPU fetches the program after reset, the bus width which is specified for 900/L is 16-bit and the bus width of the external program ROM is 8-bit, this instruction can be used.

The following table describes the conditions for using this instruction.

AM8/16 pin	Bus width of Program ROM	Bus width of the memories	LDX instruction
1	8-bit	8-bit	Not available
0	16-bit	8/16-bit	Not available
0	8-bit	8/16-bit	Available

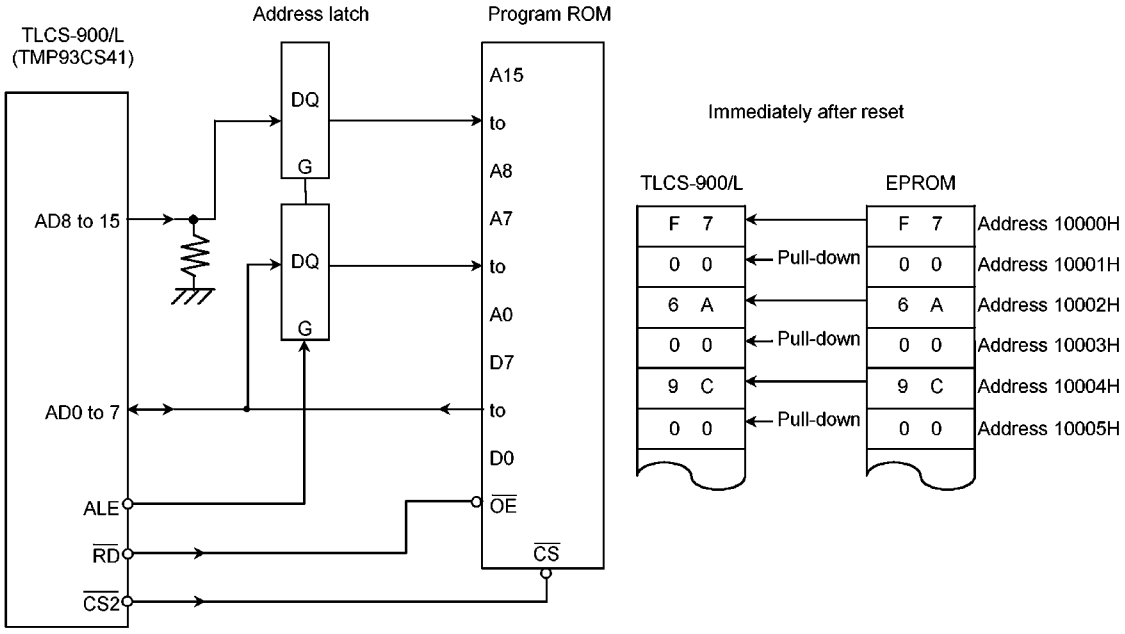
Execution example: Using the TMP93CS41, the example executes the program using a program ROM which has an 8-bit data bus when the AM8/$\overline{16}$ pin is "0" and the memories other than the program ROM are 16-bit. After reset, the reset vector is read in 16-bit data bus mode. When the program starts with the external memory which has an 8-bit data bus, it is necessary to connect the pull-up/pull-down resistor with the data bus AD8 to 15 pins of the upper side and input the PC (15:8) of the reset vector.

For example, when the reset vector is set to address 010000H, 010000H is placed on address 8000H of the program ROM and AD8 to 15 pins are pulled down. The data 00H can be input as a value of the PC (15:8). The LDX instruction is placed on address 010000H of the program ROM.



LDX (6AH), 9CH

Executing the above instruction writes of 9CH to the control register at address 6AH of the built-in programmable chip select/wait controller. As a result, memory addresses 8000H to 3FFFFFH are set to 8-bit data bus 0WAIT mode. The program is fetched and executed with the 8-bit bus width from the next instruction.



Note: The pull-up/down resistors which are connected with AD8 to 15 pins to input the reset vector PC (15:8) conflict with address and data output form AD8 to 15 pins. This causes increasing of the consumption current.
To prevent the consumption current increasing, the pull-up/down resistors should be input off after the above procedures are implemented.

LINK dst, num

<Link>

Operation: (–XSP) ← dst, dst ← XSP, XSP ← XSP + num

Description: Saves the contents of dst to the stack area. Loads the contents of stack pointer XSP to dst. Adds the contents of XSP to those of num (signed) and loads the result to XSP. Used for obtaining a local variable area in the stack area for -num bytes.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	×	O	LINK	r,d16	1 1 1 0 1 r 0 0 0 0 1 1 0 0 d<7:0> d<15:8>

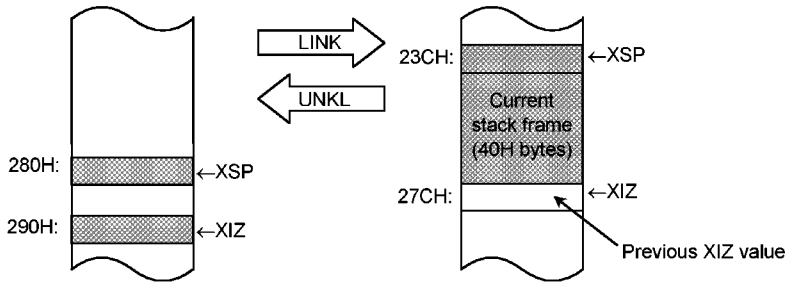
Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: LINK XIZ, –40H

When stack pointer XSP = 280H and the XIZ register = 290H, execution writes 00000290H (long data) at memory address 27CH and sets the XIZ register to 27CH and the stack pointer to XSP 23CH.



MDEC1 num, dst

<Modulo Decrement 1>

Operation: if (dst mod num) = 0 then dst ← dst + (num – 1) else dst ← dst – 1.

Description: When the modulo num of dst is 0, increments dst by num –1 .
Otherwise, decrements dst by 1. Used to operate pointers for cyclic memory table.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	×	MDEC1	#, r	1 1 0 1 1 r 0 0 1 1 1 1 0 0 #<7:0> – 1 #<15:8>

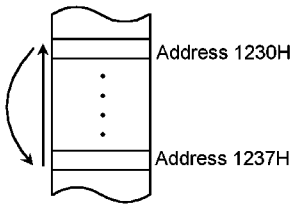
Note: The operand # must be 2 to the nth power. (n = 1 to 15)

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: Decrements the IX register by cycling from 1230H to 1237H.
MDEC1 8, IX
When the IX register = 1231H, execution sets the IX register to 1230H. Further execution increments the IX register by 8 – 1 and sets the IX register to 1237H, since the IX register modulo 8 = 0.



MDEC2 num, dst

<Modulo Decrement 2>

Operation: if (dst mod num) = 0 then dst ← dst + (num – 2) else dst ← dst – 2.

Description: When the modulo num of dst is 0, increments dst by num –2.
Otherwise, decrements dst by 2. Used to operate pointers for cyclic memory table.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	×	MDEC2	#, r	1 1 0 1 1 r 0 0 1 1 1 1 0 1 #<7:0> – 2 #<15:8>

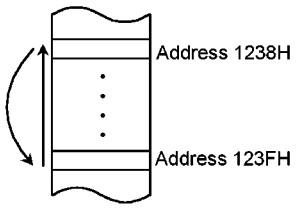
Note: The operand # must be 2 to the nth power. (n = 2 to 15)

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: Decrements the IX register by cycling from 1238H to 123FH.
MDEC2 8,IX
When the IX register = 123AH, execution sets the IX register to 1238H. Further execution increments the IX register by 8 – 2 and sets the IX register to 123EH, since the IX register modulo 8 = 0.



MDEC4 num, dst

<Modulo Decrement 4>

Operation: if (dst mod num) = 0 then dst ← dst + (num – 4) else dst ← dst – 4.

Description: When the modulo num of dst is 0, increments dst by num –4. Otherwise, decrements dst by 4. Used to operate pointers for cyclic memory table.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	×	MDEC4	#, r	1 1 0 1 1 r 0 0 1 1 1 1 1 0 #<7:0> – 4 #<15:8>

Note: The operand # must be 2 to the nth power. (n = 3 to 15)

Flags:

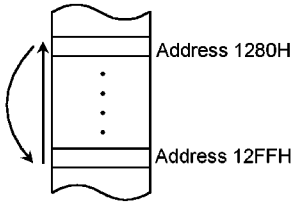
S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: Decrements the IX register by cycling from 1280H to 12FFH.

MDEC4 80H,IX

When the IX register = 1284H, execution sets the IX register to 1280H. Further execution increments the IX register by 80H – 4 and sets the IX register to 12FCH, since the IX register modulo 80H = 0.



MIN
<Minimum>

Operation: Max bit ← 0

Description: Sets the MAX bit in status register to 0. Changes the CPU operation mode to minimum.

Details:

Mnemonic	Code
MIN	0 0 0 0 0 1 0 0

Flags:

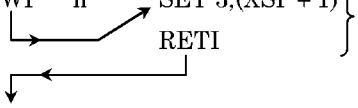
S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Note: Basically, there is no instruction for changing from minimum to maximum mode. However, if it is absolutely necessary, execute either of the following two ways.

(1) PUSH SR
SET 3, (XSP + 1)
POP SR

(2) SWI n — SET 3,(XSP + 1) / RETI — At address 8000H + n × 4H (n = 0 to 7)



MINC1 num, dst

<Modulo Increment 1>

Operation: if (dst mod num) = (num – 1) then dst ← dst – (num – 1) else dst ← dst + 1.

Description: When the modulo num of dst is num – 1, decrements dst by num – 1.
Otherwise, increments dst by 1. Used to operate pointers for cyclic memory table .

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	×	MINC1	#, r	1 1 0 1 1 r 0 0 1 1 1 0 0 0 #<7:0> – 1 #<15:8>

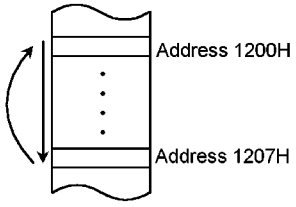
Note: The operand # must be 2 to the nth power. (n = 1 to 15)

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: Increments the IX register by cycling from 1200H to 1207H.
MINC1 8, IX
When the IX register = 1206H, execution sets the IX register to 1207H. Further execution decrements the IX register by 8 – 1 and sets the IX register to 1200H, since the IX register modulo 8 = 8 – 1.



MINC2 num, dst

<Modulo Increment 2>

Operation: if (dst mod num) = (num – 2) then dst ← dst – (num – 2) else dst ← dst + 2.

Description: When the modulo num of dst is num – 2, decrements dst by num – 2.
Otherwise, increments dst by 2. Used to operate pointers for cyclic memory table.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	×	MINC2	#, r	1 1 0 1 1 r 0 0 1 1 1 0 0 1 #<7:0> – 2 #<15:8>

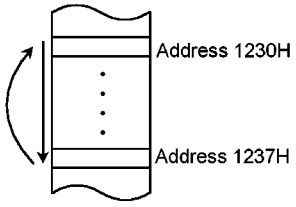
Note: The operand # must be 2 to the nth power. (n = 2 to 15)

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: Increments the IX register by cycling from 1230H to 1237H.
MINC2 8,IX
When the IX register = 1234H, execution sets the IX register to 1236H. Further execution decrements the IX register by 8 – 2 and sets the IX Register to 1230H, since the IX register modulo 8 = 8 – 2.



MINC4 num, dst

<Modulo Increment 4>

Operation: if (dst mod num) = (num – 4) then dst ← dst – (num – 4) else dst ← dst + 4.

Description: When the modulo num of dst is num – 4, decrements dst by num – 4. Otherwise, increments dst by 4. Used to operate pointers for cyclic memory table.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	×	MINC4	#, r	1 1 0 1 1 r 0 0 1 1 1 0 1 0 #<7:0> – 4 #<15:8>

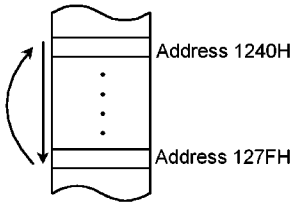
Note: The operand # must be 2 to the nth power. (n = 3 to 15)

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change

Execution example: Increments the IX register by cycling from 1240H to 127FH.
MINC4 40H,IX
When the IX register = 1278H, execution sets the IX register to 127CH. Further execution decrements the IX register by 40H – 4 and sets the IX register to 1240H, since the IX register modulo 40H = 40H – 4.



MIRR dst

<Mirror>

Operation: dst<MSB:LSB> ← dst<LSB:MSB>

Description: Mirror-exchanges the contents of dst using the bit pattern image.

Details:

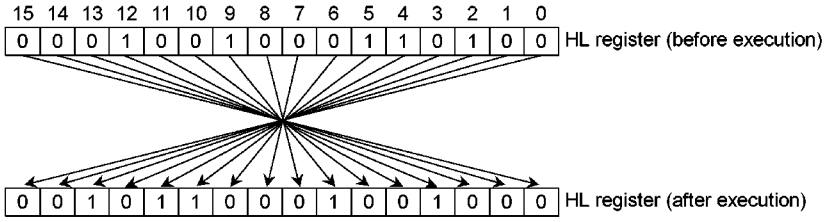
Size			Mnemonic		Code
Byte	Word	Long word			
×	O	×	MIRR	r	1 1 0 1 1 r 0 0 0 1 0 1 1 0

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change

Execution example: MIRR HL
When the HL register = 0001 0010 0011 0100B (binary), execution sets the HL register to 0010 1100 0100 1000B (binary).



MUL dst, src

<Multiply>

Operation: dst ← dst<lower half> × src (unsigned)

Description: Multiplies unsigned the contents of lower half of dst by those of src and loads the result to dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	O	×	MUL	RR, r	1 1 0 z 1 r 0 1 0 0 0 R
O	O	×	MUL	rr, #	1 1 0 z 1 r 0 0 0 0 1 0 0 0 #<7:0> #<15:8>
O	O	×	MUL	RR, (mem)	1 m 0 z m m m m 0 1 0 0 0 R

Note: When the operation is in bytes, dst (word) ← dst (byte) × src (byte).
When the operation is in words, dst (long word) ← dst (word) × src (word).
Match coding of the operand dst with the size of the result.

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change

Execution example: MUL XIX, IY
When the IX register = 1234H and the IY register = 89ABH, execution multiplies unsigned the contents of the IX register by those of the IY register and sets the XIX register to 09C9FCBCH.

Note: "RR" for the MUL RR,r and MUL RR, (mem) instructions is as listed below:

Operation size in bytes
(16 bits ← 8 bits × 8 bits)

RR	Code R
WA	001
BC	011
DE	101
HL	111
IX	Specification not possible!
IY	
IZ	
SP	

Operation size in words
(32 bits ← 16 bits × 16 bits)

RR	Code R
XWA	000
XBC	001
XDE	010
XHL	011
XIX	100
XIY	101
XIZ	110
XSP	111

*1 When the CPU is in minimum mode, XWA, XBC, XDE, or XHL cannot be used.

"rr" of the MUL rr,# instruction is as listed below.

Operation size in bytes
(16 bits ← 8 bits × 8 bits)

rr	Code r
WA	001
BC	011
DE	101
HL	111
IX	C7H : F0H
IY	C7H : F4H
IZ	C7H : F8H
SP	C7H : FCH
	1st byte 2nd byte

*2 Any other word registers can be specified in the same extension coding as those for IX to SP.

Operation size in words
(32 bits ← 16 bits × 16 bits)

rr	Code r
XWA	000
XBC	001
XDE	010
XHL	011
XIX	100
XIY	101
XIZ	110
XSP	111

*3 When the CPU is in minimum mode, XWA, XBC, XDE, or XHL cannot be used.

*4 Any other long word registers can be specified in the extension coding.

MULA dst

<Multiply and Add>

Operation: dst ← dst + (XDE) × (XHL) , XHL ← XHL – 2

Description: Multiplies signed the memory data (16 bits) specified by the XDE register by the memory data (16 bits) specified by the XHL register . Adds the result (32 bits) to the contents of dst (32 bits) and loads the sum to dst (32 bits). Then, decrements the contents of the XHL register by 2.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	×	MULA	rr	1 1 0 1 1 r 0 0 0 1 1 0 0 1

Note: Match coding of the operand dst with the operation size (long word).

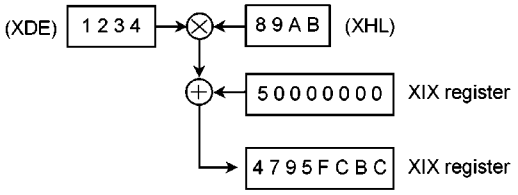
Flags:

S	Z	H	V	N	C
*	*	–	*	–	–

S = MSB value of the result is set.
Z = 1 is set when the result is 0, otherwise 0.
H = No change.
V = 1 is set when an overflow occurs as a result, otherwise 0.
N = No change.
C = No change.

Execution example: MULA XIX
Under the following conditions, execution sets the XIX register to 4795FCBCH and the XHL register to 1FEH.

Conditions:
XIX register = 50000000H
XDE register = 100H
XHL register = 200H
Memory data (word) at address 100H = 1234H
Memory data (word) at address 200H = 89ABH



MULS dst, src

<Multiply Signed>

Operation: dst ← dst<lower half> × src (signed)

Description: Multiplies signed the contents of the lower half of dst by those of src and loads the result to dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	O	×	MULS	RR, r	1 1 0 z 1 r 0 1 0 0 1 R
O	O	×	MULS	rr, #	1 1 0 z 1 r 0 0 0 0 1 0 0 1 #<7:0> #<15:8>
O	O	×	MULS	RR, (mem)	1 m 0 z m m m m 0 1 0 0 1 R

Note: When the operation is in bytes, dst(word) ← dst (byte) × src (byte).
When the operation is in words, dst (long word) ← dst (word) × src (word).
Match coding of the operand dst with the size of the result.

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change

Execution example: MULS XIX, IY
When the IX register = 1234H and the IY register = 89ABH, execution multiplies signed the contents of the IX register by those of the IY register and sets the XIX register to F795FCBCH.

Note: "RR" for the MULS RR,r and MULS RR, (mem) instructions is as listed below:

Operation size in bytes
(16 bits ← 8 bits × 8 bits)

RR	Code R
WA	001
BC	011
DE	101
HL	111
IX	Specification not possible!
IY	
IZ	
SP	

Operation size in words
(32 bits ← 16 bits × 16 bits)

RR	Code R
XWA	000
XBC	001
XDE	010
XHL	011
XIX	100
XIY	101
XIZ	110
XSP	111

*1 When the CPU is in minimum mode, XWA, XBC, XDE, or XHL cannot be used.

"rr" of the MULS rr,# instruction is as listed below.

Operation size in bytes
(16 bits ← 8 bits × 8 bits)

rr	Code r
WA	001
BC	011
DE	101
HL	111
IX	C7H : F0H
IY	C7H : F4H
IZ	C7H : F8H
SP	C7H : FCH
	1st byte 2nd byte

*2 Any other word registers can be specified in the same extension coding as those for IX to SP.

Operation size in words
(32 bits ← 16 bits × 16 bits)

rr	Code r
XWA	000
XBC	001
XDE	010
XHL	011
XIX	100
XIY	101
XIZ	110
XSP	111

*3 When the CPU is in minimum mode, XWA, XBC, XDE, or XHL cannot be used.

*4 Any other long word registers can be specified in the extension coding.

NEG dst

<Negate>

Operation: dst ← 0 – dst

Description: Decrements 0 by the contents of dst and loads the result to dst.
(Twos complement)

Details:

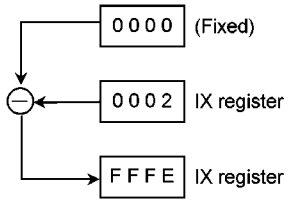
Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	NEG	r	1 1 0 z 1 r 0 0 0 0 0 1 1 1

Flags:

S	Z	H	V	N	C
*	*	*	*	1	*

S = MSB value of the result is set.
Z = 1 is set when the result is 0, otherwise 0.
H = 1 is set when a borrow from bit 3 to bit 4 occurs as a result, otherwise 0.
V = 1 is set when an overflow occurs as a result, otherwise 0.
N = 1 is set.
C = 1 is set when a borrow from the MSB occurs as a result, otherwise 0.

Execution example: NEG IX
When the IX register = 0002H, execution sets the IX register to FFFEH.



NOP

<No Operation>

Operation: None.

Description: Does nothing but moves execution to the next instruction. The object code of this instruction is 00H.

Details:

Mnemonic	Code
NOP	0 0 0 0 0 0 0 0

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change

OR dst, src

<Logical OR>

Operation: dst ← dst OR src

Description: Ors the contents of dst with those of src and loads the result to dst.

(Truth table)

A	B	A or B
0	0	0
0	1	1
1	0	1
1	1	1

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	OR	R, r	1 1 z z 1 r 1 1 1 0 0 R
○	○	○	OR	r, #	1 1 z z 1 r 1 1 0 0 1 1 1 0 #<7:0> #<15:8> #<23:16> #<31:24>
○	○	○	OR	R, (mem)	1 m z z m m m m 1 1 1 0 0 R
○	○	○	OR	(mem), R	1 m z z m m m m 1 1 1 0 1 R
○	○	×	OR<W>	(mem), #	1 m 0 z m m m m 0 0 1 1 1 1 1 0 #<7:0> #<15:8>

Flags:

S	Z	H	V	N	C
*	*	0	*	0	0

S = MSB value of the result is set.
Z = 1 is set when the result is 0, otherwise 0.
H = 0 is set.
V = 1 is set when the parity (number of 1s) of the result is even, 0 when odd.
When the operand is 32-bit, an undefined value is set.
N = Cleared to 0.
C = Cleared to 0.

Execution example: OR HL, IX
When the HL register = 7350H and the IX register is 3456H, execution sets the HL register to 7756H.

```
        0111  0011  0101  0000  ← HL register (before execution)
OR)     0011  0100  0101  0110  ← IX register (before execution)
-------------------------------
        0111  0111  0101  0110  ← HL register (after execution)
```

ORCF num, src

<OR Carry Flag>

Operation: CY ← CY OR src<num>

Description: Ors the contents of the carry flag with those of bit num of src and loads the result to the carry flag.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	O	×	ORCF	#4, r	1 1 0 z 1 r 0 0 1 0 0 0 0 1 0 0 0 0 # 4
O	O	×	ORCF	A, r	1 1 0 z 1 r 0 0 1 0 1 0 0 1
O	×	×	ORCF	#3, (mem)	1 m 1 1 m m m m 1 0 0 0 1 #3
O	×	×	ORCF	A, (mem)	1 m 1 1 m m m m 0 0 1 0 1 0 0 1

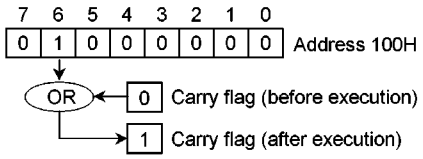
Note: When bit num is specified by the A register, the value of the lower 4 bits of the A register is used as bit num. When the operand is a byte and the value of the lower bits of bit num is from 8 to 15, the result is undefined.

Flags:

S	Z	H	V	N	C
–	–	–	–	–	*

S = No change
Z = No change
H = No change
V = No change
N = No change
C = The result of or-ing the contents of the carry flag with those of bit num of src is set.

Execution example: ORCF 6, (100H)
When the contents of memory at address 100H = 01000000B (binary) and the carry flag = 0, execution sets the carry flag to 1.



PAA dst

<Pointer Adjust Accumulator>

Operation: if dst <LSB> = 1 then dst ← dst + 1

Description: Increments dst by 1 when the LSB of dst is 1. Does nothing when the LSB of dst is 0.
Used to make the contents of dst even. With the TLCS-900 series, when accessing 16- or 32-bit data in memory, if the data are loaded from an address starting with an even number, the number of bus cycles is 1 less than that of the data loaded from an address starting with an odd number.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	O	PAA	r	1 1 z z 1 r 0 0 0 1 0 1 0 0

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change

Execution example: PAA XIZ
When the XIZ register = 00234567H, execution increments the XIZ register by 1 so that it becomes 00234568H.

POP dst

<Pop>

Operation: dst ← (XSP+)

In bytes	: dst ← (XSP), XSP ← XSP + 1
In words	: dst ← (XSP), XSP ← XSP + 2
In long words	: dst ← (XSP), XSP ← XSP + 4

Description: First loads the contents of memory address specified by the stack pointer XSP to dst. Then increments the stack pointer XSP by the number of bytes in the operand.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	×	×	POP	F	0 0 0 1 1 0 0 1
O	×	×	POP	A	0 0 0 1 0 1 0 1
×	O	O	POP	R	0 1 0 s 1 R
O	O	O	POP	r	1 1 z z 1 r 0 0 0 0 0 1 0 1
O	O	×	POP<W>	(mem)	1 m 1 1 m m m m 0 0 0 0 0 1 z 0

Flags:

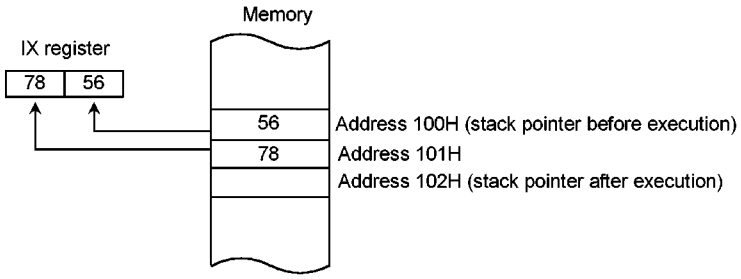
S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Note: Executing POP F changes all flags.

Execution example: POP IX

When the stack pointer XSP = 0100H, the contents of address 100H = 56H, and the contents of address 101H = 78H, execution sets the IX register to 7856H and the stack pointer XSP to 0102H.



POP SR

<Pop SR>

Operation: SR ← (XSP+)

Description: Loads the contents of the address specified by the stack pointer XSP to status register. Then increments the contents of the stack pointer XSP by 2.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	×	POP	SR	0 0 0 0 0 0 1 1

Flags:

S	Z	H	V	N	C
*	*	*	*	*	*

S =
Z =
H =
V =
N =
C =
} Contents of the memory address specified by the stack pointer XSP are set.

Note: Please execute this instruction during DI condition.
The timing for executing this instruction is delayed by several states than that for fetching the instruction. This is because an instruction queue (4 bytes) and pipeline processing method is used.

PUSH SR

<Push SR>

Operation: (–XSP) ← SR

Description: Decrements the contents of the stack pointer XSP by 2. Then loads the contents of status register to the memory address specified by the stack pointer XSP.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	O	×	PUSH	SR	0 0 0 0 0 0 1 0

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

PUSH src

<Push>

Operation: (–XSP) ← src

In bytes	: XSP ← XSP – 1, (XSP) ← src
In words	: XSP ← XSP – 2, (XSP) ← src
In long words	: XSP ← XSP – 4, (XSP) ← src

Description: Decrements the stack pointer XSP by the byte length of the operand. Then loads the contents of src to the memory address specified by the stack pointer XSP.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	×	×	PUSH	F	0 0 0 1 1 0 0 0
O	×	×	PUSH	A	0 0 0 1 0 1 0 0
×	O	O	PUSH	R	0 0 1 s 1 R
O	O	O	PUSH	r	1 1 z z 1 r 0 0 0 0 0 1 0 0
O	O	×	PUSH<W>	#	0 0 0 0 1 0 z 1 #<7:0> #<15:8>
O	O	×	PUSH<W>	(mem)	1 m 0 z m m m m 0 0 0 0 0 1 0 0

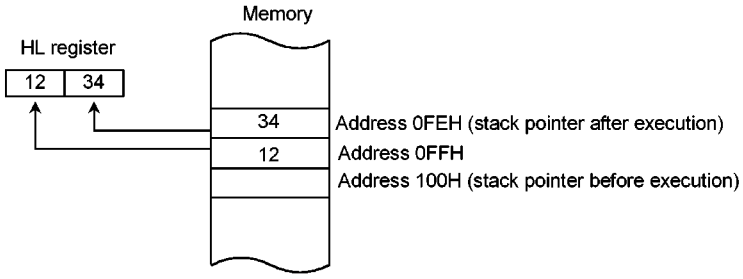
Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: PUSH HL

When the stack pointer XSP = 0100H and the HL register = 1234H, execution changes address 00FEH to 34H, address 00FFH to 12H, and sets the stack pointer XSP to 00FEH.



RCF

<Reset Carry Flag>

Operation: CY ← 0

Description: Resets the carry flag to 0.

Details:

Size	Mnemonic	Code
	RCF	0 0 0 1 0 0 0 0

Flags:

S	Z	H	V	N	C
–	–	0	–	0	0

S = No change
Z = No change
H = Reset to 0.
V = Reset to 0.
N = No change
C = Reset to 0.

RES num, dst

<Reset>

Operation: dst <num> ← 0

Description: Resets bit num of dst to 0.

Details:

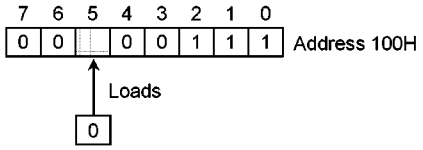
Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	RES	#4, r	1 1 0 z 1 r 0 0 1 1 0 0 0 0 0 0 0 0 # 4
○	×	×	RES	#3, (mem)	1 m 1 1 m m m m 1 0 1 1 0 #3

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: RES 5, (100H)
When the contents of memory at address 100H = 00100111B (binary), execution sets the contents to 00000111B (binary).



RET condition

<Return>

Operation: In minimum mode: If cc is true, then the 16-bit PC ← (XSP),
XSP ← XSP + 2.
In minimum mode: If cc is true, then the 32-bit PC ← (XSP),
XSP ← XSP + 2.

Description: Pops the return address from the stack area to the program counter when the operand condition is true.

Details:

Size	Mnemonic		Code
	RET		0 0 0 0 1 1 1 0
	RET	cc	1 0 1 1 0 0 0 0 1 1 1 1 c c

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: RET
When the stack pointer XSP = 0FEH and the contents of memory at address 0FCH = 9000H (word data) in minimum mode, execution sets the stack pointer XSP to 100H and jumps (returns) to address 9000H.

RETD num

<Return and Deallocate>

Operation: In minimum mode: 32-bit PC ← (XSP), XSP ← XSP + 4, XSP ← XSP + num
In maximum mode: 32-bit PC ← (XSP), XSP ← XSP + 4, XSP ← XSP + num

Description: Pops the return address from the stack area to the program counter. Then increments the stack pointer XSP by signed num.

Details:

Mnemonic		Code
RETD	d16	0 0 0 0 1 1 1 1 d<7:0> d<15:8>

Flags:

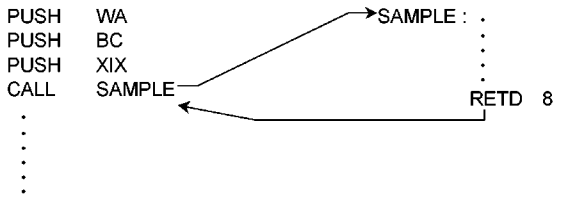
S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: RETD 8

When the stack pointer XSP = 0FCH and the contents of memory at address 0FCH = 9000H (long word data) in minimum mode, execution sets the stack pointer XSP to 0FCH + 4 + 8 → 108H and jumps (returns) to address 9000H.

Usage of the RETD instruction is shown below. In this example, the 8-bit parameter is pushed to the stack before the subroutine call. After the subroutine processing complete, the used parameter area is deleted by the RETD instruction.



RETI

<Return from Interrupt>

Operation: In minimum mode: SR ← (XSP), 16-bit PC ← (XSP + 2),
XSP ← XSP + 4
In maximum mode: SR ← (XSP), 32-bit PC ← (XSP + 2),
XSP ← XSP + 6

After the above operation is executed, the 900/L1 decrement a value of interrupt nesting counter INTNEST by 1.

Description: Pops data from the stack area to the 2-byte Temp register and program counter. Next, loads the contents of the Temp register to status register.

After the above operation is executed, the 900/L1 decrement a value of interrupt nesting counter INTNEST by 1.

Details:

Mnemonic	Code
RETI	0 0 0 0 0 1 1 1

Flags:

S	Z	H	V	N	C
*	*	*	*	*	*

S = The value popped from the stack area is set.
Z = The value popped from the stack area is set.
H = The value popped from the stack area is set.
V = The value popped from the stack area is set.
N = The value popped from the stack area is set.
C = The value popped from the stack area is set.

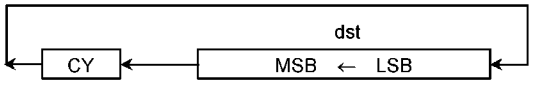
RL num, dst

<Rotate Left>

Operation: {CY & dst ← left rotates the value of CY & dst} Repeat num

Description: Rotates left the contents of the linked carry flag and dst.
Repeats the number of times specified in num.

Description figure:



Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	RL	#4, r	1 1 z z 1 r 1 1 1 0 1 0 1 0 0 0 0 0 # 4
○	○	○	RL	A, r	1 1 z z 1 r 1 1 1 1 1 0 1 0
○	○	×	RL<W>	(mem)	1 m 0 z m m m m 0 1 1 1 1 0 1 0

Note: When the number of rotates is specified by the A register, the value of the lower 4 bits of the A register is used. Specifying 0 rotates 16 times.
When dst is memory, rotating is performed only once.

Flags:

S	Z	H	V	N	C
*	*	0	*	0	*

S = MSB value of dst after rotate is set.
Z = 1 is set when the contents of dst after rotate is 0, otherwise 0.
H = Reset to 0.
V = 1 is set when the parity (number of 1s) of dst is even after rotate, otherwise 0.
If the perand is 32 bits, an undefined value is set.
N = Reset to 0.
C = The value after rotate is set.

Execution example: RL 4, HL
When the HL register = 6230H and the carry flag = 1, execution sets the HL register to 230BH and the carry flag to 0.

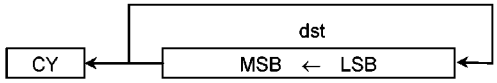
RLC num, dst

<Rotate Left without Carry>

Operation: {CY ← dst <MSB>, dst ← left rotate value of dst} Repeat num

Description: Loads the contents of the MSB of dst to the carry flag and rotates left the contents of dst. Repeats the number of times specified in num.

Description figure:



Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	RLC	#4, r	1 1 z z 1 r 1 1 1 0 1 0 0 0 0 0 0 0 # 4
○	○	○	RLC	A, r	1 1 z z 1 r 1 1 1 1 1 0 0 0
○	○	×	RLC<W>	(mem)	1 m 0 z m m m m 0 1 1 1 1 0 0 0

Note: When the number of rotates is specified by the A register, the value of the lower 4 bits of the A register is used. Specifying 0 rotates 16 times.
When dst is memory, rotating is performed only once.

Flags:

S	Z	H	V	N	C
*	*	0	*	0	*

S = MSB value of dst after rotate is set.
Z = 1 is set when the contents of dst after rotate is 0, otherwise, 0.
H = Reset to 0.
V = 1 is set when the parity (number of 1s) of dst is even after rotate.
If the operand is 32 bits, an undefined value is set.
N = Reset to 0.
C = MSB value of dst before the last rotate is set.

Execution example: RLC 4, HL
When the HL register = 1230H, execution sets the HL register to 2301H and the carry flag to 1.

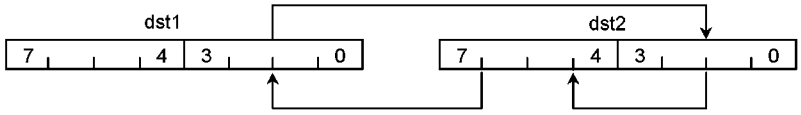
RLD dst1, dst2

<Rotate Left Digit>

Operation: dst1<3:0> ← dst2<7:4>, dst2<7:4> ← dst2<3:0>, dst2<3:0> ← dst1 <3:0>

Description: Rotates left the lower 4 bits of dst1 and the contents of dst2 in units of 4 bits.

Description figure:



Details:

Size			Mnemonic		Code
Byte	Word	Long word			
O	×	×	RLD	[A,] (mem)	1 m 0 0 m m m m 0 0 0 0 0 1 1 0

Flags:

S	Z	H	V	N	C
*	*	0	*	0	–

S = MSB value of the A register after rotate is set.
Z = 1 is set when the contents of the A register after the rotate are 0, otherwise 0.
H = Reset to 0.
V = 1 is set when the parity (number of 1s) of the A register is even after the rotate, otherwise 0.
N = Reset to 0.
C = No change

Execution example: RLD A, (100H)
When the A register = 12H and the contents of memory at address 100H = 34H, execution sets the A register to 13H and the contents of memory at address 100H to 42H.

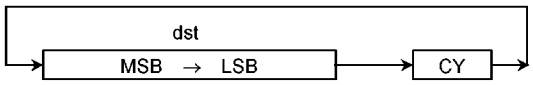
RR num, dst

<Rotate Right>

Operation: {CY & dst ← right rotates the value of CY & dst} Repeat num

Description: Rotates right the linked contents of the carry flag and dst.
Repeats the number of times specified in num.

Description figure:



Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	RR	#4, r	1 1 z z 1 r 1 1 1 0 1 0 1 1 0 0 0 0 # 4
○	○	○	RR	A, r	1 1 z z 1 r 1 1 1 1 1 0 1 1
○	○	×	RR<W>	(mem)	1 m 0 z m m m m 0 1 1 1 1 0 1 1

Note: When the number of rotates is specified by the A register, the value of the lower 4 bits of the A register is used. Specifying 0 rotates 16 times.
When dst is memory, rotating is performed only once.

Flags:

S	Z	H	V	N	C
*	*	0	*	0	*

S = MSB value of dst after rotate is set.
Z = 1 is set when the contents of dst after rotate is 0, otherwise 0.
H = Reset to 0.
V = 1 is set when the parity (number of 1s) of dst is even after the rotate, otherwise 0.
If the operand is 32 bits, an undefined value is set.
N = Reset to 0.
C = The value after rotate is set.

Execution example: RR 4, HL
When the HL register = 6230H and the carry flag = 1, execution sets the HL register to 1623H and the carry flag to 0.

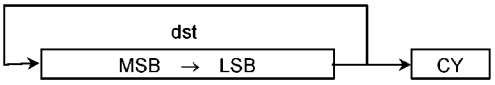
RRC num, dst

<Rotate Right without Carry>

Operation: {CY ← dst <LSB>, dst ← right rotate value of dst} Repeat num

Description: Loads the contents of the LSB of dst to the carry flag and rotates the contents of dst to the right. Repeats the number of times specified in num.

Description figure:



Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	RRC	#4, r	1 1 z z 1 r 1 1 1 0 1 0 0 1 0 0 0 0 # 4
○	○	○	RRC	A, r	1 1 z z 1 r 1 1 1 1 1 0 0 1
○	○	×	RRC<W>	(mem)	1 m 0 z m m m m 0 1 1 1 1 0 0 1

Note: When the number of rotates num is specified by the A register, the value of the lower 4 bits of the A register is used as the number of rotates.
Specifying 0 rotates 16 times. When dst is memory, rotating is only once.

Flags:

S	Z	H	V	N	C
*	*	0	*	0	*

S = MSB value of dst after rotate is set.
Z = 1 is set when the contents of dst after rotate is 0, otherwise 0.
H = Reset to 0.
V = 1 is set when the parity (number of 1s) of dst is even after rotate, otherwise 0.
If the operand is 32 bits, an undefined value is set.
N = Reset to 0.
C = MSB value of dst before the last rotate is set.

Execution example: RLC 4, HL
When the HL register = 1230H, execution sets the HL register to 0123H and the carry flag to 0.

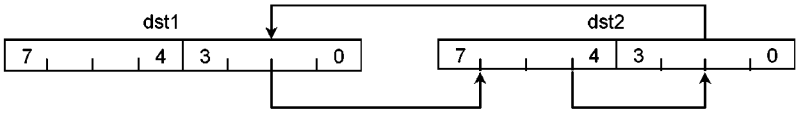
RRD dst1, dst2

<Rotate Right Digit>

Operation: dst1<3:0> ← dst2<3:0>, dst2<7:4> ← dst1<3:0>, dst2<3:0> ← dst2<7:4>

Description: Rotates right the lower 4 bits of dst1 and the contents of dst2 in units of 4 bits.

Description figure:



Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	×	×	RRD	[A,](mem)	1 m 0 0 m m m m 0 0 0 0 0 1 1 1

Flags:

S	Z	H	V	N	C
*	*	0	*	0	–

S = MSB value of the A register after rotate is set.
Z = 1 is set when the contents of the A register after rotate is 0, otherwise 0.
H = Reset to 0.
V = 1 is set when the parity (number of 1s) of the A register is even after rotate, otherwise 0.
N = Reset to 0.
C = No change

Execution example: RRD A, (100H)
When the A register = 12H and the contents of memory at address 100H = 34H, execution sets the A register to 14H and the contents of memory at address 100H to 23H.

SBC dst, src

<Subtract with Carry>

Operation: dst ← dst – src – CY

Description: Subtracts the contents of src and the carry flag from those of dst, and loads the result to dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	SBC	R, r	1 1 z z 1 r 1 0 1 1 0 R
○	○	○	SBC	r, #	1 1 z z 1 r 1 1 0 0 1 0 1 1 #<7:0> #<15:8> #<23:16> #<31:24>
○	○	○	SBC	R, (mem)	1 m z z m m m m 1 0 1 1 0 R
○	○	○	SBC	(mem), R	1 m z z m m m m 1 0 1 1 1 R
○	○	×	SBC<W>	(mem), #	1 m 0 z m m m m 0 0 1 1 1 0 1 1 #<7:0> #<15:8>

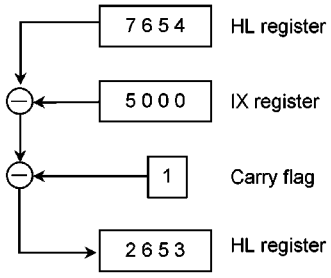
Flags:

S	Z	H	V	N	C
*	*	*	*	1	*

S = MSB value of the result is set.
Z = 1 is set when the result is 0, otherwise 0.
H = 1 is set when a borrow from bit 3 to bit 4 occurs as a result, otherwise 0.
When the operand is 32 bits, an undefined value is set.
V = 1 is set when an overflow occurs as a result, otherwise 0.
N = 1 is set.
C = 1 is set when a borrow from the MSB occurs as a result, otherwise 0.

Execution example: SBC HL, IX

When the HL register is 7654H, the IX register = 5000H, and the carry flag = 1, execution sets the HL register to 2653H.



SCC condition, dst

<Set Condition Code>

Operation: If cc is true, then dst ← 1 else dst ← 0.

Description: Loads 1 to dst when the operand condition is true; when false, 0 is loaded to dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	SCC	cc, r	1 1 0 z 1 r 0 1 1 1 c c

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: SCC OV, HL
When the contents of the V flag = 1, execution sets the HL register to 0001H.

SCF
<Set Carry Flag>

Operation: CY ← 1

Description: Sets the carry flag to 1.

Details:

Mnemonic	Code
SCF	0 0 0 1 0 0 0 1

Flags:

S	Z	H	V	N	C
–	–	0	–	0	1

S = No change
Z = No change
H = Reset to 0.
V = No change
N = Reset to 0.
C = Set to 1.

SET num, dst

<Set>

Operation: dst <num> ← 1

Description: Sets bit num of dst to 1.

Details:

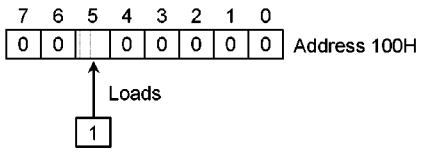
Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	SET	#4, r	1 1 0 z 1 r 0 0 1 1 0 0 0 1 0 0 0 0 # 4
○	×	×	SET	#3, (mem)	1 m 1 1 m m m m 1 0 1 1 1 #3

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: SET 5, (100H)
When the contents of memory at address 100H = 00000000B (binary), execution sets the contents of memory at address 100H to 00100000B (binary).



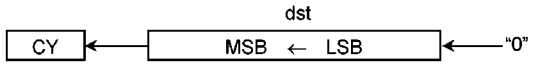
SLA num, dst

<Shift Left Arithmetic>

Operation: {CY ← dst<MSB>, dst ← left shift value of dst, dst<LSB> ← 0} Repeat num

Description: Loads the contents of the MSB of dst to the carry flag, shifts left the contents of dst, and loads 0 to the LSB of dst. Repeats the number of times specified in num.

Description chart:



Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	SLA	#4, r	1 1 z z 1 r 1 1 1 0 1 1 0 0 0 0 0 0 # 4
○	○	○	SLA	A, r	1 1 z z 1 r 1 1 1 1 1 1 0 0
○	○	×	SLA<W>	(mem)	1 m 0 z m m m m 0 1 1 1 1 1 0 0

Note: When the number of shifts, num, is specified by the A register, the value of the lower 4 bits of the A register is used. Specifying 0 shifts 16 times. When dst is memory, shifting is performed only once.

Flags:

S	Z	H	V	N	C
*	*	0	*	0	*

S = MSB value of dst after shift is set.
Z = 1 is set when the contents of dst after shift is 0, otherwise 0.
H = Reset to 0.
V = 1 is set when the parity (number of 1s) of dst is even after shifting, otherwise 0.
If the operand is 32 bits, an undefined value is set.
N = Reset to 0.
C = MSB value of dst before the last shift is set.

Execution example: SLA 4, HL
When the HL register = 1234H, execution sets the HL register to 2340H and the carry flag to 1.

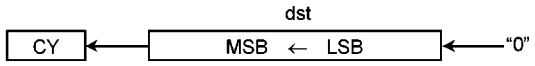
SLL num, dst

<Shift Left Logical>

Operation: {CY ← dst<MSB>, dst ← left shift value of dst, dst<LSB> ← 0} Repeat num

Description: Loads the contents of the MSB of dst to the carry flag, shifts left the contents of dst, and loads 0 to the MSB of dst. Repeats the number of times specified in num.

Description chart:



Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	SLL	#4, r	1 1 z z 1 r 1 1 1 0 1 1 1 0 0 0 0 0 # 4
○	○	○	SLL	A, r	1 1 z z 1 r 1 1 1 1 1 1 1 0
○	○	×	SLL<W>	(mem)	1 m 0 z m m m m 0 1 1 1 1 1 1 0

Note: When the number of shifts, num, is specified by the A register, the value of the lower 4 bits of the A register is used. Specifying 0 shifts 16 times. When dst is memory, shifting is performed only once.

Flags:

S	Z	H	V	N	C
*	*	0	*	0	*

S = MSB value of dst after shift is set.
Z = 1 is set when the contents of dst after shift is 0, otherwise 0.
H = Reset to 0.
V = 1 is set when the parity (number of 1s) of dst is even after shifting, otherwise 0.
If the operand is 32 bits, an undefined value is set.
N = Reset to 0.
C = MSB value of dst before the last shift is set.

Execution example: SLL 4, HL
When the HL register = 1234H, execution sets the HL register to 2340H and the carry flag to 1.

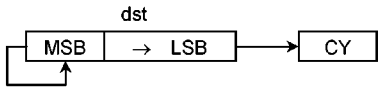
SRA num, dst

<Shift Right Arithmetic>

Operation: {CY ← dst<MSB>, dst ← right shift value of dst, dst <MSB> is fixed} Repeatnum

Description: Loads the contents of the LSB of dst to the carry flag and shifts right the contents of dst (MSB is fixed). Repeats the number of times specified in num.

Description chart:



Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	SRA	#4, r	1 1 z z 1 r 1 1 1 0 1 1 0 0 0 0 0 0 # 4
○	○	○	SRA	A, r	1 1 z z 1 r 1 1 1 1 1 1 0 1
○	○	×	SRA<W>	(mem)	1 m 0 z m m m m 0 1 1 1 1 1 0 1

Note: When the number of shifts, num, is specified by the A register, the value of the lower 4 bits of the A register is used. Specifying 0 shifts 16 times. When dst is memory, shifting is performed only once.

Flags:

S	Z	H	V	N	C
*	*	0	*	0	*

S = MSB value of dst after shift is set.
Z = 1 is set when the contents of dst after shift is 0, otherwise 0.
H = Reset to 0.
V = 1 is set when the parity (number of 1s) of dst is even after shift, otherwise 0.
If the operand is 32 bits, an undefined value is set.
N = Reset to 0.
C = LSB value of dst before the last shift is set.

Execution example: SRA 4, HL
When the HL register = 8230H, execution sets the HL register to F823H and the carry flag to 0.

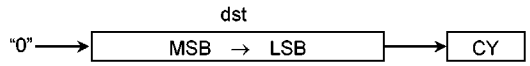
SRL num, dst

<Shift Right Logical>

Operation: {CY ← dst<LSB>, dst ← right shift value of dst, dst <MSB> ← 0} Repeat num

Description: Loads the contents of the LSB of dst to the carry flag, shifts right the contents of dst, and loads 0 to the MSB of dst. Repeats the number of times specified in num.

Description chart:



Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	SRL	#4, r	1 1 z z 1 r 1 1 1 0 1 1 1 1 0 0 0 0 # 4
○	○	○	SRL	A, r	1 1 z z 1 r 1 1 1 1 1 1 1 1
○	○	×	SRL<W>	(mem)	1 m 0 z m m m m 0 1 1 1 1 1 1 1

Note: When the number of shifts, num, is specified by the A register, the value of the lower 4 bits of the A register is used. Specifying 0 shifts 16 times. When dst is memory, shifting is performed only once.

Flags:

S	Z	H	V	N	C
*	*	0	*	0	*

S = MSB value of dst after shift is set.
Z = 1 is set when the contents of dst after shift is 0, otherwise 0.
H = Reset to 0.
V = 1 is set when the parity (number of 1s) of dst is even after shift, otherwise 0.
If the operand is 32 bits, an undefined value is set.
N = Reset to 0.
C = LSB value of dst before the last shift is set.

Execution example: SRL 4, HL
When the HL register = 1238H, execution sets the HL register to 0123H and the carry flag to 1.

STCF num, dst

<Store Carry Flag>

Operation: dst<num> ← CY

Description: Loads the contents of the carry flag to bit num of dst.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	STCF	#4, r	1 1 0 z 1 r 0 0 1 0 0 1 0 0 0 0 0 0 # 4
○	○	×	STCF	A, r	1 1 0 z 1 r 0 0 1 0 1 1 0 0
○	×	×	STCF	#3, (mem)	1 m 1 1 m m m m 1 0 1 0 0 #3
○	×	×	STCF	A, (mem)	1 m 1 1 m m m m 0 0 1 0 1 1 0 0

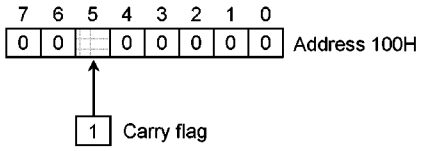
Note: When bit num is specified by the A register, the value of the lower 4 bits of the A register is used. When the operand is a byte and the value of the lower 4 bits of bit num is from 8 to 15, the operand value does not change.

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: STCF 5, (100H)
When the contents of memory at address 100H = 00H and the carry flag = 1, execution sets the contents of memory at address 100H to 00100000B (binary).



SUB dst, src

<Subtract>

Operation: dst ← dst – src

Description: Subtracts the contents of src from those of dst and loads the result to dst.

Details:

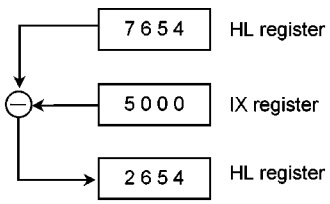
Size			Mnemonic		Code
Byte	Word	Long word			
O	O	O	SUB	R, r	1 1 z z 1 r 1 0 1 0 0 R
O	O	O	SUB	r, #	1 1 z z 1 r 1 1 0 0 1 0 1 0 #<7:0> #<15:8> #<23:16> #<31:24>
O	O	O	SUB	R, (mem)	1 m z z m m m m 1 0 1 0 0 R
O	O	O	SUB	(mem), R	1 m z z m m m m 1 0 1 0 1 R
O	O	×	SUB<W>	(mem), #	1 m 0 z m m m m 0 0 1 1 1 0 1 0 #<7:0> #<15:8>

Flags:

S	Z	H	V	N	C
*	*	*	*	1	*

S = MSB value of the result is set.
Z = 1 is set when the result is 0, otherwise 0.
H = 1 is set when a borrow from bit 3 to bit 4 occurs as a result, otherwise 0.
When the operand is 32 bits, an undefined value is set.
V = 1 is set when an overflow occurs as a result, otherwise 0.
N = 1 is set.
C = 1 is set when a borrow from MSB occurs as a result, otherwise 0.

Execution example: SUB HL, IX
When the HL register = 7654H and the IX register = 5000H, execution sets the HL register to 2654H.



SWI num

<Software Interrupt>

Operation:
1) XSP ← XSP – 4...in minimum mode
   or
   XSP ← XSP – 6...in maximum mode
2) (XSP) ← SR
3) (XSP + 2) ← 16 bit PC...in minimum mode
   or
   (XSP + 2) ← 32 bit PC...in maximum mode
4) PC ← (Address refer to vector + num × 4)
5) INTNEST ← INTNEST + 1

Note: Address refer to vector is defined for each product.

Description: Saves to the stack area the contents of the status register and contents of the program counter which indicate the address next to the SWI instruction. Finally, jumps to vector is indicated address refer to vector.

The SWI0 to 7 are interrupts of level 7. The interrupt level mask register IFF2 to 0 are not changed by executing this instruction. For example, when an interrupt of level 1 is requested in progress of the SWI interrupt routine at IFF2 to 0 = 1, acceptance of the interrupt is nesting.

Details:

Size	Mnemonic		Code
	SWI	[#3]	1 1 1 1 1 #3

Note 1: A value from 0 to 7 can be specified as the operand value. When the operand coding is omitted, SWI 7 is assumed.

Note 2: The status register structure is as shown below.

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
SYSM	1FF2	1FF1	1FF0	MAX	RFP2	RFP1	RFP0	S	Z	"0"	H	"0"	V	N	C

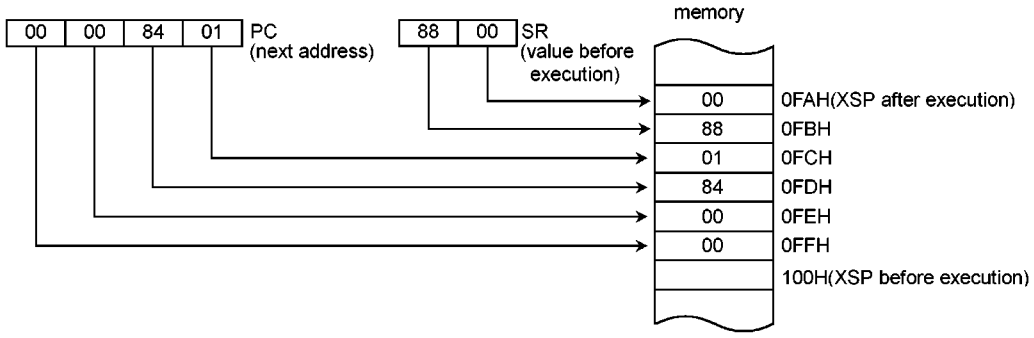
Flags:

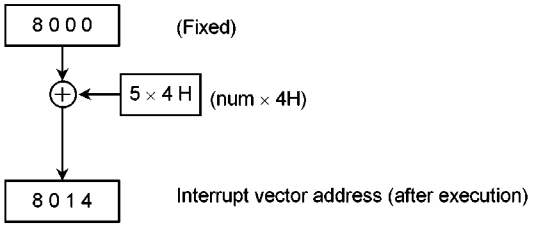
S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: SWI 5

When the stack pointer XSP = 100H and the status register SR is 8800H, executing writes the contents of the previous status register SR (8800H) into address 00FAH, and the contents of the program counter PC (00008401H) into address 00FCH, and the interrupt vector which is defined at address 8014 (in TMP93CS40) is read. Finally, jumps to address 8014H.





TSET num, dst

<Test and Set>

Operation: Z flag ← inverted value of dst <num>
dst <num> ← 1

Description: Loads the inverted value of the bit num of dst to the Z flag.
Then the bit num of dst is set to "1".

Details:

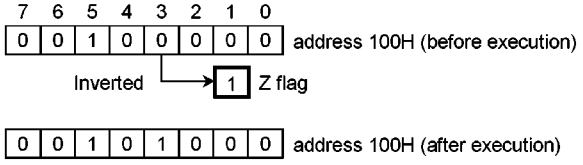
Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	TSET	#4, r	1 1 z z 1 r 0 0 1 1 0 1 0 0 0 0 0 0 # 4
○	×	×	TSET	#3, (mem)	1 m 1 1 m m m m 1 0 1 0 1 #3

Flags:

S	Z	H	V	N	C
×	*	1	×	0	–

S = An undefined value is set.
Z = The inverted value of the src <num> is set.
H = Set to 1
V = An undefined value is set.
N = Set to 0
C = No change

Execution example: TSET 3, (100H)
When the contents of memory at address 100H = 00100000B (binary), execution sets the Z flag to 1, the contents of memory at address 100H = 00101000B (binary).



UNLK dst

<Unlink>

Operation: XSP ← dst, dst ← (XSP+)

Description: Loads the contents of dst to the stack pointer XSP, then pops long word data from the stack area to dst. Used paired with the Link instruction.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
×	×	O	UNLK	r	1 1 1 0 1 r 0 0 0 0 1 1 0 1

Flags:

S	Z	H	V	N	C
–	–	–	–	–	–

S = No change
Z = No change
H = No change
V = No change
N = No change
C = No change

Execution example: UNLK XIZ

As a result of executing this instruction after executing the Link instruction, the stack pointer XSP and the XIZ register revert to the same values they had before the Link instruction was executed.

(For the details, see the Link instruction, page 100)

XOR dst, src

<Exclusive OR>

Operation: dst ← dst XOR src

Description: Exclusive ors the contents of dst with those of src and loads the result to dst.

(Truth table)

A	B	A XOR B
0	0	0
0	1	1
1	0	1
1	1	0

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	○	XOR	R, r	1 1 z z 1 r 1 1 0 1 0 R
○	○	○	XOR	r, #	1 1 z z 1 r 1 1 0 0 1 1 0 1 #<7:0> #<15:8> #<23:16> #<31:24>
○	○	○	XOR	R, (mem)	1 m z z m m m m 1 1 0 1 0 R
○	○	○	XOR	(mem), R	1 m z z m m m m 1 1 0 1 1 R
○	○	×	XOR<W>	(mem), #	1 m 0 z m m m m 0 0 1 1 1 1 0 1 #<7:0> #<15:8>

Flags:

S	Z	H	V	N	C
*	*	0	*	0	0

S = MSB value of the result is set.
Z = 1 is set when the result is 0, otherwise 0.
H = Reset to 0.
V = 1 is set when the parity (number of 1s) of dst is even as a result, otherwise 0.
If the operand is 32 bits, an undefined value is set.
N = Cleared to 0.
C = Cleared to 0.

Execution example: XOR HL, IX
When the HL register = 7350H and the IX register = 3456H, execution sets the HL register to 4706H.

```
        0111  0011  0101  0000  ← HL register (before execution)
XOR)    0011  0100  0101  0110  ← IX register (before execution)
-------------------------------
        0100  0111  0000  0110  ← HL register (after execution)
```

XORCF num, src

<Exclusive OR Carry Flag>

Operation: CY ← CY XOR src<num>

Description: Exclusive ors the contents of the carry flag and bit num of src, and loads the result to the carry flag.

Details:

Size			Mnemonic		Code
Byte	Word	Long word			
○	○	×	XORCF	#4, r	1 1 0 z 1 r 0 0 1 0 0 0 1 0 0 0 0 0 # 4
○	○	×	XORCF	A, r	1 1 0 z 1 r 0 0 1 0 1 0 1 0
○	×	×	XORCF	#3, (mem)	1 m 1 1 m m m m 1 0 0 1 0 #3
○	×	×	XORCF	A, (mem)	1 m 1 1 m m m m 0 0 1 0 1 0 1 0

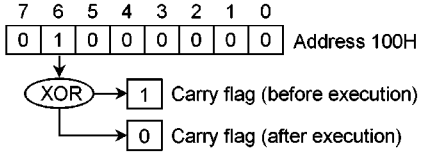
Note: When bit num is specified by the A register, the value of the lower 4 bits of the A register is used. When the operand is a byte and the value of the lower 4 bits of bit num is from 8 to 15, the result is undefined.

Flags:

S	Z	H	V	N	C
–	–	–	–	–	*

S = No change
Z = No change
H = No change
V = No change
N = No change
C = The value obtained by exclusive or-ing the contents of the carry flag with those of bit num of src is set.

Execution example: XORCF 6, (100H)
When the contents of memory at address 100H = 01000000B (binary) and the carry flag = 1, execution sets the carry flag to 0.



ZCF

<Zero flag to Carry Flag>

Operation: CY ← inverted value of Z flag

Description: Loads the inverted value of the Z flag to the carry flag.

Details:

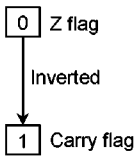
Mnemonic	Code
ZCF	0 0 0 1 0 0 1 1

Flags:

S	Z	H	V	N	C
–	–	×	–	0	*

S = No change
Z = No change
H = An undefined value is set.
V = No change
N = Reset to 0.
C = The inverted value of the Z flag is set.

Execution example: ZCF
When the Z flag = 0, execution sets the carry flag to 1.



Appendix B Instruction Lists

- Explanation of symbols used in this document

1. Size

B	The operand size is in bytes (8 bits)
W	The operand size is in word (16 bits)
L	The operand size is in long word (32 bits)

2. Mnemonic

R	Eight general-purpose registers including 8/16/32-bit current bank registers. 8 bit register : W, A, B, C, D, E, H, L 16 bit register : WA, BC, DE, HL, IX, IY, IZ, SP 32 bit register : XWA, XBC, XDE, XHL, XIX, XIY, XIZ, XSP
r	8/16/32-bit general-parpose registers
cr	All 8/16/32-bit CPU control registers DMAS0 to 3, DMAD0 to 3, DMAC0 to 3, DMAM0 to 3, INTNEST
A	A register (8 bits)
F	Flag registers (8 bits)
F'	Inverse flag registers (8 bits)
SR	Status registers (16 bits)
PC	Program Counter (in minimum mode,16 bits; in maximum mode, 32 bits)
(mem)	8/16/32-bit memory data
mem	Effective address value
<W>	When the operand size is a word, "W" must be specified.
[ ]	Operands enclosed in square brackets can be omitted.
#	8/16/32-bit immediate data.
#3	3-bit immediate data: 0 to 7 or 1 to 8 ····· for abbreviated codes.
#4	4-bit immediate data: 0 to 15 or 1 to 16
d8	8-bit displacement: –80H to + 7FH
d16	16-bit displacement: –8000H to + 7FFFH
cc	Condition code
(#8)	Direct addressing : (00H) to (0FFH) ··· 256-byte area
(#16)	64K-byte area addressing : (0000H) to (0FFFFH)
$	A start address of the instruction is located

3. Cord

Z	The code crepresent the operand sizes. byte (8 bit) = 0 word (16 bit) = 2 long word (32 bit) = 4
ZZ	The code represent the operand sizes. byte (8 bit) = 00H word (16 bit) = 10H long word (32 bit) = 20H

4. Flag (SZHVNC)

–	Flag doesn't change.
*	Flag changes by executing instruction.
0	Flag is cleared to "0".
1	Flag is set to "1".
P	Flag changes by executing instruction (It works as parity flag).
V	Flag changes by executing instruction (It works as overflow flag).
X	An undefined value is set in flag.

5. Instruction length

Instruction length is represented in byte unit.

+#	adds immediate data length.
+M	adds addressing code length.
+#M	adds immediate data length and addressing code length.

6. State

Execution processing time of instruction are shown in order of 8 bit, 16 bit, 32 bit processing in status unit.

- Instruction Lists of 900/L (1/10)

(1) Load

Group	Size	Mnemonic	Codes (hex.)	Function	SZHVNC	Length (byte)	State
LD	BWL	LD R, r	C8 + zz + r : 88 + R	R ← r	------	2	4. 4. 4
	BWL	LD r, R	C8 + zz + r : 98 + R	r ← R	------	2	4. 4. 4
	BWL	LD r, #3	C8 + zz + r : A8 + #3	r ← #3	------	2	4. 4. 4
	BWL	LD R, #	20 + zz + R : #	R ← #	------	1 + #	2. 3. 5
	BWL	LD r, #	C8 + zz + r : 03 : #	r ← #	------	2 + #	4. 4. 6
	BWL	LD R, (mem)	80 + zz + mem : 20 + R	R ← (mem)	------	2 + M	4. 4. 6
	BWL	LD (mem), R	B0 + mem : 40 + zz + R	(mem) ← R	------	2 + M	4. 4. 6
	BW–	LD<W> (#8) , #	08 + z : #8 : #	(#8) ← #	------	2 + #	5. 6. –
	BW–	LD<W> (mem), #	B0 + mem : 00 + z : #	(mem) ← #	------	2 + M#	5. 6. –
	BW–	LD<W> (#16), (mem)	80 + zz + mem : 19 : #16	(#16) ← (mem)	------	4 + M	8. 8. –
	BW–	LD<W> (mem), (#16)	B0 + mem : 14 + z : #16	(mem) ← (#16)	------	4 + M	8. 8. –
PUSH	B—	PUSH F	18	(–XSP) ← F	------	1	3. –. –
	B–	PUSH A	14	(–XSP) ← A	------	1	3. –. –
	–WL	PUSH R	18 + zz + R	(–XSP) ← R	------	1	–. 3. 5
	BWL	PUSH r	C8 + zz + r : 04	(–XSP) ← r	------	2	5. 5. 7
	BW–	PUSH<W> #	09 + z : #	(–XSP) ← #	------	1 + #	4. 5. –
	BW–	PUSH<W> (mem)	80 + zz + mem : 04	(–XSP) ← (mem)	------	2 + M	7. 7. –
POP	B–	POP F	19	F ← (XSP+)	******	1	4. –. –
	B—	POP A	15	A ← (XSP+)	------	1	4. –. –
	–WL	POP R	38 + zz + R	R ← (XSP+)	------	1	–. 4. 6
	BWL	POP r	C8 + zz + r : 05	r ← (XSP+)	------	2	6. 6. 8
	BW–	POP<W> (mem)	B0 + mem : 04 + z	(mem) ← (XSP+)	------	2 + M	6. 6. –
LDA	–WL	LDA R, mem	B0 + mem : 10 + zz + R	R ← mem	------	2 + M	–. 4. 4
LDAR	–WL	LDAR R, $ + 4 + d16	F3:13:d16 : 20 + zz + R	R ← PC + d16	------	5	–. 11. 11

(2) Exchange

Group	Size	Mnemonic	Codes (hex.)	Function	SZHVNC	Length (byte)	State
EX	B—	EX F, F'	16	F ↔ F'	******	1	2. –. –
	BW–	EX R, r	C8 + zz + r : B8 + R	R ↔ r	------	2	5. 5. –
	BW–	EX (mem), R	80 + zz + mem : 30 + R	(mem) ↔ R	------	2 + M	6. 6. –
MIRR	–W–	MIRR r	D8 + r : 16	r<0:MSB> ← r<MSB : 0>	------	2	–. 4. –

- Instruction Lists of 900/L (2/10)

(3) Load/Increment/Decrement & Compare Increment/Decrement Size

Group	Size	Mnemonic	Codes (16 hex)	Function	SZHVNC	Length (byte)	State
LDxx	BW–	LDI<W> [(XDE+), (XHL+)]	83 + zz : 10	(XDE+) ← (XHL+) BC ← BC – 1	–-01.0–	2	10. 10. –
	BW–	LDI<W> (XIX+), (XIY+)	85 + zz : 10	(XIX+) ← (XIY+) BC ← BC – 1	–-01.0–	2	10. 10. –
	BW–	LDIR<W> [(XDE+), (XHL+)]	83 + zz : 11	repeat (XDE+) ← (XHL+) BC ← BC – 1 until BC = 0	–-000–	2	10. 10. – (end) 14. 14. – (repeat)
	BW–	LDIR<W> (XIX+), (XIY+)	85 + zz : 11	repeat (XIX+) ← (XIY+) BC ← BC – 1 until BC = 0	–-000–	2	10. 10. – (end) 14. 14. – (repeat)
	BW–	LDD<W> [(XDE–), (XHL–)]	83 + zz : 12	(XDE–) ← (XHL–) BC ← BC – 1	–-01.0–	2	10. 10. –
	BW–	LDD<W> (XIX–), (XIY–)	85 + zz : 12	(XIX–) ← (XIY–) BC ← BC – 1	–-01.0–	2	10. 10. –
	BW–	LDDR<W> [(XDE–), (XHL–)]	83 + zz : 13	repeat (XDE–) ← (XHL–) BC ← BC – 1 until BC = 0	–-000–	2	10. 10. – (end) 14. 14. – (repeat)
	BW–	LDDR<W> (XIX–), (XIY–)	85 + zz : 13	repeat (XIX–) ← (XIY–) BC ← BC – 1 until BC = 0	–-000–	2	10. 10. – (end) 14. 14. – (repeat)
CPxx	BW–	CPI [A/WA, (R+)]	80 + zz + R : 14	A/WA – (R+) BC ← BC – 1	*2)*1)1–	2	8. 8. –
	BW–	CPIR [A/WA, (R+)]	80 + zz + R : 15	repeat A/WA – (R+) BC ← BC – 1 until A/WA = (R) or BC=0	*2)*1)1–	2	10. 10. – (end) 14. 14. – (repeat)
	BW–	CPD [A/WA, (R–)]	80 + zz + R : 16	A/WA – (R–) BC ← BC – 1	*2)*1)1–	2	8. 8. –
	BW–	CPDR [A/WA, (R–)]	80 + zz + R : 17	Repeat A/WA – (R–) BC ← BC – 1 until A/WA = (R) or BC = 0	*2)*1)1–	2	10. 10. – (end) 14. 14. – (repeat)

Note 1: 1); If BC = 0 after execution, the P/V flag is set to 0, otherwise 1.
2); If A/WA = (R), the Z flag is set to 1, otherwise, 0 is set.

Note 2: When the operand is omitted in the CPI, CPIR, CPD, or CPDR instruction, A, (XHL+/–) is used as the default value.

- Instruction Lists of 900/L (3/10)

(4) Arithmetic Operations

Group	Size	Mnemonic		Codes (16 hex)		Function	SZHVNC	Length (byte)	State
ADD	BWL	ADD	R, r	C8 + zz + r	: 80 + R	R ← R + r	***V0*	2	4. 4. 7
	BWL	ADD	r, #	C8 + zz + r	: C8 : #	r ← r + #	***V0*	2 + #	4. 4. 7
	BWL	ADD	R, (mem)	80 + zz + mem	: 80 + R	R ← R + (mem)	***V0*	2 + M	4. 4. 6
	BWL	ADD	(mem), R	80 + zz + mem	: 88 + R	(mem) ← (mem) + R	***V0*	2 + M	6. 6. 10
	BW–	ADD<W>	(mem), #	80 + zz + mem	: 38 : #	(mem) ← (mem) + #	***V0*	2 + M#	7. 8. –
ADC	BWL	ADC	R, r	C8 + zz + r	: 90 + R	R ← R + r + CY	***V0*	2	4. 4. 7
	BWL	ADC	r, #	C8 + zz + r	: C9 : #	r ← r + # + CY	***V0*	2 + #	4. 4. 7
	BWL	ADC	R, (mem)	80 + zz + mem	: 90 + R	R ← R + (mem) + CY	***V0*	2 + M	4. 4. 6
	BWL	ADC	(mem), R	80 + zz + mem	: 98 + R	(mem) ← (mem) + R + CY	***V0*	2 + M	6. 6. 10
	BW–	ADC<W>	(mem), #	80 + zz + mem	: 39 : #	(mem) ← (mem) + # + CY	***V0*	2 + M#	7. 8. –
SUB	BWL	SUB	R, r	C8 + zz + r	: A0 + R	R ← R – r	***V1*	2	4. 4. 7
	BWL	SUB	r, #	C8 + zz +r	: CA : #	r ← r – #	***V1*	2 + #	4. 4. 7
	BWL	SUB	R, (mem)	80 + zz + mem	: A0 + R	R ← R – (mem)	***V1*	2 + M	4. 4. 6
	BWL	SUB	(mem), R	80 + zz + mem	: A8 + R	(mem) ← (mem) – R	***V1*	2 + M	6. 6. 10
	BW–	SUB<W>	(mem), #	80 + zz + mem	: 3A : #	(mem) ← (mem) – #	***V1*	2 + M#	7. 8. –
SBC	BWL	SBC	R, r	C8 + zz + r	: B0 + R	R ← R – r – CY	***V1*	2	4. 4. 7
	BWL	SBC	r, #	C8 + zz + r	: CB : #	r ← r – # – CY	***V1*	2 + #	4. 4. 7
	BWL	SBC	R, (mem)	80 + zz + mem	: B0 + R	R ← R – (mem) – CY	***V1*	2 + M	4. 4. 6
	BWL	SBC	(mem), R	80 + zz + mem	: B8 + R	(mem) ← (mem) – R – CY	***V1*	2 + M	6. 6. 10
	BW–	SBC<W>	(mem), #	80 + zz + mem	: 3B : #	(mem) ← (mem) – # – CY	***V1*	2 + M#	7. 8. –
CP	BWL	CP	R, r	C8 + zz + r	: F0 + R	R – r	***V1*	2	4. 4. 7
	BW–	CP	r, #3	C8 + zz + r	: D8 + #3	r – #3	***V1*	2	4. 4. –
	BWL	CP	r, #	C8 + zz + r	: CF : #	r – #	***V1*	2 + #	4. 4. 7
	BWL	CP	R, (mem)	80 + zz + mem	: F0 + R	R – (mem)	***V1*	2 + M	4. 4. 6
	BWL	CP	(mem), R	80 + zz + mem	: F8 + R	(mem) – R	***V1*	2 + M	6. 6. 6
	BW–	CP<W>	(mem), #	80 + zz + mem	: 3F : #	(mem) – #		2 + M#	6. 6. –
INC	B––	INC	#3, r	C8 + r	: 60 + #3	r ← r + #3	***V0–	2	4. –. –
	–WL	INC	#3, r	C8 + zz + r	: 60 + #3	r ← r + #3	------	2	–. 4. 4
	BW–	INC<W>	#3, (mem)	80 + zz + mem	: 60 + #3	(mem) ← (mem) + #3	***V0–	2 + M	6. 6. –
DEC	B––	DEC	#3, r	C8 + r	: 68 + #3	r ← r – #3	***V1–	2	4. –. –
	–WL	DEC	#3, r	C8 + zz + r	: 68 + #3	r ← r – #3	------	2	–. 4. 5
	BW–	DEC<W>	#3, (mem)	80 + zz + mem	: 68 + #3	(mem) ← (mem) – #3	***V1–	2 + M	6. 6. –
NEG	BW–	NEG	r	C8 + zz + r	: 07	r ← 0 – r	***V1*	2	5. 5. –
EXTZ	–WL	EXTZ	r	C8 + zz + r	: 12	r<high> ← 0	------	2	–. 4. 4
EXTS	–WL	EXTS	r	C8 + zz + r	: 13	r<high> ← r<low. MSB>	------	2	–. 5. 5
DAA	B––	DAA	r	C8 + r	: 10	Decimal adjustment after addition or subtraction	***P–*	2	6. –. –
PAA	–WL	PAA	r	C8 + zz + r	: 14	if r<0> = 1 then INC r	------	2	–. 4. 4

Note 1: With the INC/DEC instruction, when the code value of #3 = 0, functions as +8/–8.

Note 2: When the ADD R, r (word type) instruction is used in the TLCS-90, the S, Z, and V flags do not change. In the TLCS-900, these flags change.

- Instruction Lists of 900/L (4/10)

Group	Size	Mnemonic	Codes (hex.)	Function	SZHVNC	Length (byte)	State
MUL	BW–	MUL RR, r	C8 + zz + r : 40 + R	RR ← R × r	------	2	18. 26. –
	BW–	MUL rr, #	C8 + zz + r : 08 : #	rr ← r × #	------	2 + #	18. 26. –
	BW–	MUL RR, (mem)	80 + zz + mem : 40 + R	RR ← R × (mem)	------	2 + M	18. 26. –
MULS	BW–	MULS RR, r	C8 + zz + r : 48 + R	RR ← R × r ;signed	------	2	18. 26. –
	BW–	MULS rr, #	C8 + zz + r : 09 : #	rr ← r × # ;signed	------	2 + #	18. 26. –
	BW–	MULS RR, (mem)	80 + zz + mem : 48 + R	RR ←R × (mem);signed	------	2 + M	18. 26. –
DIV	BW–	DIV RR, r	C8 + zz + r : 50 + R	R ← RR ÷ r	--V--	2	22. 30. –
	BW–	DIV rr, #	C8 + zz + r : 0A : #	r ← rr ÷ #	--V--	2 + #	22. 30. –
	BW–	DIV RR, (mem)	80 + zz + mem : 50 + R	R ← RR ÷ (mem)	--V--	2 + M	22. 30. –
DIVS	BW–	DIVS RR, r	C8 + zz + r : 58 + R	R ← RR ÷ r ;signed	--V--	2	24. 32. –
	BW–	DIVS rr, #	C8 + zz + r : 0B : #	r ← rr ÷ # ;signed	--V--	2 + #	24. 32. –
	BW–	DIVS RR, (mem)	80 + zz + mem : 58 + R	R ← RR ÷ (mem);signed	--V--	2 + M	24. 32. –
MULA	–W–	MULA rr	D8 + r : 19	Multiply and add signed rr ← rr + (XDE) × (XHL) 32 bit 32 bit 16 bit 16 bit XHL ← XHL–2	**–V--	2	–.31. –
MINC	–W–	MINC1 #, r (# = 2**n) (1< = n< = 15)	D8 + r : 38 : # – 1	modulo increment ;+1 if (r mod #) = (# – 1) then r ← r – (# – 1) else r ← r + 1	------	4	–. 8. –
	–W–	MINC2 #, r (# = 2**n) (2< = n< = 15)	D8 + r : 39 : # – 2	modulo increment ;+2 if (r mod #) = (# – 2) then r ← r – (# – 2) else r ← r + 2	------	4	–. 8. –
	–W–	MINC4 #, r (# = 2**n) (3< = n< = 15)	D8 + r : 3A : # – 4	modulo increment ;+4 if (r mod #) = (# – 4) then r ← r – (# – 4) else r ← r + 4	------	4	–. 8. –
MDEC	–W–	MDEC1 #, r (# = 2**n) (1< = n< = 15)	D8 + r : 3C : # – 1	modulo decrement ;–1 if (r mod #) = 0 then r ← r + (# – 1) else r ← r – 1	------	4	–. 7. –
	–W–	MDEC2 #, r (# = 2**n) (2< = n< = 15)	D8 + r : 3D : # – 2	modulo decrement ;–2 if (r mod #) = 0 then r ← r + (# – 2) else r ← r – 2	------	4	–. 7. –
	–W–	MDEC4 #, r (# = 2**n) (3< = n< = 15)	D8 + r : 3E : # – 4	modulo decrement ;–4 if (r mod #) = 0 then r ← r + (# – 4) else r ← r – 4	------	4	–. 7. –

Note: Operand RR of the MUL, MULS, DIV, and DIVS instructions indicates that a register twice the size of the operation is specified. When the operation is in bytes (8 bits × 8 bits, 16/8 bits), word register (16 bits) is specified; when the operation is in words (16 bits × 16 bits, 32/16 bits), long word register (32 bits) is specified.

- Instruction Lists of 900/L (5/10)

(5) Logical operations

Group	Size	Mnemonic	Codes (hex.)	Function	SZHVNC	Length (byte)	State
AND	BWL	AND R, r	C8 + zz + r : C0 + R	R ← R and r	**1P00	2	4. 4. 7
	BWL	AND r, #	C8 + zz + r : CC : #	r ← r and #	**1P00	2 + #	4. 4. 7
	BWL	AND R, (mem)	80 + zz + mem : C0 + R	R ← R and (mem)	**1P00	2 + M	4. 4. 6
	BWL	AND (mem), R	80 + zz + mem : C8 + R	(mem) ← (mem) and R	**1P00	2 + M	6. 6. 10
	BW–	AND<w> (mem), #	80 + zz + mem : 3C : #	(mem) ← (mem) and #	**1P00	2 + M#	7. 8. -
OR	BWL	OR R, r	C8 + zz + r : E0 + R	R ← R or r	**0P00	2	4. 4. 7
	BWL	OR r, #	C8 + zz + r : CE : #	r ← r or #	**0P00	2 + #	4. 4. 7
	BWL	OR R, (mem)	80 + zz + mem : E0 + R	R ← R or (mem)	**0P00	2 + M	4. 4. 6
	BWL	OR (mem), R	80 + zz + mem : E8 + R	(mem) ← (mem) or R	**0P00	2 + M	6. 6. 10
	BW–	OR<W> (mem), #	80 + zz + mem : 3E : #	(mem) ← (mem) or #	**0P00	2 + M#	7. 8. –
XOR	BWL	XOR R, r	C8 + zz + r : D0 + R	R ← R xor r	**0P00	2	4. 4. 7
	BWL	XOR r, #	C8 + zz + r : CD : #	r ← r xor #	**0P00	2 + #	4. 4. 7
	BWL	XOR R, (mem)	80 + zz + mem : D0 + R	R ← R xor (mem)	**0P00	2 + M	4. 4. 6
	BWL	XOR (mem), R	80 + zz + mem : D8 + R	(mem) ← (mem) xor R	**0P00	2 + M	6. 6. 10
	BW–	XOR<W> (mem), #	80 + zz + mem : 3D : #	(mem) ← (mem) xor #	**0P00	2 + M#	7. 8. –
CPL	BW–	CPL r	C8 + zz + r : 06	r ← not r	—1–1–	2	4. 4. –

- Instruction Lists of 900/L (6/10)

(6) Bit operations

Group	Size	Mnemonic	Codes (hex.)	Function	SZHVNC	Length (byte)	State
LDCF	BW–	LDCF #4, r	C8 + zz + r : 23 : #4	CY ← r<#4>	———*	3	4. 4. –
	BW–	LDCF A , r	C8 + zz + r : 2B	CY ← r<A>	———*	2	4. 4. –
	B—	LDCF #3, (mem)	B0 + men : 98 + #3	CY ← (mem)<#3>	———*	2 + M	8. –. –
	B—	LDCF A , (mem)	B0 + men : 2B	CY ← (mem)<A>	———*	2 + M	8. –. –
STCF	BW–	STCF #4, r	C8 + zz + r : 24 : #4	r<#4> ← CY	———	3	4. 4. –
	BW–	STCF A , r	C8 + zz + r : 2C	r<A> ← CY	———	2	4. 4. –
	B—	STCF #3, (mem)	B0 + mem : A0 + #3	(mem)<#3> ← CY	———	2 + M	8. –. –
	B—	STCF A , (mem)	B0 + mem : 2C	(mem)<A> ← CY	———	2 + M	8. –. –
ANDCF	BW–	ANDCF #4, r	C8 + zz + r : 20 : #4	CY ← CY and r<#4>	———*	3	4. 4. –
	BW–	ANDCF A , r	C8 + zz + r : 28	CY ← CY and r<A>	———*	2	4. 4. –
	B—	ANDCF #3, (mem)	B0 + mem : 80 + #3	CY ← CY and (mem)<#3>	———*	2 + M	8. –. –
	B—	ANDCF A , (mem)	B0 + mem : 28	CY ← CY and (mem)<A>	———*	2 + M	8. –. –
ORCF	BW–	ORCF #4, r	C8 + zz + r : 21 : #4	CY ← CY or r<#4>	———*	3	4. 4. –
	BW–	ORCF A , r	C8 + zz + r : 29	CY ← CY or r<A>	———*	2	4. 4. –
	B—	ORCF #3, (mem)	B0 + mem : 88 + #3	CY ← CY or (mem)<#3>	———*	2 + M	8. –. –
	B—	ORCF A , (mem)	B0 + mem : 29	CY ← CY or (mem)<A>	———*	2 + M	8. –. –
XORCF	BW–	XORCF #4, r	C8 + zz + r : 22 : #4	CY ← CY xor r<#4>	———*	3	4. 4. –
	BW–	XORCF A , r	C8 + zz + r : 2A	CY ← CY xor r<A>	———*	2	4. 4. –
	B—	XORCF #3, (mem)	B0 + mem : 90 + #3	CY ← CY xor (mem)<#3>	———*	2 + M	8. –. –
	B—	XORCF A , (mem)	B0 + mem : 2A	CY ← CY xor (mem)<A>	———*	2 + M	8. –. –
RCF	—	RCF	10	CY ← 0	—0–00	1	2
SCF	—	SCF	11	CY ← 1	—0–01	1	2
CCF	—	CCF	12	CY ← not CY	—X–0*	1	2
ZCF	—	ZCF	13	CY ← not Z flag	—X–0*	1	2
BIT	BW–	BIT #4, r	C8 + zz + r : 33 : #4	Z ← not r<#4>	X*1X0–	3	4. 4. –
	B—	BIT #3, (mem)	B0 + mem : C8 + #3	Z ← not (mem)<#3>	X*1X0–	2 + M	8. –. –
RES	BW–	RES #4, r	C8 + zz + r : 30 : #4	r<#4> ← 0	———	3	4. 4. –
	B—	RES #3, (mem)	B0 + mem : B0 + #3	(mem)<#3> ← 0	———	2 + M	8. –. –
SET	BW–	SET #4, r	C8 + zz + r : 31 : #4	r<#4> ← 1	———	3	4. 4. –
	B—	SET #3, (mem)	B0 + mem : B8 + #3	(mem)<#3> ← 1	———	2 + M	8. –. –
CHG	BW–	CHG #4, r	C8 + zz + r : 32 : #4	r<#4> ← not r<#4>	———	3	4. 4. –
	B—	CHG #3, (mem)	B0 + mem : C0 + #3	(mem)<#3> ← not (mem)<#3>	———	2 + M	8. –. –
TSET	BW–	TSET #4, r	C8 + zz + r : 34 : #4	Z ← not r<#4> : r<#4> ← 1	X*1X0–	3	6. 6. –
	B—	TSET #3, (mem)	B0 + mem : A8 + #3	Z ← not (mem)<#3> (mem)<#3> ← 1	X*1X0–	2 + M	10. –. –
BS1	–W–	BS1F A, r	D8 + r : 0E	A ← 1 search r ; Forward	—–1.—	2	–. 4. –
	–W–	BS1B A, r	D8 + r : 0F	A ← 1 search r ; Backward	—–1.—	2	–. 4. –

Note: 1) ; 0 is set when the bit searched for is found, otherwise 1 is set and an undefined value is set in the A register.

- Instruction Lists of 900/L (7/10)

(7) Special operations and CPU control

Group	Size	Mnemonic	Codes (hex.)	Function	SZHVNC	Length (byte)	State
NOP	–––	NOP	00	no operation	–––––	1	2
MIN	–––	MIN	00	Changes to minimum mode. MAX←0	–––––	1	4
EI	–––	EI [#3]	06 : #3	Sets interrupt enable flag. IFF ← #3	–––––	2	5
DI	–––	DI	06 : 07	Disables interrupt. IFF ← 7	–––––	2	5
PUSH	–W–	PUSH SR	02	(–XSP) ← SR	–––––	1	–. 4. –
POP	–W–	POP SR	03	SR ← (XSP+)	******	1	–. 6. –
SWI	–––	SWI [#3]	F8 + #3	Software interrupt PUSH PC&SR JP (8000H + 4H × #3)	–––––	1	16
HALT	–––	HALT	05	CPU halt	–––––	1	8
LDC	BWL BWL	LDC cr, r LDC r, cr	C8 + zz + r : 2E : cr C8 + zz + r : 2F : cr	cr ← r r ← cr	––––– –––––	3 3	8. 8. 8 8. 8. 8
LDX	B––	LDX (#8), #	F7:00 : #8 : 00 : # : 00	(#8) ← #	–––––	6	9. –. –

LINK	––L	LINK r, d16	E8 + r : 0C : d16	PUSH r
				LD r, XSP
				ADD XSP, d16
					–––––	4	–. –. 10

UNLK	––L	UNLK r	E8 + r : 0D	LD XSP, r POP r	–––––	2	–. –. 8
LDF	–––	LDF #3	17 : #3	Sets register bank. RFP ← #3 (0 at reset)	–––––	2	2
INCF	–––	INCF	0C	Switches register banks. RFP ← RFP + 1	–––––	1	2
DECF	–––	DECF	0D	Switches register banks. RFP ← RFP – 1	–––––	1	2
SCC	BW–	SCC cc, r	C8 + zz + r : 70 + cc	if cc then r ← 1 else r ← 0	–––––	2	6. –6. –

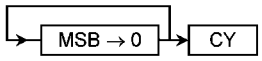
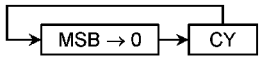
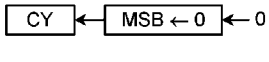
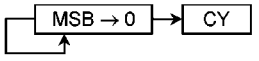
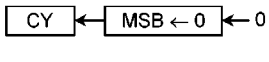
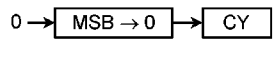
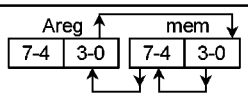
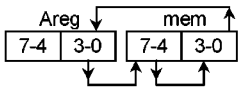
Note 1: When operand #3 coding in the EI instruction is omitted, 0 is used as the default value.

Note 2: When operand #3 coding in the SWI instruction is omitted, 7 is used as the default value.

Note 3: The value in the state column for the SWI instruction represents the number of states when the CPU is in maximum mode. In minimum mode, subtract - 2.

- Instruction Lists of 900/L1(8/10)

(8) Rotate and Shift

Group	Size	Mnemonic		Codes (hex.)		Function	SZHVNC	Length (byte)	State
RLC	BWL	RLC	#4, r	C8 + zz + r	: E8 : #4		**0P0*	3	6.6.8+2n
	BWL	RLC	A, r	C8 + zz + r	: F8		**0P0*	2	6.6.8+2n
	BW–	RLC<W>	(mem)	80 + zz + mem	: 78		**0P0*	2 + M	8.8.–
RRC	BWL	RRC	#4, r	C8 + zz + r	: E9 : #4		**0P0*	3	6.6.8+2n
	BWL	RRC	A, r	C8 + zz + r	: F9		**0P0*	2	6.6.8+2n
	BW–	RRC<W>	(mem)	80 + zz + mem	: 79		**0P0*	2 + M	8.8.–
RL	BWL	RL	#4, r	C8 + zz + r	: EA : #4		**0P0*	3	6.6.8+2n
	BWL	RL	A, r	C8 + zz + r	: FA		**0P0*	2	6.6.8+2n
	BW–	RL<W>	(mem)	80 + zz + mem	: 7A		**0P0*	2 + M	8.8.–
RR	BWL	RR	#4, r	C8 + zz + r	: EB : #4		**0P0*	3	6.6.8+2n
	BWL	RR	A, r	C8 + zz + r	: FB		**0P0*	2	6.6.8+2n
	BW–	RR<W>	(mem)	80 + zz + mem	: 7B		**0P0*	2 + M	8.8.–
SLA	BWL	SLA	#4, r	C8 + zz + r	: EC : #4		**0P0*	3	6.6.8+2n
	BWL	SLA	A, r	C8 + zz + r	: FC		**0P0*	2	6.6.8+2n
	BW–	SLA<W>	(mem)	80 + zz + mem	: 7C		**0P0*	2 + M	8.8.–
SRA	BWL	SRA	#4, r	C8 + zz + r	: ED : #4		**0P0*	3	6.6.8+2n
	BWL	SRA	A, r	C8 + zz + r	: FD		**0P0*	2	6.6.8+2n
	BW–	SRA<W>	(mem)	80 + zz + mem	: 7D		**0P0*	2 + M	8.8.–
SLL	BWL	SLL	#4, r	C8 + zz + r	: EE : #4		**0P0*	3	6.6.8+2n
	BWL	SLL	A, r	C8 + zz + r	: FE		**0P0*	2	6.6.8+2n
	BW–	SLL<W>	(mem)	80 + zz + mem	: 7E		**0P0*	2 + M	8.8.–
SRL	BWL	SRL	#4, r	C8 + zz + r	: EF : #4		**0P0*	3	6.6.8+2n
	BWL	SRL	A, r	C8 + zz + r	: FF		**0P0*	2	6.6.8+2n
	BW–	SRL<W>	(mem)	80 + zz + mem	: 7F		**0P0*	2 + M	8.8.–
RLD	B––	RLD	[A,](mem)	80 + mem	: 06		**0P0–	2 + M	12. –. –
RRD	B––	RRD	[A,](mem)	80 + mem	: 07		**0P0–	2 + M	12. –. –

Note 1: When #4/A is used to specify the number of shifts, the values of lower 4 bits (0 to 15) is used. Code 0 means 16 shifts.

Note 2: When the following instructions are used in the TLCS-90, the S, Z and V flags do not change.
RLCA, RRCA, RLA, RRA, SLAA, SRAA, SLLA, and SRLA
In the TLCS-900, these flags change.

- Instruction Lists of 900/L (9/10)

(9) Jump, Call and Return

Group	Size	Mnemonic	Codes (hex.)	Function	SZHVNC	Length (byte)	State
JP	–––	JP #16	1A : #16	PC ← #16	–––––––	3	7
	–––	JP #24	1B : #24	PC ← #24	–––––––	4	7
	–––	JR [cc,]$ + 2 + d8	60 + cc : d8	if cc then PC ← PC + d8	–––––––	2	8/4 (T/F)
	–––	JRL [cc,]$ + 3 + d16	70 + cc : d16	if cc then PC ← PC + d16	–––––––	3	8/4 (T/F)
	–––	JP [cc,]mem	B0 + mem : DO + cc	if cc then PC ← mem	–––––––	2 + M	9/6 (T/F)
CALL	–––	CALL #16	1C : #16	PUSH PC : JP #16	–––––––	3	12
	–––	CALL #24	1D : #24	PUSH PC : JP #24	–––––––	4	12
	–––	CALR $ + 3 + d16	1E : d16	PUSH PC : JR $ + 3 + d16	–––––––	3	12
	–––	CALL [cc,]mem	B0 + mem : E0 + cc	if cc then PUSH PC : JP mem	–––––––	2 + M	12/6(T/F)
DJNZ	BW–	DJNZ [r,]$ + 3/4 + d8	C8 + zz + r : 1C : d8	r ← r – 1 if r ≠ 0 then JR $ + 3 + d8	–––––––	3	11(r ≠ 0) 7(r = 0)
RET	–––	RET	0E	POP PC	–––––––	1	9
	–––	RET cc	B0 : F0 + cc	if cc then POP PC	–––––––	2	12/6(T/F)
	–––	RETD d16	0F : d16	RET : ADD XSP, d16	–––––––	3	9
	–––	RETI	07	POP SR&PC	******	1	12

Note 1: The value in the state column for the CALL, CALR, RET, RETD, and RETI instructions represents the number of states when the CPU is in maximum mode. In minimum mode, subtract –2.

Note 2: (T/F) represents the number of states at true/false.

- Instruction Lists of 900/L (10/10)

(10) Addressing mode

type	mode	State (addition)
R	R	+0
r	r	+1
(mem)	(R)	+0
	(R + d8)	+2
	(#8)	+2
	(#16)	+2
	(#24)	+3
	(r)	+5
	(r + d16)	+5
	(r + r8)	+8
	(r + r16)	+8
	(–r)	+3
	(r+)	+3

(11) Interrupt

mode		operation	state
General-purpose interrupt processing		PUSH PC PUSH SR IFF ← accepted level + 1 INTNEST ← INTNEST + 1 JP (FFFF00H + vector)	25 (MAX mode) 23 (MIN mode)
μDMA	I/O to MEM	(DMADn+) ← (DMASn)	16. 16. –
	I/O to MEM	(DMADn–) ← (DMASn)	16. 16. –
	MEM to I/O	(DMADn) ← (DMASn+)	16. 16. –
	MEM to I/O	(DMADn) ← (DMASn–)	16. 16. –
	I/O to I/O	(DMADn) ← (DMASn)	16. 16. –
	DRAM Refresh	Dummy ← (DMASn+)	–. –. 14
	Counter	DMASn ← DMASn + 1	–. –. 11

Note: For details of interrupt processing, refer to Chapter4 “3.4 Interrupts”.

Appendix C Instruction Code Maps (1/4)

1-byte op code instructions

H/L	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0	NOP		PUSH SR	POP SR	MIN	HALT	EI n	RETI	LD (n) , n	PUSH n	LDW (n) , nn	PUSHW nn	INCF	DECF	RET	RETD dd
1	RCF	SCF	CCF	ZCF	PUSH A	POP A	EX F, F'	LDF n	PUSH F	POP F	JP nn	JP nnn	CALL nn	CALL nnn	CALR PC+dd	
2	LD R, n								PUSH RR							
3	LD RR, nn								PUSH XRR							
4	LD XRR, nnnn								POP RR							
5									POP XRR							
6	F	LT	LE	ULE	PE/OV	M/MI	Z	JR C	cc,PC + d (T)	GE	GT	UGT	PO/NOV	P/PL	NZ	NC
7	F	LT	LE	ULE	PE/OV	M/MI	Z	JRL C	cc,PC + dd (T)	GE	GT	UGT	PO/NOV	P/PL	NZ	NC
8	scr. B (XWA)	(XBC)	(XDE)	(XHL)	(XIX)	(XIY)	(XIZ)	(XSP)	scr. B (XWA +d)	(XBC +d)	(XDE +d)	(XHL +d)	(XIX +d)	(XIY +d)	(XIZ +d)	(XSP +d)
9	scr. W (XWA)	(XBC)	(XDE)	(XHL)	(XIX)	(XIY)	(XIZ)	(XSP)	scr. W (XWA +d)	(XBC +d)	(XDE +d)	(XHL +d)	(XIX +d)	(XIY +d)	(XIZ +d)	(XSP +d)
A	scr. L (XWA)	(XBC)	(XDE)	(XHL)	(XIX)	(XIY)	(XIZ)	(XSP)	scr. L (XWA +d)	(XBC +d)	(XDE +d)	(XHL +d)	(XIX +d)	(XIY +d)	(XIZ +d)	(XSP +d)
B	dst (XWA)	(XBC)	(XDE)	(XHL)	(XIX)	(XIY)	(XIZ)	(XSP)	dst (XWA +d)	(XBC +d)	(XDE +d)	(XHL +d)	(XIX +d)	(XIY +d)	(XIZ +d)	(XSP +d)
C	scr. B (n)	(nn)	(nnn)	(mem)	(–xrr)	(xrr+)		reg. B r	reg. B W	A	B	C	D	E	H	L
D	scr. W (n)	(nn)	(nnn)	(mem)	(–xrr)	(xrr+)		reg. W rr	reg. W WA	BC	DE	HL	IX	IY	IZ	SP
E	scr. L (n)	(nn)	(nnn)	(mem)	(–xrr)	(xrr+)		reg. L xrr	reg. L XWA	XBC	XDE	XHL	XIX	XIY	XIZ	XSP
F	dst (n)	(nn)	(nnn)	(mem)	(–xrr)	(xrr+)		LDX (n), n	SWI n 0	1	2	3	4	5	6	7

Note 1: Codes in blank parts are undefined instructions (i.e., illegal instructions).

Note 2: Dummy instructions are assigned to code 01H. Do not use it.

Appendix C Instruction Code Maps (2/4)

1st byte: reg

H/L	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0				LD r,#	PUSH r	POP r	CPL BW r	NEG BW r	MUL rr,#	MULS rr,#	DIV rr,#	DIVS BW rr,#	LINK L r, dd	UNLK L r	BS1F A, r	BS1B W A, r
1	DAA B r		EXTZ WL r	EXTS WL r	PAA WL r		MIRR W r			MULA W r			DJNZ BW r, d		LDC cr, r	LDC r, cr
2	ANDCF #, r	ORCF #, r	XORCF #, r	LDCF #, r	STCF BW #, r				ANDCF A, r	ORCF A, r	XORCF A, r	LDCF A, r	STCF BW A, r		LDC cr, r	LDC r, cr
3	RES #, r	SET #, r	CHG #, r	BIT #, r	TSET BW #, r				MINC1	MINC2 #, r	MINC4 W		MDEC1	MDEC2 #, r	MDEC4 W	
4				MUL	R, r			BW				MULS	R, r			BW
5				DIV	R, r			BW				DIVS	R, r			BW
6	8	1	2	INC 3	#3, r 4	5	6	7	8	1	2	DEC 3	#3, r 4	5	6	7
7	F	LT	LE	ULE	PE/OV	M/MI	Z	SCC C	cc, r (T)	GE	GT	UGT	PO/NOV	P/PL	NZ	BW NC
8				ADD	R, r							LD	R, r			
9				ADC	R, r							LD	r, R			
A				SUB	R, r				0	1	2	LD 3	r, #3 4	5	6	7
B				SBC	R, r							EX	R, r			BW
C				AND	R, r				ADD r, #	ADC r, #	SUB r, #	SBC r, #	AND r, #	XOR r, #	OR r, #	CP r, #
D				XOR	R, r				0	1	2	CP 3	r, #3 4	5	6	BW 7
E				OR	R, r				RLC #, r	RRC #, r	RL #, r	RR #, r	SLA #, r	SRA #, r	SLL #, r	SRL #, r
F				CP	R, r				RLC A, r	RRC A, r	RL A, r	RR A, r	SLA A, r	SRA A, r	SLL A, r	SRL A, r

r: Register specified by the 1st byte code. (Any CPU registers can be specified.)

R: Register specified by the 2nd byte code. (Only eight current registers can be specified.)

B: Operand size is a byte.

W: Operand size is a word.

L: Operand size is a long word.

Note: Dummy instructions are assigned to codes 1AH, 1BH, 3BH, and 3FH. Do not use them.

Appendix C Instruction Code Maps (3/4)

1st byte: src (mem)

H/L	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0					PUSH BW (mem)		RLD A,	RLD B (mem)								
1	LDI	LDIR	LDD	LDDR BW	CPI	CPIR	CPD	CPDR BW		LD BW (nn), (m)						
2	W	A	B	LD C	R,(mem) D	E	H	L								
3				EX	(mem),R			BW	ADD	ADC	SUB	SBC (mem), #	AND	XOR	OR	CP BW
4				MUL	R,(mem)			BW				MULS	R,(mem)			BW
5				DIV	R,(mem)			BW				DIVS	R,(mem)			BW
6	8	1	2	INC 3	#3, (mem) 4	5	6	BW 7	8	1	2	DEC 3	#3, (mem) 4	5	6	BW 7
7									RLC	RRC	RL	RR (mem)	SLA	SRA	SLL	SRL BW
8				ADD	R,(mem)							ADD	(mem),R			
9				ADC	R,(mem)							ADC	(mem),R			
A				SUB	R,(mem)							SUB	(mem),R			
B				SBC	R,(mem)							SBC	(mem),R			
C				AND	R,(mem)							AND	(mem),R			
D				XOR	R,(mem)							XOR	(mem),R			
E				OR	R,(mem)							OR	(mem),R			
F				CP	R,(mem)							CP	(mem),R			

B: Operand size is a byte.

W: Operand size is a word.

Appendix C Instruction Code Maps (4/4)

1st byte: dst (mem)

H/L	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0	LD B (m), #		LD W (m), #		POP B (mem)		POP W (mem)									
1					LD B (m), (nn)		LD W (m), (nn)									
2	LDA R,(mem) W								ANDC A, (mem)	ORCF A, (mem)	XORCF A, (mem)	LDCF A, (mem)	STCF B A, (mem)			
3	LDA R,(mem) L															
4	LD (mem),R B															
5	LD (mem),R W															
6	LD (mem),R L															
7																
8	ANDCF #3, (mem) B								ORCF #3, (mem) B							
	0	1	2	3	4	5	6	7	0	1	2	3	4	5	6	7
9	XORCF #3, (mem) B								LDCF #3, (mem) B							
	0	1	2	3	4	5	6	7	0	1	2	3	4	5	6	7
A	STCF #3, (mem) B								TSET #3, (mem) B							
	0	1	2	3	4	5	6	7	0	1	2	3	4	5	6	7
B	RES #3, (mem) B								SET #3, (mem) B							
	0	1	2	3	4	5	6	7	0	1	2	3	4	5	6	7
C	CHG #3, (mem) B								BIT #3, (mem) B							
	0	1	2	3	4	5	6	7	0	1	2	3	4	5	6	7
D								JP cc, mem								
	F	LT	LE	ULE	PE/OV	M/MI	Z	C	(T)	GE	GT	UGT	PO/NOV	P/PL	NZ	NC
E								CALL cc, mem								
	F	LT	LE	ULE	PE/OV	M/MI	Z	C	(T)	GE	GT	UGT	PO/NOV	P/PL	NZ	NC
F								RET cc	(1st byte code is BOH.)							
	F	LT	LE	ULE	PE/OV	M/MI	Z	C	(T)	GE	GT	UGT	PO/NOV	P/PL	NZ	NC

B: Operand size is a byte.

W: Operand size is a word.

L: Operand size is a long word.

Appendix D Differences between TLCS-90 and TLCS-900 Series

Item \ Series	TLCS-90	TLCS-900/L
CPU architecture	8-bit CPU	16-bit CPU
Built-in ROM/built-in RAM	8-bit data bus	16-bit data bus
Built-in I/O	8-bit data bus	8-bit data bus
External data bus	8-bit data bus	8-bit/16-bit data bus (can be mixed)
Program space (except devices with MMU)	64 KB	16MB (linear)
Data space	16 MB (bank)	16 MB (linear)
Instruction set/instruction mnemonic	TLCS-90	TLCS-90 + α α = enhancement of 16-bit multiply / divide instructions and bit operation instruction. 32-bit load/operation instructions, C compiler instructions, register bank operation instructions, etc.
Instruction code (object code)	Unique to TLCS-90	Unique to TLCS-900 (Different from TLCS-90.)
Addressing mode	TLCS-90	TLCS-90 + α α = (–Reg), (Reg+), (Reg + disp16), (Reg + Reg16), (nnn)
General-purpose register	TLCS-90	TLCS-90 + α α = Uses as 32 bits and register bank, and adds a system stack pointer.
Flag (F)	S Z I H X V N C	S Z "0" H "0" V N C I flag is extended to IFF2 to 0 of status register.X flag is deleted.
Reset	PC ← 0000H (SP does not change.)	PC ← (Vector base address) XSP ← 100H
Built-in ROM address	0000H to	undefined
Built-in RAM address	to FFxxH	0080H to
Built-in I/O address	FFxxH to FFFFH	0000H to 007FH
Direct addressing area (n)	FF00H to FFFFH	0000H to 00FFH
Interrupt		
Interrupt start address	0000H + (8 × V)	Vector base address + 4 × V
Register to be saved	PC & AF	PC & SR
Mask register	IFF	IFF2 to 0
Mask level	0 to 1	0 to 7

Item \ Series	TLCS-90	TLCS-900/L
Instruction		
(1). ADD R, r (word type)	S/Z/V flags don't change. [S/Z/V flag changes expect add 16 bit register.]	S/Z/V flag changes.
(2). Shift of A register	[RLCA RRCA RLA RRA SLAA SRAA SLLA SRLA] S/Z/V flags don't change in these instruction. [RLC A RRC A RL A RR A SLA A SRA A SLL A SRL A] S/Z/V flag changes in these instruction.	S/Z/V flag changes.

Note: The TLCS-900/L is essentially the same as the TLCS-90 but with a 16-bit CPU.
However, six types of instructions used in the TLCS-90 do not directly correspond with those used in the TLCS-900/L. Thus, when transfering programs designed for the TLCS-90 to the TLCS-900/L, replace them with equivalents as follows:

Instructions in TLCS-90 but not in TLCS-900/L	Equivalent instructions in TLCS-900/L
EXX	EX BC, BC' EX DE, DE' EX HL, HL'
EX AF, AF'	EX A, A' EX F, F'
PUSH AF	PUSH A PUSH F
POP AF	POP F POP A
INCX	(32-bit INC instruction)
DECX	(32-bit DEC instruction)

Some TLCS-900/L instructions, though basically the same as TLCS-90 instructions, have more functions and more specification items in their operands. They are listed below.

TLCS-90	TLCS-900
INC reg INC mem	INC imm3, reg INC imm3, mem
DEC reg DEC mem	DEC imm3, reg DEC imm3, mem
RLC reg RRC reg RL reg RR reg SLA reg SRA reg SLL reg SRL reg	RLC imm, reg RRC imm, reg RL imm, reg RR imm, reg SLA imm, reg SRA imm, reg SLL imm, reg SRL imm, reg